How to adopt a logic

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Abstract

What is commonly referred to as the Adoption Problem is a challenge to the idea that the principles of logic can be rationally revised. The argument is based on a reconstruction of unpublished work by Saul Kripke. As the reconstruction has it, Kripke extends the scope of Willard van Orman Quine’s regress argument against conventionalism to the possibility of adopting new logical principles. In this paper we want to discuss the scope of this challenge. Are all revisions of logic subject to the Adoption Problem? If not, are there significant cases of logical revision that are subject to the Adoption Problem? We will argue that both questions should be answered negatively.

Keywords— Adoption Problem, Logical Pluralism, Nonclassical Logic, Epistemology of Logic, Logical Revisionism

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What is commonly referred to as the Adoption Problem\footnote{This label for the problem is due to Padro (2015).} is often considered a challenge to the idea that the principles for logic can be rationally revised. The argument is based on a reconstruction of unpublished work by Saul Kripke.\footnote{See Stairs (2006), Padro (2015), Finn (2019a), Finn (2019b), Kripke (2020) and Devitt and Roberts (npub). Since the basis of this discussion is an unpublished manuscript that is not authorized, we decided to refer to it in the following way: when providing textual evidence for Kripke’s views, we quote from papers that are published and directly quote Kripke; in cases in which we want to give credit to Kripke for an observation or argument, we refer to (Kripke, 2020).} As the reconstruction has it, Kripke essentially extends the scope of Willard Van Orman Quine’s regress argument (Quine, 1976) against conventionalism to the possibility of adopting new logical principles or rules. According to the reconstruction, the Adoption Problem is that new logical rules cannot be adopted unless one already can infer with these rules, in which case the adoption of the rules is unnecessary (Padro, 2015, 18).

In this paper we want to discuss the scope of this challenge. Are all revisions of logic subject to the Adoption Problem? If not, are there significant cases of logical revision that are subject to the Adoption Problem? We will argue that both questions should be answered negatively. Kripke’s regress does not arise for all rules of inference and not even for the adoption of those rules that are of relevance for the discussion of the rational revisability of logic.

We will begin the paper in Section 1 with a brief summary of the use that Quine made of the regress argument against a conventionalist conception of logic and sketch Quine’s own view on the revisability of logic. Kripke seems to claim that the point that Quine makes against conventionalism should equally apply to Quine’s own view on the rational revisability of logic. In Section 2 we will look at which logical principles are at all subject to a potential regress or circularity problem and we will discuss whether the principles that are potentially subject to such problem are principles that are of relevance for the discussion of the rational revisability of logic.\footnote{Since our discussion can’t cover all possible revision to logic that one can come up with, we will limit our discussion to logics that are plausible alternatives to classical logic. We will motivate this choice in due course.} Our arguments in Section 2 will thereby follow the specific setup that Kripke introduced for the discussion of the Adoption Problem. In Section 3 we will investigate actual cases of proposed logical revisions in order to show how the more abstract considerations of the previous sections may apply to “real life” examples.

Since we arrive at a negative answer to the two questions above, we will close the paper in Section 4 by considering alternative targets for for the Adoption Problem. Perhaps it doesn’t primarily target Quine’s view on the revisability of logic but some
other aspect of Quine’s view on logic. However, as we will argue in that, also for these alternative targets Kripke’s argument doesn’t pose a real challenge.

The main claims of the paper are then that there is no adoption problem that would compromise rational revision of our logic, provided that we already possess some basic reasoning skills. This is the case both for the thought experiment considered by Padro and Kripke, and for more realistic scenarios of logical revision. Moreover, that basic reasoning skills are unadoptable is consistent with a Quinean philosophy of logic.

1 The Adoption Problem

According to Padro (2015), Kripke uses the following example to illustrate the problem of adoption:

Let’s try to think of someone—and let’s forget any questions about whether he can really understand the concept of “all” and so on—who somehow just doesn’t see that from a universal statement each instance follows. But he is quite willing to accept my authority on these issues—at least, to try out or adopt or use provisionally any hypotheses that I give him. So I say to him, ‘Consider the hypothesis that from each universal statement, each instance follows.’ Now, previously to being told this, he believed it when I said that all ravens are black because I told him that too. But he was unable to infer that this raven, which is locked in a dark room, and he can’t see it, is therefore black. And in fact, he doesn’t see that that follows, or he doesn’t see that that is actually true. So I say to him, ‘Oh, you don’t see that? Well, let me tell you, from every universal statement each instance follows.’ He will say, ‘Okay, yes. I believe you.’ Now I say to him, ‘“All ravens are black” is a universal statement, and “This raven is black” is an instance. Yes?’ ‘Yes,’ he agrees. So I say, ‘Since all universal statements imply their instances, this particular universal statement, that all ravens are black, implies this particular instance.’ He responds: ‘Well, Hmm, I’m not entirely sure. I don’t really think that I’ve got to accept that.’ (Padro, 2015, fn. 49)

1.1 Quine against conventionalism

Lewis Carroll’s similar dialogue between a tortoise and Achilles (Carroll, 1895) has famously been used by Quine (1976) in order to show that the logical positivists’ conventionalism about logic is in trouble. Conventionalism about logic (of the kind that Quine

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4Who the target of Quine’s paper ‘Truth by convention’ eventually is, is not clear. Quine doesn’t explicitly say that it is Carnap and there are reasons to think he targeted his own view (Ebbs (2011)) and that of C.I. Lewis (Morris (ming)).
considers) explains why logic should have a special status: Logical principles are knowable \textit{a priori} and necessarily true. According to conventionalism, we decide to maintain the statements of logic “independently of our observations of the world” and thus assign them a truth-value by convention. This accounts for their epistemic and modal status.

Although Quine expresses considerable sympathy for the view (granting that it is “perhaps neither empty nor uninteresting nor false”), he nevertheless sees it facing a difficulty that he summarizes as follows:

Each of these conventions [Quine refers here to the schematic axioms of propositional logic] is general, announcing the truth of every one of an infinity of statements conforming to a certain description; derivation of the truth of any specific statement from the general convention thus requires a logical inference, and this involves us in an infinite regress. (Quine, 1976, 103)

In Carroll’s dialogue, the tortoise challenges Achilles to get it to infer in accordance with Modus Ponens. Achilles fails to achieve this even though the tortoise is ready to accept an explicit statement of Modus Ponens as a true principle. For Quine, the upshot of that dialogue is that logic can’t be based on convention alone, since it seems that we need to have the ability to apply the supposed conventions and derive consequences from them in order to follow them. But then logic must be prior to such conventions (rather than the other way around).

In a word, the difficulty is that if logic is to proceed \textit{mediately} from conventions, logic is needed for inferring logic from the conventions. (Quine, 1976, 104)

Quine does see a way for the conventionalist to address this difficulty. What if we can adopt a convention “through behaviour” (Quine, 1976, 105) instead of adopting it via explicitly announcing it first? Perhaps the explicit formulation of these conventions can come later, once we have language and logic and all that at our disposal. For Quine this is a live option, but not one that he is still willing to describe as logic being based on “convention”. From Quine’s behaviorist point of view, behavior that follows a conventional rule is indistinguishable from behavior that displays firmly held beliefs.\footnote{In fact, Quine only makes the much weaker observation that it would be “difficult to distinguish” a behavioral adoption of conventions from behavior that displays firmly held beliefs.}

Since the label ‘convention’ is then without explanatory power, we can drop it from our account of logic.\footnote{See Azzouni (2014) for a discussion of conventionalism and Quinean arguments against it. Thanks}
1.2 Kripke against Quine

As Padro (2015) explains, Kripke now turns the regress7 argument against Quine himself. Quine had famously suggested in ‘Two dogmas of empiricism’ (Quine, 1953) that not even logic is immune to revision. Empirico-pragmatic considerations may lead us to the adoption of a new logic. A view that is, of course, quite compatible with the idea that logic is nothing but firmly held belief in the first place. Perhaps—so Quine’s own example—we may decide to adopt a logic that drops the principle of excluded middle because it may help to simplify quantum mechanics (Quine, 1953). However, Kripke seems to believe that Quine’s picture, viz. that we can treat principles of logic just like any other empirical hypothesis, is prone to the exact same objection that Quine mounted against conventionalism. Padro cites Kripke as follows:

... the Carnapian tradition about logic maintained that one can adopt any kind of laws for the logical connectives that one pleases. This is a principle of tolerance, only some kind of scientific utility should make you prefer one to the other, but one is completely free to choose. Of course, a choice of a different logic is a choice of a different language form.

Now, here we already have the notion of adopting a logic, which is what I directed my remarks against last time. As I said, I don’t think you can adopt a logic. Quine also criticizes this point of view and for the very same reason I did. He said, as against Carnap and this kind of view, that one can’t adopt a logic because if one tries and sets up the conventions for how one is going to operate, one needs already to use logic to deduce any consequences from the conventions, even to understand what these alleged conventions mean.

This is all very familiar as a criticism of Carnap. Somehow people haven’t realized how deep this kind of issue cuts. It seems to me, as I said last time, obviously to go just as strongly against Quine’s own statements that logical laws are just hypotheses within the system which we accept just like any other laws, because then, too, how is one going to deduce anything from them? I cannot for the life of me, see how he criticizes this earlier view and then presents an alternative which seems to me to be subject to exactly the same difficulty. (Padro, 2015, 113)

Stairs (2006) and Devitt and Roberts (npub) interpret the adoption problem as targeting to the work of David Lewis and others we now have a much clearer idea of how behavior that is based on firmly held belief can be distinguished from behavior that is guided by an implicitly adopted convention.  

7In Carroll’s original argument, the structure of the problem is a regress: the tortoise requires always new meta-principles in order to apply Modus Ponens at a given level. The regress is provoked, because the very rule that is supposed to be adopted is the rule that is necessary to apply that new rule. In that sense, the regress obtains because of that circularity. In what follows we will sometimes refer to that argument/problem as a regress or a circularity argument/problem.
in particular Quine’s idea that logic is revisable and that we can adopt a new logic. Padro (Padro, 2015) seems to see the adoption problem as a problem for adopting a logic in the first place and Kripke (2020) is vague about the target of the argument. Kripke discusses the adoption problem in a paper on Putnam’s views on the possibility of revising logic for empirical reasons and clearly seems to think that the adoption problem should have some relevance for the revisability of logic. His main target is the use that Putnam (1968)—and others who follow Quine’s views on the revisability of logic—makes of the phrase that we can “adopt” a new logic (Stairs, 2006, 2016). Thus, we take it that discussing the adoption problem as a problem in the context of logic’s revisability gets at what is ultimately at stake in Kripke’s original argument. (However, as we will discuss in the last section of this paper, Kripke also makes some remarks in his paper that suggest that he, too, may have the adoption of a first logic in mind.)

We will begin this paper by considering the Adoption Problem as a problem for Quine’s idea that logic is revisable but will discuss in the last section of this paper whether that is the best interpretation of Kripke’s attack on Quine. We hope that this brings some clarity into what the adoption problem possibly is and for which view this might be a problem. The idea that logic is rationally revisable is broader than the idea that a first logic can be adopted via the acceptance of logical principles, thus the former seems to be the natural starting point for our analysis.

According to this reconstruction of the argument, logic is not only not based on convention, but logic can’t be rationally revised either, because whatever empirico-pragmatic reasons we may have for preferring some alternative logic, we can’t adopt a new logic. Presumably the argument is then that the adoption of a new logical principle (as in Kripke’s example) would already presuppose the logical competence that allows us to apply such principle. However, as in Kripke’s example, if that competence is in fact the very rule we are supposed to adopt, then this can’t work.

A prima facie reasonable reaction to the argument so understood—due to Devitt and Roberts (npub), for instance—is to distinguish the way in which we come to know the propositional form of a logical principle, its representation, such as ‘from a universal statement, each instance follows’, and the way in which an agent can come to be governed by such logical principle, a state that may not necessarily require a representational form. The first kind of knowledge may be dubbed declarative, the second procedural. According to this first reaction, therefore, the sort of revision involved in Carroll’s example concerns

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8Finn (2019a) interprets Kripke to pose a problem for “anti-exceptionalism” about logic, but leaves it vague what aspect of anti-exceptionalism is the target. Revisability is, however, a central aspect of the anti-exceptionalist doctrine and clearly a potential target if there was a problem with adopting new rules.
the fact that declarative knowledge of a rule alone may not be sufficient to rationally revise one’s logical beliefs. But this does not rule out the possibility of training someone in acquiring procedural knowledge of a new logical principle.

A similar position is assumed by Graham Priest (2014), although framed in his distinction between the *logica docens*, *utens*, and *ens*. The logic we teach (*docens*) can be revised by means of a broadly abductive methodology. What is commonly called a ‘logic’, for Priest, should in fact better be seen as a ‘logical theory’, namely a substantial body of knowledge concerning some notion of logical consequence. Now, a logical theory can be rationally revised in the same way as other scientific theories can be revised, namely by comparing it with alternatives according to theory-choice criteria such as explanatory power, strength, adequacy to data, unifying power, and whatever else these may be. The logical theory we teach, therefore, can be rationally revised, and so can the logical theory we use. How? Simply by training oneself in a chosen *logica docens*. To connect Priest’s approach to rational revisability of logic with the Carroll-Kripke example, what seems to be clear is that for Priest the process of acquisition of a rule is not a local procedure, but rather a global process of acceptance of a logical theory that goes well beyond the rules of a formal system. This point will be further expanded in §3.

In the next three sections we leave aside these attempts to undermine the Adoption Problem that deny a significant role to the declarative knowledge of a rule or principle. We will work under the assumption that the declarative knowledge of a logical principle does indeed play a role in one’s actual adoption, and consider in more detail how such process could actually work. This is indeed how (Padro, 2015, p. 31) understands ‘adoption’: we adopt a way of inferring (for example, in accordance with Modus Ponens), if we pick it up on the basis of the acceptance of the corresponding logical principle alone.

As it will turn out in sections 2 and 3, there is no problem of adoption that would arise for the *revision* of logic (as Kripke seems to claim). It is true that one needs *some* basic reasoning skills in order to be able to adopt and apply new ones, but in pretty much all cases in which one has already a logic, these will be available.

### 1.3 Logica Utens

Although we will set aside Priest’s solution to the problem of adoption, it will still be useful for our discussion to help ourselves to a distinction between *logica docens* and *logica utens*. The former is an explicit theory that may or may not be formalized in precise mathematical terms.
A *logica utens*, on the other hand, is—in our terminology—the logic that we reason with under suitably idealized circumstances. What matters is that the *logica utens* is not just a description of all of our actual inferences (including all inferences we would ourselves accept to be mistakes) but rather a reconstruction of the rules we recognize as normatively governing correct reasoning. While Aristotle is widely credited with having started the business of developing a *logica docens*, *homo sapiens* much earlier started to develop a *logica utens*.

*Logica utens* will play an important role in our analysis of the Adoption Problem. We will argue that Kripke’s thought experiment is best understood as the attempt to revise one’s *logica utens*, and we will pinpoint precisely when this task is bound to fail, and when it is instead unproblematic. Even the more general context of revision of one’s logical theory can be thought of as an attempt to revise one’s *logica utens*: in those cases revision of *logica utens* amounts to a revision of one’s logical metatheory, and we will investigate whether this is a feasible task also in that context.

## 2 Patterns of adoption

### 2.1 What can we adopt?

As noticed already in Cohnitz and Estrada-González (2019), when one looks carefully at the Carroll-Kripke example, it becomes clear that not all principles are equally problematic. To see this, let us frame our discussion in a logical formalism in which one has finitely many rules for introduction and elimination for a finite set of logical connectives (natural deduction or sequent calculi are both adequate options). Consider the following version of our original dialogue in which universal instantiation is now replaced by the introduction of the existential quantifier. It involves subjects A and B and we assume, for the sake of the argument, that B is not able to perform inferences according to *Existential Introduction*. As before, we assume that B is willing to cooperate in accepting and reasoning according to the hypotheses that A provides.

A. Consider the hypothesis that, if some predicate $\varphi$ holds of $t$, then there is something that satisfies $\varphi$.

B. OK, I am considering it.

A. This piece of paper is white, isn’t it?

B. Yes.
A. Therefore, since if some predicate \( \varphi \) holds of an individual \( t \), then there is something that satisfies \( \varphi \), it follows that \textit{there is something that is white}.

B. Sure, thanks!

In the above dialogue, unlike what happens in the Kripke case, nothing prevents B from following and accepting A’s instructions. The reason is that no prior understanding of Existential Introduction is needed for B to follow the instructions given by A.

However, there is something else that needs to be presupposed by B. First of all they need the ability of inferring via Modus Ponens, as we learnt from Carroll’s example. To be clear, we employ the label ‘Modus Ponens’ for a rule of inference akin to the standard natural deduction rule, or the cut rule in a sequent calculus. A choice between one or the other may depend, for instance, on whether we conceive of the ‘if..., then...’ in A’s hypothesis as an entailment sign—in which case one needs cut –, or as an object linguistic conditional—in which case on needs a rule for the elimination of such a conditional. Of course we are not fixing a specific system in our discussion, and therefore these are at best structural analogies. We will come back to this point below.

In the light of Kripke’s example, it would prima facie seem that also Universal Instantiation is required. However, both in Kripke’s example and here we need much less than Universal Instantiation in full generality. Consider A’s last sentence: it presupposes the capability of recognizing the validity of the step that goes from an argument of the form \( \varphi(t/v) \therefore \exists v \varphi \), for all \( \varphi \), to an argument of the form \( P(t/v) \therefore \exists v P \) for a particular \( P \). Similarly, in Kripke’s example, the step that prevents the receiver of the instructions from agreeing on the desired conclusion is her incapability of recognizing the validity of the inference from an argument of the form \( \forall v \varphi \therefore \varphi(t/v) \) to one of the form \( \forall v P \therefore P(t/v) \). In both cases, it is a form of universal instantiation that is at stake. But at a closer look, the inferences under considerations are in fact of the form:

\[(\text{scs}) \text{ for any formula } \varphi, \text{ if } \Phi(\varphi), \text{ then } \Phi(P/\varphi), \text{ for some fixed argument pattern } \Phi.\]

\((\text{scs})\) is a very distinguished form of Universal Instantiation. First, quantifiers range over a fixed set of formulae of the language under consideration. Under the natural assumption that the languages we speak are countable, the size of such set is then countable too, whereas no such assumption is required for the general form of Universal Instantiation. Moreover, \((\text{scs})\) has a form that is well-known to logicians: it is a \textit{schematic substitution principle}—whence the label \((\text{scs})\)—, according to which, by accepting the schema, one accepts all its specific instances in the language under consideration.
This discussion can be generalized by formulating a more abstract RECIPE FOR ADOPTION in the box below.

RECIPE FOR ADOPTION:

1. One starts with a schematic logical principle of the form

   \( \text{(1)} \quad \text{if} \quad \Phi_1(\vec{X}; \vec{z}) \quad \text{and} \quad \ldots \quad \text{and} \quad \Phi_k(\vec{X}; \vec{z}), \quad \text{then} \quad \Psi(\vec{X}; \vec{z}), \)

   with \( \vec{X} \) and \( \vec{z} \) possibly empty strings of variables of finite length. Here the \( X_i \)'s are one sort of variables to be replaced with formulae, and the \( z_j \)'s are meta-variables for terms possibly including a different sort of variables for objects. Some machinery for renaming variables is also assumed.

2. One is then given a schematic instance of the antecedent of the conditional

   \( \Phi_1(\vec{A}; \vec{t}) \quad \text{and} \quad \ldots \quad \text{and} \quad \Phi_k(\vec{A}; \vec{t}) \)

   for \( \vec{A} \) formulae of the language and \( \vec{t} \) actual terms in the language.\(^9\)

3. (scs) enables one to go from (1) to

   \( \text{if} \quad \Phi_1(\vec{A}; \vec{t}) \quad \text{and} \quad \ldots \quad \text{and} \quad \Phi_k(\vec{A}; \vec{t}), \quad \text{then} \quad \Psi(\vec{A}; \vec{t}), \)

4. by Modus Ponens applied to (2) and (3), one concludes \( \Psi(\vec{A}; \vec{t}) \), thereby inferring according to (1).

A few comments to the RECIPE FOR ADOPTION are in order. First, we are analysing Kripke’s pattern for adoption. As such, the intended application of our pattern is the scenario envisaged by Kripke: we are not putting forward a recipe to adopt any possible logical principle, but a list of notable examples. That being said, the recipe possesses some degree of flexibility intended to deliver fruitful applications under several specific formalisms. As anticipated, a first (deliberate) scope of manoeuvre is given by the way in which premisses of inferences are gathered in (1). The most straightforward way to understand ‘and’ is as a metatheoretic juxtaposition sign, very much like commas in a
sequent calculus formulation. In this way, the final detaching step that we call ‘Modus Ponens’ becomes akin to an application of the structural rule of cut. One then easily sees that, under this reading, the principle of conjunction introduction ‘if \( \varphi \) and \( \psi \), infer \( \varphi \land \psi \)’ is unaffected by the adoption problem.

As noticed by Kripke (2020) himself, if instead one identifies ‘and’ with the object linguistic conjunction, conjunction introduction might acquire a status analogous to Modus Ponens and Schematic Substitution, because gathering premisses via conjunction presupposes the rule of conjunction introduction. An alternative may be to dispense with conjunction, and consider the operation of gathering premisses via nested conditionals (e.g. ‘if \( \varphi \) and \( \psi \), then \( \chi \)’ is turned into ‘if \( \varphi \), then \( \psi \) only if \( \chi \)’). Under this assumption, other principles will become unadoptable, such as the principle of conditional introduction.\(^{11}\)

The extent to which (scs) is a logical rule can be debated at length: it can even be argued that it is the logical rule, as it is possible to axiomatize, say, classical logic, by resorting to axioms involving specific predicate letters—and not axiom schemata or rule schemata—and some principle akin to (scs). For our concerns, however, what matters is that the form of universal instantiation that Kripke suggests is presupposed by our capability of acquiring Universal Instantiation is not as strong. Rather, it is a very specific form of universal instantiation that has much to do with our ability of recognizing and combining syntactic patterns.

The problems encountered with the adoption of a logical rule—as far as Kripke’s example is concerned—boil down, therefore, to the necessity of certain presuppositions to the process. Under a plausible reading of the pattern isolated by Kripke, such presuppositions amount to competence with Modus Ponens and the validity and a very specific form of universal instantiation (scs).\(^{12}\)

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\(^{10}\) Again, some vagueness concerning different implementations of this idea is assumed: we do not take a stance on whether commas should be understood as distinguishing elements in a set, a multiset, or a sequence.

\(^{11}\) We would like to thank an anonymous referee for asking for a clarification of the status of what we call ‘Modus Ponens’. Given our purpose, any choice that is more specific than our current proposal would lead to specific choices that are not compatible with the general analysis of Kripke’s project that is the main aim of the paper.

\(^{12}\) A recent paper by Suki Finn (2019b) makes use of the same idea, but erroneously assumes that the ingredients of this “recipe” are Modus Ponens and Universal Instantiation and that both of these rules are individually necessary and jointly sufficient for the adoption of any other logical rule. As we argue here, the recipe doesn’t require Universal Instantiation in full generality but only a very restricted form. Also, depending on the logical rule in question, Modus Ponens is not always necessary either (just consider rules that allow adding theorems to any step in the reasoning). As explained, those two rules may also be not jointly sufficient.
2.2 Where can we adopt?

In Kripke’s example, the receiver of the instructions may not be able to perform any inference. The scenario is compatible with a tabula rasa adoption. Let us now consider a more realistic, although still highly idealized, scenario in which an agent is in possession of some inferential abilities that are in need of revision. In general, revisions can reasonably involve either (i) dropping some principle from the set of one’s logical beliefs, or (ii) adding principles to it.\(^\text{13}\) We call the former process DROP, and the latter ADD.

Most cases of proposed logical revision at the heart of modern and contemporary debates involve DROP. Starting with classical reasoning, intuitionists proposed to drop the law of excluded middle or, equivalently, to weaken one of the rules for negation. Para-complete and paraconsistent logicians also propose to drop one of the rules for negation, although their weakening of classical negation is more severe than the one proposed by the intuitionists. Some subtler proposals are also possible. Supervaluationists, for instance, agree with all inferences of classical logic of the form \(\langle \Gamma, \varphi \rangle\), but disagree on inferences with multiple conclusions.\(^\text{14}\)

Let us start with DROP. There are various scenarios compatible with dropping a logical principle. In the trivial case, revision simply amounts to disregarding some principle, previously regarded as logical. There is no adoption involved in the revision, and a fortiori there is no adoption problem. In a slightly less trivial case, the rule that needs to be adopted is not one of the principles that fall under our understanding of Modus Ponens.\(^\text{15}\) In this case, it is clear that the patter of adoption straightforwardly applies (modulo some adjustments required by the specific formalism employed and discussed above). For instance, an agent who is able to infer according to Modus Ponens and scs is in the position to adopt the familiar principles involving conjunction, disjunction, negation.

Another case of DROP may concern the adoption of a new rule by restricting the scope of previously acquired rules. The crucial (and non-trivial) case involves adoption of restricted versions of Modus Ponens. Some paraconsistent logics, Priest’s LP for instance, result from classical logic via the restriction of the elimination rule for the conditional to formulae that are not truth value gluts (Priest, 2008). Similarly, non-transitive logics restrict the meta-inference of Cut (Ripley, 2015), by allowing it only for

\(^{13}\)Of course it is possible that the proposed adoption in question leads from a set of logical beliefs to another which is inconsistent with the previous one, but in the reasonable cases in which this happens one can always describe this process as the result of first dropping some rule and then adding to the remaining principles some other principles.

\(^{14}\)For instance, they drop the classical inference \(\langle \{\varphi \lor \neg \varphi\}, \{\varphi, \neg \varphi\} \rangle\).

\(^{15}\)We are leaving out scs from the picture, because of its special status.
some non-pathological sentences. In such a scenario, a crucial issue concerns whether the pattern of adoption should be itself revised to feature such restricted detachment principles instead on the original form of Modus Ponens. Luckily, the answer is positive. If one wanted to apply the pattern for adoption to the restricted Modus Ponens, schematic substitution and the restricted form of Modus Ponens would suffice. The (re)adoption of other principles by means of restricted Modus Ponens may be more problematic. For instance, paraconsistent logics such as LP feature unrestricted principles governing conjunction and disjunction, and therefore the adoption of such principles will not involve only sentences with a classical truth value.

Problems can occur only, if the logical resources become to weak to apply the principle even with suitably restricted rules. Whether there are interesting cases of that kind, will be explored below.

Let us now turn to add. Prima facie there are good reasons to doubt the significance of add, if one assumes that the process of adoption has classical logic as its starting point and restricts oneself to the propositional case. The Post completeness of classical propositional logic tells us that the only consequence relation that properly extends it is the trivial one. On the other hand, when we move to first-order classical logic, which isn’t Post-complete, it is also clear that Modus Ponens and Universal Instantiation are already in place. Therefore, any revision that follows our schema for adoption is also unproblematic—new rules can be adopted and applied by following the pattern for adoption isolated above. For instance, we might consider a higher-order version of the rule of existential introduction:

\[ \text{from } \varphi(R), \text{ infer } \exists X \varphi(X) \]

with \( R \) a set variable which is free for \( X \) in \( \varphi \). As before, the adoption of such a rule would require the capability of applying (scs). In the specific case of (2), the schematic variable needs to be of a suitable type; it should be capable of taking variables like \( X \) as arguments. This process, however, is still carried out once a suitable language is fixed. The substitution involved in the adoption of (2) does not require any substantial decision on the semantic status of the different types of variables. Similarly, a higher-order version of the rule of (monadic) Universal Instantiation

\[ \text{from } \forall X \varphi(X), \text{ infer } \varphi(P/X) \]

can be accommodated in our framework via (scs) once a suitable language is fixed.
What is only required is that the schematic variable $\varphi$ can be instantiated to a specific formula of the higher-order language one is considering. In other words, in the pattern of adoption for (2) and (3), one always assumes a specific domain of syntactic entities on which (scs) operates. And this is all that seems to be required.

As expected, the only problematic candidates in the context of ADD are logics that either don’t have what we called Modus Ponens or do not have (scs). It is fair to say that, if one operates in Kripke’s idealized scenario of a tabula rasa adoption, our analysis deems the unrestricted rule of Modus Ponens as unadoptable. However, it is equally fair to say that the debate is still open on whether logics that do not feature Modus Ponens satisfy some fundamental adequacy requirements for playing the role of a logica utens, i.e. whether a logic without such a rule could be an adequate formal model for any possible form of natural reasoning. We rest content with the claim that, for the overwhelming majority of case studies, the last step (3 to 4) of our pattern of adoption applies.

What about (scs)? It is a common assumption in much of contemporary semantics that natural languages must (in some way, (Cohnitz, 2005)) be compositional. How else could it be explained that we can use and understand new sentences with novel meanings? However, compositionality requires some form of systematic syntactic decomposition and of keeping track of how, for example, argument places of predicates are filled. It is hard to see why such capacity shouldn’t already be sufficient for the kind of schematic substitution that Kripke’s example requires. Compositionality by itself guarantees that competence with a sentence like ‘Sam kisses Martin’ entails competence with ‘Martin kisses Sam’, ‘Reinold kisses Julie’—this fact is behind the systematicity argument for compositionality (Szabó, 2000). But then the basic skills involved in processing a compositional language (treating linguistic items as schematic and (re)combinable with other linguistic items of certain syntactic categories) already allow one to reason in accordance with (scs). This skill doesn’t seem to be in need of “adoption”.

(scs) is weaker than the rule of Universal Instantiation. It is a basic (logical or linguistic) skill that is presupposed by reasoning of any kind. Not just any logical rule we learn, but learning any new compositional phrase requires mastery of schematic substitution. A fortiori, any logic that is supposed to model an actual logica utens will have to contain (scs) then.

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16To be precise, for the application of (scs) in reasoning, we need not only the ability to compose new expressions, but also to decompose them. This requires compositionality, as well as inverse compositionality (Pagin, 2003).

17And, as we argued above, schematic substitution is implicit in our mastery of composing and decomposing complex expressions in general.
Again, there can be formal systems that are weaker than classical logic and that
do not contain Modus Ponens or (scs). But the real question is whether there is any
formal system that models a \textit{logica utens} but fails to enable the reasoner to adopt a new
rule. If any application of logical rules requires some (suitably restricted form of) Modus
Ponens and (scs), and if from that a reasoner can obtain a (suitably generalized) form
of Modus Ponens and (scs) that is sufficient for grasping the application conditions for a
new rule, then every logic that is a possible \textit{logica utens} will allow upwards adoption (as
well as downwards adoption to any logic that is a possible \textit{logica utens}). If this is right,
then Kripke’s “adoption problem” does not actually pose a problem for the adoption of
a new logic.

But Kripke’s scenario is anyway highly artificial. No one adopts a logic simply
because some oracle told them that the principle behind it is logically valid. We may
come to reason in new ways, because we adopted a new theoretical perspective on matters
of validity.

3 Adaption in a logical theory

We have argued that the revision of logic by adoption of a new logical principle is best
understood as a revision of one’s \textit{logica utens}. In this section we consider the patterns
of adoption isolated earlier in the arguably more realistic context of a \textit{logical theory},
typically defined as a collection of principles governing the core notions involved in
one’s specific account of logical consequence: truth-preservation, predication, negation,
implication, assertion, formality, consistency, provability and so on. Therefore, giving a
full account of one’s preferred logical theory is often a highly non-trivial matter. That
the Adoption Problem discussed by Kripke should carry over to these more realistic
contexts is clear from the discussion of empirically motivated logical revision found in
Kripke (2020).

3.1 Deflationary views of logical theories

The preliminary characterization of logical theories just given is not the only one con-
sidered in the literature. It more or less aligns to what Hjortland (2017) calls \textit{non-
deflationary} logical theories. Following this terminology, a typically deflationary ac-
count is the one articulated in Williamson (2017), which holds that the ultimate task
of logical theories is to unravel general claims about the world. Meta-linguistic notions
such as truth and validity are not the primary concern of logic, which is essentially a
non-metalinguistic enterprise pointed at discovering absolutely general laws of reality. In this, logic does not differ from physics, or from metaphysics; it only proceeds at a much higher level of abstraction.

Williamson suggests that a logical theory is a collection of nonmetalinguistic generalizations corresponding to logical truths. This picture is motivated by the following process: Williamson starts from valid inferences in some logic $\mathcal{S}$ in a language $\mathcal{L}_S$—e.g. $\neg\neg\varphi \vdash \varphi$. It proceeds by extending $\mathcal{L}_S$ with new, higher-order variables of the same type as formulae of $\mathcal{L}_S$ and by replacing the entailment relation with a conditional—in our example, this turns $\neg\neg\varphi \vdash \varphi$ into $\neg\neg X \rightarrow X$. The process is then completed by universally quantifying over the free higher-order variables of the translation of the logical claim under considerations. A logic, in this view, is a collection of claims such as $\forall X (\neg\neg X \rightarrow X)$. Endorsing a logic is endorsing a collection of universally quantified claims: since there is no reason to consider higher-order quantification as more metalinguistic than first-order quantification (Williamson, 2017, p. 329), a logical theory is no more metalinguistic than any other theoretical enterprise seeking universal laws, such as physics itself.

Given our analysis, the problem of adoption in a deflationary logical theory of the kind just sketched does not arise. Already the process of turning a purported valid inference into a universal generalization of the appropriate type requires a prior understanding of quantification. It is hard to see how this understanding may not involve something as basic as (scs): this is especially clear in the step that requires the expansion of one’s language with variables of the appropriate type. The very adequacy of this process seems to rest on the capability of instantiating such variables with formulae of $\mathcal{L}_S$, as required by (scs). Moreover, the substitution of the entailment sign with a suitable conditional certainly presupposes a conditional that satisfies Modus Ponens. How can the reduction be put to use, if one cannot retrieve the original inference by assuming an instance of the antecedent of the law-like conditional and conclude its consequent via Modus Ponens? The structural assumptions required by Williamson’s view of logical theories therefore presuppose both (scs) and Modus Ponens; our analysis of the pattern for adoption entails that the circularity involved for the adoption of a new rule does not arise in the presence of such principles.\footnote{Williamson ultimately rejects this Tarski-Bolzano procedure of bringing inferences to their normal form as a tool to compare logical consequences. This is because the procedure requires a strong conditional, and many of the logics involved in the comparison will not have it. What we said however still stands: on this view of logical theories Modus Ponens and (scs) are essential requirements.}

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3.2 Logical theories

Logical theories, in the abstract—and more substantial—sense considered in this section, can be seen as the formal counterpart of logicae utenses. In the same way as a logicae utenses encodes the agent’s dispositions towards a class of inferences (or meta-inferences), a logical theory enriches this acceptance of a class of validities with a collection of metatheoretic claims concerning semantic and proof-theoretic notions associated with such inferences. For instance, the logical theory of intuitionistic logic includes an account of what is a canonical or direct method of verification, as opposed to an indirect one. Similarly, the logical theory of paraconsistent logic involves a characterization of negation and falsity (and truth) that substantially differs from the classical exclusive approach to negation. Taken at face value, claims of the sort just described belong to the metatheory of one’s logic. And such a metatheory typically amounts to a fragment of classical or intuitionistic mathematics.

There are at least two possibilities to formulate the adoption problem in this richer framework, depending on what one considers to be the class of logical principles that can be adopted/revised. On the one hand, one might consider revision and adoption of the purely logical part of one’s metatheory, which may not align with the object-theoretic logical principles. On the other, one can extend the status of logical principle to core metatheoretic principles such as consequence and truth, and consider their adoption and revision.\textsuperscript{19}

Let us consider the scenario in which one wishes to revise/adopt logical principles of one’s overall logical theory, including the logic of metalinguistic concepts. In the abstract case, it is clear that this is no more nor less problematic than allowing for a revision of object-linguistic logical principles: the logical component of one’s logical theory is simply a collection of inference patterns that one recognizes as valid in one’s metatheory. There seem to be no substantial differences between the analysis of the

\textsuperscript{19}One might also think about a third option, in which one’s logical theory plays a purely instrumental role. In this scenario, one would keep all metatheoretic principles fixed, consider them in a purely instrumental role, and take into account only adoption and revision for the object-theoretic logical inferences. The discussion of the previous section would then largely transfer to this case, with possibly a further complication. Suppose we are in the crucial case of the absence of Modus Ponens on one’s object-theoretic logical toolbox. In this case the instrumentalist about metatheory may find herself in the position of not accepting (yet) object-theoretic claims of the form \( \varphi, \varphi \rightarrow \psi \; \vdash \psi \), but accepting—given a standard set-theoretic semantics:

\[(4) \quad \text{If } '\varphi', \text{ is true and (if '}'\varphi', \text{ is true, then '}'\psi', \text{ is true), then '}'\psi', \text{ is true.} \]

where ‘is true’ is a standard Tarskian truth predicate for the object language. Therefore, the instrumentalist would have to argue that, even though she is able to infer on the basis of principles such as (4), she is in no position to adopt Modus Ponens at the object linguistic level.
local adoption problem above and the present case: again, the only problematic cases might be cases of ADD, in which from a weaker metatheory one moves to a stronger metatheory. For instance, one might ask whether the intuitionistic logician is able to adopt a classical perspective on validity. In the current setting, this can simply be reduced to the problem of whether one can instruct an intuitionist to infer according to, say, double negation elimination \( \neg\neg \varphi : \neg \varphi \). But in the presence of (scs) and Modus Ponens, we have seen that this is unproblematic: one starts with exhibiting a specific doubly negated instance \( \neg\neg A \) of \( \neg\neg \varphi \); by (scs), one provides the intuitionist with the concrete instance of—a suitable translation of—the original principle ‘if \( \neg\neg A \), then \( A' \). From \( \neg\neg A \) and ‘if \( \neg\neg A \), then \( A' \), the agent that possesses the general capability of inferring by Modus Ponens can immediately conclude \( A \). Under the assumption that intuitionistic or classical foundations are the only reasonable candidates for the logic of the metalinguistic components of one’s logical theory, we can safely conclude that no worries of circularity can arise in this second reading of logical theories.

The assumption that one’s logical metatheory is framed in classical or intuitionistic set theory may be questioned. There have been interesting attempts, in the context of some approaches to the semantic paradoxes, to align a weaker nonclassical approach—generally substantially weaker than intuitionistic logic, since semantic paradoxes affect classical and intuitionistic logic alike—in the object theory with a nonclassical metatheory (Leitgeb, 2007; Bacon, 2013; Weber et al., 2016). Such attempts, however, are at best at an initial stage and cannot yet be considered to be actual rivals of a classical or intuitionistic metatheory. For instance, most of these meta-theoretic results heavily rely on a classical meta-meta-theory. What would be required is a non-classical set theory (or of an alternative foundational framework) in which all metatheoretic reasoning could be performed.

The status of non-classical set theories, however, is controversial. Let us consider for instance on some paracomplete and paraconsistent options. Partial set theories have been developed by Gilmore, Aczel, Feferman (Gilmore, 1974; Aczel and Feferman, 1980; Feferman, 1984): the naive comprehension principle is built on top of a three valued logic such as Strong Kleene logic. Consistency is obtained by showing (in a classical metatheory) that membership can be interpreted by means of a positive inductive definition. The main drawback of such attempts consists in their deductive weakness: the theories are able to recover only a fragment of predicative mathematics.

Paraconsistent set theories have also been extensively studied in recent years. Several combinations of set-theoretic and logical principles are possible. One option is to formulate naive comprehension on top of the LP (Restall, 1992; Priest, 2006). Due to
the weakness of the conditional of LP, it is not clear whether this option can deliver standard set-theoretic results such as Cantor’s theorem, or even the existence of two objects (Weir, 2004). An alternative is to replace the conditional of the paraconsistent logic with a relevant conditional. In this way, a substantial amount of standard results of classical set theory can be obtained (Weber, 2012). However, doubts still remain about the adequacy of such an option: as argued in Incurvati (2020), the relevant conditional is insufficiently motivated, and the fundamental extensional nature of the set concept is compromised of in such approaches—there are sets that have the same members but that are not identical (Incurvati, 2020, Ch. 4).

We are left with the possibility of adopting/revising quasi-logical principles such as truth and falsity. This is, arguably, the option that is closest to actual cases of revision of one’s logical assumptions. Paraconsistent and paracomplete logicians motivated by semantic or logical paradoxes, for instance, aim at a revision also of foundational tools, such as comprehension axioms, that are needed to define their notion of logical consequence. In this context, one considers not only a collection of logical inferences, but also the principles of quasi-logical notions such as truth, property predication, and consequence as possible candidates for revision. Can the worries of circularity/regress adumbrated in the local case of adoption in the previous sections have some bearing on such cases of revision?

If the adoption/revision process is a local process involving some specific quasi-logical rules and follows the blueprint of Kripke’s setup, our analysis in §2 can be transferred with only little modifications. For instance, if one’s logical theory makes essential use of the notion of truth, one might want to adopt/revise suitable principles for the truth predicate, e.g. a disquotational rule of the form ‘from \( \varphi \), infer \( \text{True}("\varphi") \).’ If Modus Ponens and (scs) are available, one can essentially follow the pattern outlined above for the case of adopting a logical rule such as double negation in an intuitionistic logical theory. The only step that requires care is the selection of a suitable range instance of instances of (scs). In the case unrestricted schemata such as double negation, in fact, specifying a range of instances of (scs) is a trivial affair: all sentences of the language are allowed. By contrast, due to the Liar Paradox, selecting a suitable range for the instances of \( \varphi \) in the truth rules might prove to be involve resources that are very complex in computational terms. We cannot choose all instances whatsoever to avoid inconsistency, and a more sophisticated procedure is needed. Now, if this procedure is

\[ A \text{ couple of qualifications about the example: first, the rule should be intended to apply also to } \varphi \text{ that we have assumed, and not only proved. Secondly, this rule should be intended to be adopted together with other truth rules. These qualifications are needed to ensure that the rule characterizes truth, and not weaker notions such as provability.} \]
purely syntactic, it can be easily implemented in the pattern for adoption stated above without any ad hoc move. For instance, if one intends to adopt the rule ‘from \( \varphi \), infer True(⌜\( \varphi \)⌝)’ for instances of \( \varphi \) that do not contain ‘True’, the relevant specification of the range of (scs) is a fairly simple procedure—at most primitive recursive—and can be reasonably taken to be part of the conceptual toolbox of anyone that understands the syntax of the language of their logical theory.

If the specification of the relevant instances of (scs) is not syntactic, it may result in a more complex procedure. If, for instance, this involves selecting the grounded sentences in the sense of Kripke (1975), or the set of stable truths in the sense of the revision theory of truth Gupta and Belnap (1993), this would involve a highly non-computable process (McGee, 1991; Burgess, 1986). Therefore, we might have a situation in which there is no Kripke-style circularity in adopting ‘from \( \varphi \), infer True(⌜\( \varphi \)⌝)’, but simply the absence of a suitable schematic substitution rule to implement in the pattern of adoption for such rule. It should be clear, however, that this scenario is perfectly compatible with our analysis of the problem of adoption/revision. Whereas the adoption problem concerns the one’s (seeming) impossibility of inferring according to a rule that is available to her, in the scenario under consideration the agent does not have at her disposal a suitable version of (scs) to perform inferences, because its range may be too complex to be specified.

We are then left with the familiar scenario in which one would like to adopt/revise quasi-logical rules but does not possess Modus Ponens. We have already cast some doubts on the availability of a workable logica utens in the absence of Modus Ponens. In the context considered here, this is even more so, since a logical theory may involve complex semantic constructions couched in classical mathematics, which require a substantial use of classical logic.

4 Alternative Quinean Targets for Kripke’s Argument

For all we have argued so far it seems that there is no adoption problem that would pose an obstacle or challenge to the idea that we can rationally revise our logica utens, provided that prior to the revision we already possess some basic reasoning skills and that our revision is supposed to preserve these. Neither in the abstract scenario that Kripke presents nor in more realistic cases is it plausible to assume that we lack the resources to apply new logical rules in reasoning.

As we explained in Section 1, we started with discussing the case of revision, since that seemed to us the broadest target for the adoption problem. In this last section,
we will look at other aspects of a broadly Quinean philosophy of logic that could be potential targets of an adoption problem.

We could identify four possible alternative targets that are part of Quine’s conception of logic and may, at least *prima facie*, be affected by the proposed regress. The candidates are in turn the adoption of a first logic, the transition from the acceptance of a principle to the adoption of certain behavior, the problem of the missing normative force of purely descriptive logical principles, and the *knowledge that/knowledge how*-distinction. We will discuss the candidates in this order.

### 4.1 The Adoption of the First Logic

So far we have considered the adoption problem as as a challenge for Quine’s idea that we can adopt a *new* logic. So it was legitimate in our argument to suppose that some logic and some language is already in place and that an individual has on the basis of some *reasoning* arrived at the conviction that she should adopt a different way of reasoning, that she should adopt a new logic.

But perhaps is best understood in close similarity to Quine’s original point against conventionalism and concerns the question how—on Quine’s view—logic could have ever gotten off the ground (Quine, 1976). After all, also on the conception that logic is just general, firmly held belief (Quine, 1953), there seems to be the issue that firmly believing Modus Ponens does not yet allow you to reason with it, if you don’t yet have that capacity. Thus, as a general theory of what logic is, Quine’s theory isn’t better than conventionalism, since it still is open to the challenge that it can’t explain how the first logical principles could have been adopted in absence of an already existing logic.21

Although this well may be so, it is not clear that this is a challenge that Quine needs to address. Or, in other words, it seems to us that Quine, quite clearly, does not have to address it. Quine (1976) presents a picture according to which the first principles of logic are not adopted as a result of engaging with some explicit formulation of the principles (as conventionalism has it), but where they get adopted in behavior and only later are reconstructed in terms of explicit reasoning principles or rules. This adoption in behavior does not require that Quine’s theory of belief revision applies to it, so he does not at all need to explain how *homo sapiens* managed to develop structured reasoning that is describable in terms of schematic inference principles. This should be part of a general naturalistic account of how higher cognition and reasoning in general developed. To require that Quine’s conception of logic provides some detailed explanation of this

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21 This seems to be how (Padro, 2015) understands the adoption problem.
process is entirely inadequate.

It is worth emphasizing that conceding that the acceptance of logical principles cannot explain how reasoning got off the ground takes nothing away from the idea that principles of logic are (as far as epistemology is concerned) just like other hypotheses. By developing a *logica docens* (as a formal representation of our most general ways of reasoning) we can critically study the way we think about most general matters (or matters most generally) and maybe decide to make revisions to those central aspects of our web of beliefs. Just as we would do with other hypotheses. How we could then “adopt” the so revised logic, we have described above.

### 4.2 From Belief to Behavior

A second potential target for the regress argument is Quine’s idea of the status of logic in the web of belief. Quine (1953) considers logic to be nothing but firmly held belief, statements that are just like any other statements in the web of belief, with the only difference, that they are more central than others, and thus less likely to be given up. But adopting a logic is not just adopting some belief. It is adopting a way of reasoning. There are two ways to make that challenge. The first would be to see this as a critique of Quine’s behaviorism. For a behaviorist, having a certain belief (for example, the belief that Modus Ponens is valid) just means to show certain forms of behavior (for example to reason in ways that are licensed by Modus Ponens). But perhaps that’s too short-sighted. As Kripke’s thought experiment shows (on this interpretation), one may accept a belief (viz. that Modus Ponens is valid) and yet fail to show the appropriate behavior (e.g. to assent to implications that are licensed by Modus Ponens). The thought experiment then doesn’t show that there indeed is a regress or circularity problem, but that there may be a problem of a certain kind of “stubbornness”: someone may count as having grasped and accepted a certain belief, but just doesn’t act in a way that may be canonical for the ascription of that belief.

This may be a reasonable challenge to the idea that ‘*S* believes that *p*’ can be analysed as ‘*S* is disposed to assent to this and that under conditions such and such’. But this

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22To see this, maybe it helps to consider an analogy with scepticism about our senses. Maybe we need to have already default trust in our senses in order to be able to learn anything from them. That may make the hypothesis ‘I can trust my senses’ special (in comparison to other empirical hypothesis) insofar as my knowledge of the world wouldn’t get off the ground without it. But even if that were so, *this* would not make this hypothesis immune to revision, not even immune to revision via information that I receive through my senses.

23Quine’s behaviorism is a well-known aspect of much of his work. We already encountered it in Quine (1976), when Quine argues that there is no difference between firm belief and implicit and spontaneous acceptance of a convention. Most famously, Quine’s behaviorism shows in his arguments in Quine (1960).
doesn’t seem to be a specific problem for Quine’s theory of logic than rather a problem for Quine’s theory of belief. However, while the regress/circularity argument displays the problem, it doesn’t actually establish anything that could seriously be regarded as an argument for the claim that such an analysis must fail. It seems still perfectly reasonable to just respond to such argument that it merely shows that the person in the dialog who doesn’t reason in accordance with, for example, Modus Ponens has not yet actually adopted the relevant belief.

4.3 The Normative Force of Logical Principles

A closely related challenge (one that actually makes use of a regress) is to interpret Kripke’s argument as revealing that Quine overlooked the normative nature of logic (if one believes that is has such a normative nature). Logic, on this view, tells us how we ought to reason. However, the general principles that are featured in Kripke’s thought experiment are not norms or imperatives. They don’t say anything about how anyone should reason. Therefore there is a gap between adopting the belief that a certain logical principle is true and adopting the norm that one ought to reason in a certain way. Quine, who takes logical principles to be just like any other general scientific hypotheses, overlooks this.

There are two reasons why this is not a plausible target for Kripke’s argument. First, not all logical norms or imperatives will hold unconditionally. But if they are norms that apply under certain conditions, then also a conditional norm could do nothing about the regress. The reasoner in the Kripke scenario would still need to be already following that norm in order to apply it under the current conditions. Thus, just adding deontic force to a rule doesn’t help with the regress at all, our hypothetical reasoner would still have to instantiate the general norm to the current case and then detach a consequence concerning what they now ought to infer.

The second reason is that the plausible normative force of logical principles is in fact too weak to be of any help in Kripke’s thought experiment. As Besson (2016) explains, the recent discussion of the normative force of logic strongly suggests that in order for the argument to go through, we’d need an imperative or a rule that would “move” a subject to reason in accordance with the logical principle at issue. However, as we have learned from Harman (1986) and others, logical principles can’t give rise to such rules. It simply isn’t always rational to use Modus Ponens and endorse \( q \) whenever you believe \( p \) and \( p \supset q \) for some \( p \) and \( q \). However, a weaker principle that would, say, allow that it is rationally permissible to believe \( q \) whenever you believe \( p \) and \( p \supset q \) for some \( p \) and
$q$ is plausible, but would not lead to a plausible regress (see Besson (2016) for details).

Once you know the principle

\[(5)\]

Given your beliefs $P$ and (if $P$, then $Q$), you are rationally permitted to reason to $Q$.

We can explain why you should be rationally permitted to reason with Modus Ponens. If the regress argument is supposed to make a point about normativity, it simply operates with the wrong deontic force.

4.4 Knowledge that and knowledge how

This leaves us with a last candidate which again tries to explain the problem of the regress by a certain insufficiency of the merely propositional knowledge that we acquire, when we accept the claim that Modus Ponens is valid. We mentioned in the beginning in Section 1 that Priest as well as Devitt and Roberts both see the problem of adoption as primarily an issue of acquiring certain knowledge how after one has convinced oneself of the relevant knowledge that. Stairs (2006) also seems to understand Kripke in this way.

Take a familiar analogy: from reading a book about how one rides a bike, one doesn’t know yet how to ride a bike in the sense that one won’t be able (yet) to ride a bike. The latter will require certain practical competence, a skill, that can not be acquired by simply reading a description of what that skill involves. Instead, the acquisition of that skill might require training.\(^{24}\) In the regress argument, the subject accepts Modus Ponens but doesn’t have the skill to apply it, she thus gets a new bit of propositional knowledge which she doesn’t know how to apply either, and so forth.

Priest as well as Devitt and Roberts seem to think that also the adoption of a new logic requires that we train ourselves in the application of a rule in order to be able to apply it. However, as our discussion above shows, the competence that rule application of logical principles requires is merely the competence with basic inferences like Modus Ponens or scs for the examples that we considered. For instance, one could infer according to Conjunction Elimination by just plugging the conditional rule into our recipe above. The relevant knowledge how, in these cases, is a certain basic capacity to reason in the first place. Adoption of a new rule thus does not require training in new

\(^{24}\)We don’t distinguish here between knowledge how and a skill, for the purpose of our argument it is sufficient to note that there are skills for which it is true that they can’t be acquired by just understanding an instruction.
rules.

Another question may be what it takes to “see” new implications that one didn’t see as implications with the “old” logic, or how one can get to stop seeing implications that aren’t implications according to a new logic. This seems to be what Kripke has in mind when he is complaining that a merely formal account of logic would not be the same as an intuitive form of reasoning:

What I mean is this: you can’t undermine intuitive reasoning in the case of logic and try to get everything on a much more rigorous basis. One has just to think not in terms of some formal set of postulates but intuitively. That is, one has to reason. [...] One can only reason as we always did, independently of any special set of rules called “logic”, in setting up a formal system or in doing anything else. (Stairs, 2016)²⁵

This version of the adoption problem seems to be what Kripke originally had in mind, but it neither leads to a regress, nor is it very convincing. The regress is irrelevant, since the problem is not that a logical rule is missing and requires the introduction by some explicit statement of the rule (the application of which again requires the rule, and so on ad infinitum). The problem is rather that any formal statement of logical laws is not the same as a way of reasoning. Thus, whether such a formal account is stronger or weaker than our actual way of reasoning, or in our terminology, whether revision goes via DROP or ADD, is irrelevant; if a formal logic does not agree with our intuitive way of reasoning, we will not be able to adopt such logic. Seeing that a consequence follows is as impossible to adopt as unseeing that a consequence follows, according to that view.

The point is then not that we need training to be able to apply a new rule (i.e. to be able to apply a new general rule to a new concrete case). As we argued above, application of the rules is easy once you have the skill necessary to follow our recipe. The problem is rather that such a form of application of an explicit rule does not count as reasoning.

But why should it not? Why should the habituation of a logic have any special status? Kripke presumably does not want to give the same value to all our dispositions to draw inferences intuitively. We often make mistakes in our intuitive reasoning. Maybe reasoning is a complex cluster of dispositions for Kripke; dispositions to draw inferences as well as dispositions to retract them after reflection. But if reasoning is such a wider cluster, then reasoning is malleable. And once reasoning is malleable in light of new

²⁵Stairs (2016) and Stairs (2006) are also discussions of Kripke’s lectures, but focus primarily on his case against quantum mechanics and less on a reconstruction of the adoption problem.
information about our inferences not being valid (maybe on the basis of a formal representation of that inference), why stop there? Why only consider reasoning as a set of dispositions stable under such a narrow equilibrium, rather than stable under a wider equilibrium that considers more general principles of theory choice, e.g. fruitfulness, etc. The latter is just the anti-exceptionalist, Quinean view.

Carroll/Quine-type considerations do not provide support for excluding a wide equilibrium view, nor an argument against the possibility of habituation or the malleability of reasoning.

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