Mathematical Logic



Largest initial segments pointwise fixed by automorphisms of models of set theory

Ali Enayat¹ · Matt Kaufmann² · Zachiri McKenzie¹

Received: 12 June 2016 / Accepted: 7 November 2016 / Published online: 12 September 2017 © The Author(s) 2017. This article is an open access publication

Abstract Given a model \mathcal{M} of set theory, and a nontrivial automorphism j of \mathcal{M} , let $\mathcal{I}_{\mathrm{fix}}(j)$ be the submodel of \mathcal{M} whose universe consists of elements m of \mathcal{M} such that j(x) = x for every x in the transitive closure of m (where the transitive closure of m is computed within \mathcal{M}). Here we study the class \mathcal{C} of structures of the form $\mathcal{I}_{\mathrm{fix}}(j)$, where the ambient model \mathcal{M} satisfies a frugal yet robust fragment of ZFC known as MOST, and j(m) = m whenever m is a finite ordinal in the sense of \mathcal{M} . Our main achievement is the calculation of the theory of \mathcal{C} as precisely MOST + $\Delta_0^{\mathcal{P}}$ -Collection. The following theorems encapsulate our principal results:

Theorem A. Every structure in C satisfies $MOST + \Delta_0^{\mathcal{P}}$ -Collection.

Theorem B. Each of the following three conditions is sufficient for a countable structure \mathcal{N} to be in \mathcal{C} :

- (a) \mathcal{N} is a transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -Collection.
- (b) \mathcal{N} is a recursively saturated model of MOST + $\Delta_0^{\mathcal{P}}$ -Collection.
- (c) \mathcal{N} is a model of ZFC.

Theorem C. Suppose \mathcal{M} is a countable recursively saturated model of ZFC and I is a proper initial segment of $\mathrm{Ord}^{\mathcal{M}}$ that is closed under exponentiation and contains $\omega^{\mathcal{M}}$. There is a group embedding $j \mapsto \check{j}$ from $\mathrm{Aut}(\mathbb{Q})$ into $\mathrm{Aut}(\mathcal{M})$ such that I is

> Matt Kaufmann kaufmann@cs.utexas.edu

Zachiri McKenzie zach.mckenzie@gmail.com

Department of Computer Science, University of Texas at Austin, Austin, TX, USA



Department of Philosophy, Linguistics and Theory of Science, Gothenburg University, Gothenburg, Sweden

the longest initial segment of $\mathrm{Ord}^{\mathcal{M}}$ that is pointwise fixed by \check{j} for every nontrivial $j \in \mathrm{Aut}(\mathbb{Q})$.

In Theorem C, Aut(X) is the group of automorphisms of the structure X, and \mathbb{Q} is the ordered set of rationals.

Keywords Automorphisms · Models of set theory · Recursively saturated

Mathematics Subject Classification Primary 03C62; Secondary 03H9

1 Introduction

We study automorphisms of models of set theory, focusing on structures of the form $\mathcal{I}_{fix}(j)$, where j is a nontrivial automorphism of a model \mathcal{M} of MOST (Mostowski Set Theory) such that j(m) = m whenever m is a finite ordinal in the sense of \mathcal{M} , and $\mathcal{I}_{fix}(j)$ is the submodel of \mathcal{M} whose universe consists of elements m in M such that j(x) = x for every x in the transitive closure of m (where the transitive closure is calculated within \mathcal{M}). MOST is a frugal yet robust fragment of ZFC, introduced and studied in Mathias' majestic paper [17] (see Definition 2.4).

A principal source of motivation for our work is to be found in the study of automorphisms of models of ZF's sister theory PA (Peano arithmetic). For example: by a theorem of Smoryński ([23], [14, Theorem 8.4.2]), if \mathcal{I} is submodel of a countable recursively saturated model of PA whose elements form a proper initial segment of \mathcal{M} , and \mathcal{I} is closed under the exponential function of \mathcal{M} , then there is an automorphism j of \mathcal{M} such that $\mathcal{I}_{\mathrm{fix}}(j) = \mathcal{I}$, where $\mathcal{I}_{\mathrm{fix}}(j)$ is defined in this context as the submodel of \mathcal{M} whose universe consists of element m such that j(x) = x for all $x < \mathcal{M} m$. Furthermore, it is known that:

- (1) [4, Lemmas A.0 & A.2] If \mathcal{M} is a model of the fragment $I\Delta_0$ of PA, and j is a nontrivial automorphisms of \mathcal{M} , then $\mathcal{I}_{fix}(j)$ satisfies $I\Delta_0 + Exp + B\Sigma_1$ (where Exp expresses the totality of the exponential function, and $B\Sigma_1$ is the scheme of Σ_1 -Collection).
- (2) [4, Theorem A] Every countable model of $I\Delta_0 + Exp + B\Sigma_1$ arises as $\mathcal{I}_{fix}(j)$ for some nontrivial automorphism j of a model of $I\Delta_0$.

Our Theorem 3.4 is an analogue of (1); while Theorems 5.5, 5.6, and 5.15 provide analogues of (2). More generally, our results can be viewed as contributing to the project initiated in [3] of investigating the extent to which core results about automorphisms of model of arithmetic can be extended to the set-theoretic realm.

Another source of inspiration for our work is the metamathematics of NFU, an urelement-variant of Quine's system "New Foundations", NF. Jensen's pioneering work [11] unveiled a magical link between models of NFU and automorphisms of models of ZF-style set theories, a link that has captured the imagination of other researchers, e.g., Holmes [9], Solovay [24], and two of the authors of the present paper ([3,16]). Our results here have a number of implications for NFU, for example, Theorem 5.15 can be used to show that every countable model of ZFC can be realized as the strongly cantorian part of a model of NFU. However, the precise implications of our results to the NFU setting is yet to be worked out and will be pursued elsewhere.



The plan of the paper is as follows. After dealing with preliminaries in Sect. 2, we introduce the key notion of an "H-cut" of a model of set theory in Sect. 3, where we establish two important facts, namely: (1) every H-cut of a model of MOST satisfies $MOST + \Delta_0^{\mathcal{P}}$ -Collection; and (2) $\mathcal{I}_{fix}(j)$ is an H-cut of \mathcal{M} , if \mathcal{M} is a model of MOST, and j is a nontrivial automorphisms of \mathcal{M} such that $\mathcal{I}_{fix}(j)$ includes $\omega^{\mathcal{M}}$ (equivalently: j(m) = m whenever whenever m is a finite ordinal in the sense of \mathcal{M}). The central result of Sect. 4 is Theorem 4.1, whose iterated ultrapower proof involves a rather intricate set-theoretical adaptation of a machinery that was invented in the arithmetical context by Paris and Mills ([19], [14, Theorem 3.5.5]) and was further elaborated in [4]. One of the remarkable consequences of Theorem 4.1 is that if \mathcal{I} is an H-cut of a countable model \mathcal{M} of MOST + $\Delta_0^{\mathcal{P}}$ -Collection, then \mathcal{M} has a cofinal extension \mathcal{N} that (1) \mathcal{N} does not add any new members to elements of I; (2) \mathcal{N} satisfies MOST; and (3) \mathcal{N} carries an automorphism j such that $\mathcal{I} = \mathcal{I}_{fix}(j)$. This consequence of Theorem 4.1 is put to work together with a key construction in Sect. 5 (Theorem 5.6) to show that every countable recursively saturated model of MOST + $\Delta_0^{\mathcal{P}}$ -Collection can be realized as $\mathcal{I}_{fix}(j)$; a result that, together with our work in Sect. 3, yields the central theorem of our paper (Theorem 5.8) that identifies the theory of the class of models of the form $\mathcal{I}_{fix}(j)$ to be precisely MOST + $\Delta_0^{\mathcal{P}}$ -Collection. In Sect. 5 we also use our work in Sect. 4 together with some classical results of Friedman [8] and Hutchinson [10] to identify two other sufficient conditions for a countable model \mathcal{N} to be realizable as $\mathcal{I}_{fix}(j)$, namely: (1) \mathcal{N} is a transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -Collection and (2) \mathcal{N} is a model of ZFC. Finally, in Sect. 6, we fine-tune a theorem of Togha [25] on automorphisms of countable recursively saturated models of ZFC in a manner reminiscent of a refinement of Smoryński's aforementioned theorem established in [4, Theorem B].

2 Background and definitions

Throughout this paper \mathcal{L} will denote the language of set theory—first-order logic endowed with a binary relation symbol ∈ whose intended interpretation is membership. Structures will usually be denoted using upper-case calligraphy roman letters $(\mathcal{M}, \mathcal{N}, \ldots)$ and the corresponding plain font letter (M, N, \ldots) will be used to denote the underlying set of that structure. If \mathcal{M} is an \mathcal{L}' -structure where $\mathcal{L}' \supseteq \mathcal{L}$ and $a \in M$ then we will use a^* to denote the class $\{x \in M \mid \mathcal{M} \models (x \in a)\}$. As usual $\Delta_0, \Sigma_1, \Pi_1, \ldots$ with be used to denote the Lévy classes of \mathcal{L} -formulae. We will also have cause to consider the class $\Delta_0^{\mathcal{P}}$ which is the smallest class of \mathcal{L} -formulae that contains all atomic formulae, contains all compound formulae formed using the connectives of first-order logic, and is closed under quantification in the form $Qx \in y$ and $Qx \subseteq y$ where x and y are distinct variables and Q is \exists or \forall . If $\mathcal{L}' \supset \mathcal{L}$ then we use $\Delta_0(\mathcal{L}')$ ($\Delta_0^{\mathcal{P}}(\mathcal{L}')$) to denote the smallest class of formulae that contains all atomic formulae, all compound formulae formed using the connectives of first-order logic, and is closed under quantification in the form $Qx \in t$ (and $Qx \subseteq t$) where t is an \mathcal{L}' -term and x is a variable that does not appear in t, and \mathcal{Q} is \exists or \forall . The classes $\Sigma_1^{\mathcal{P}}, \Pi_1^{\mathcal{P}}, \dots (\Sigma_1(\mathcal{L}'), \Pi_1(\mathcal{L}'), \dots$ and $\Sigma_1^{\mathcal{P}}(\mathcal{L}'), \Pi_1^{\mathcal{P}}(\mathcal{L}'), \dots)$ are defined inductively from the class $\Delta_0^{\mathcal{P}}$ (respectively $\Delta_0(\mathcal{L}')$ and $\Delta_0^{\mathcal{P}}(\mathcal{L}')$) in the same way as the classes Σ_1, Π_1, \ldots are defined from Δ_0 .



Let $\mathcal{L}'\supseteq\mathcal{L}$ and let \mathcal{M} and \mathcal{N} be \mathcal{L}' -structures. If \mathcal{M} is a substructure of \mathcal{N} then we will write $\mathcal{M}\subseteq\mathcal{N}$. If Γ is a class of \mathcal{L}' -formulae then we will write $\mathcal{M}\prec_{\Gamma}\mathcal{N}$ if $\mathcal{M}\subseteq\mathcal{N}$ and for every $\vec{a}\in\mathcal{M}$, \vec{a} satisfies the same Γ -formulae in both \mathcal{M} and \mathcal{N} . If Γ is \mathcal{L}' or $\Sigma_n(\mathcal{L}')$ then we will abbreviate this notation by writing $\mathcal{M}\prec\mathcal{N}$ and $\mathcal{M}\prec_n\mathcal{N}$ respectively. If $\mathcal{M}\subseteq\mathcal{N}$ and for all $x\in\mathcal{M}$ and $y\in\mathcal{N}$,

if
$$\mathcal{N} \models (y \in x)$$
 then $y \in M$,

then we say that \mathcal{N} is an *end-extension* of \mathcal{M} and write $\mathcal{M} \subseteq_e \mathcal{N}$. It is well-known that if $\mathcal{M} \subseteq_e \mathcal{N}$ then $\mathcal{M} \prec_{\Delta_0} \mathcal{N}$. If \mathcal{N} is an end-extension of \mathcal{M} and $\mathcal{M} \prec \mathcal{N}$ then we write $\mathcal{M} \prec_e \mathcal{N}$. In contrast, if $\mathcal{M} \subseteq \mathcal{N}$ and for all $x \in \mathcal{N}$, there exists $y \in \mathcal{M}$ such that $\mathcal{N} \models (x \in y)$, then we say that \mathcal{N} is a *cofinal extension* of \mathcal{M} and we write $\mathcal{M} \subseteq_{\operatorname{cf}} \mathcal{N}$. And, if $\mathcal{M} \subseteq_{\operatorname{cf}} \mathcal{N}$ and $\mathcal{M} \prec \mathcal{N}$ then we write $\mathcal{M} \prec_{\operatorname{cf}} \mathcal{N}$.

Let $\mathcal{L}' \supseteq \mathcal{L}$ and let Γ be a class of \mathcal{L}' formulae. We will use Γ -separation and Γ -collection to abbreviate the separation and collection schemes, respectively, restricted to formulae in Γ . We will also use the following axioms and schemes to axiomatize and study weak variants of ZFC:

(transitive containment) $\forall x \exists y (\bigcup y \subseteq y \land x \subseteq y)$,

 $(\forall \kappa \exists \kappa^+)$ for every initial ordinal κ , there exists a least initial ordinal $> \kappa$,

 $(\forall \kappa (2^{\kappa} \text{ exists}))$ for every initial ordinal κ , there exists an initial λ such that $\lambda = |\mathcal{P}(\kappa)|$,

(Γ-foundation) for all Γ-formulae $\phi(x, \vec{z})$,

$$\forall \vec{z} (\exists x \phi(x, \vec{z}) \Rightarrow \exists y (\phi(y, \vec{z}) \land (\forall x \in y) \neg \phi(x, \vec{z})))$$

When Γ only contains the formula $x \in z$ then we will refer to the single axiom in this scheme as *set foundation*.

The following subsystems of ZFC are studied in [17]:

Definition 2.1 Mac is the \mathcal{L} -theory with axioms: extensionality, pair, emptyset, union, infinity, powerset, transitive containment, Δ_0 -separation, set foundation, and the axiom of choice in the form: every set can be well-ordered.

Definition 2.2 KP is the \mathcal{L} -theory with axioms: extensionality, pair, emptyset, union, Δ_0 -separation, Δ_0 -collection and Π_1 -foundation.

Definition 2.3 KP^{\mathcal{P}} is the \mathcal{L} -theory with axioms: extensionality, pair, emptyset, union, infinity, powerset, Δ_0 -separation, $\Delta_0^{\mathcal{P}}$ -collection and $\Pi_1^{\mathcal{P}}$ -foundation.

Note that Mac proves that for every set x, there exists a unique smallest transitive set, which we will denote TC(x), that contains x. We will use Ord and Card to denote the classes (definable in any extension of Mac) of ordinals and cardinals respectively. We record the following well-known facts about fragments of the collection scheme:

Lemma 2.1 *Let* $n \in \omega$.

(i) Mac + Π_n -collection $\vdash \Sigma_{n+1}$ -collection,



(ii) Mac +
$$\Pi_n^{\mathcal{P}}$$
-collection $\vdash \Sigma_{n+1}^{\mathcal{P}}$ -collection.

Definition 2.4 *The theory* MOST *is obtained from* Mac *by adding* Σ_1 *-separation and* Δ_0 *-collection.*

Mathias [17] extensively studies the following axiom, originally proposed by Mitchell [18], which asserts that for every set u, there exists a universal transitive set which contains every transitive set that is of size $\leq |u|$:

(Axiom H)
$$\forall u \exists T (| T \subseteq T \land \forall z (| Z \subseteq z \land |z| \le |u| \Rightarrow z \subseteq T)).$$

Mathias [17, Proposition 3.14] shows that adding Axiom H to Mac yields MOST.

Lemma 2.2 Mac
$$+$$
 Axiom H $=$ MOST.

The set theory MOST is capable of defining the rank function, which we will denote ρ . It should be noted, however, that MOST does not prove enough recursion to ensure the totality of the function $\alpha \mapsto V_{\alpha}$. The following consequences of MOST are proved in [17, Theorem scheme 3.13, Proposition 3.14, Theorem 3.18]:

Lemma 2.3 MOST proves the following:

- (i) all instances of $\Delta_0^{\mathcal{P}}$ -separation,
- (ii) all instances of Π_1 -separation,
- (iii) all instances of Σ_1 -collection,
- (iv) all instances of Π_1 -foundation,
- (v) every well-ordering is isomorphic to an ordinal,
- (vi) $\forall \kappa \exists \kappa^+$,
- (vii) for all cardinals κ , there is a set H_{κ} of all sets whose transitive closure has cardinality less than κ .
- (v) is a special case of what is known as Mostowski Isomorphism Theorem. Mathias [17, Lemma 3.15] also proves that a more general special case of this theorem which deals with well-founded extensional relations that are sets is provable in MOST.

Lemma 2.4 MOST proves that if $R \subseteq X \times X$ is well-founded and extensional then the function $\varpi_R : X \longrightarrow V$ defined by

$$\varpi_R(x) = \{ \varpi_R(y) \mid y \in X \land (\langle y, x \rangle \in R) \}$$

is defined on all of X. Moreover, ϖ_R "X is the unique transitive set that is isomorphic to R.

In [7] Thomas Forster and Richard Kaye introduce the notion of a powerset preserving end-extension. Here we give a slightly more general version of this notion that does not require the powerset axiom to hold in the structures that are being compared.

Definition 2.5 *Let* \mathcal{M} *and* \mathcal{N} *are* \mathcal{L}' -structures where $\mathcal{L}' \supseteq \mathcal{L}$. We say that \mathcal{N} is a powerset preserving end-extension of \mathcal{M} and write $\mathcal{M} \subseteq_{e}^{\mathcal{P}} \mathcal{N}$ if

- (i) $\mathcal{M} \subseteq_{e} \mathcal{N}$
- (ii) for all $x \in N$ and for all $y \in M$, if $\mathcal{N} \models (x \subseteq y)$ then $x \in M$.



Just as end-extensions preserve Δ_0 properties, powerset preserving end-extensions preserve $\Delta_0^{\mathcal{P}}$ properties. The following is a slight modification of a result that appears in [7]:

Lemma 2.5 *Let* \mathcal{M} *and* \mathcal{N} *be* \mathcal{L} -structures that satisfy extensionality. If $\mathcal{M} \subseteq_e^{\mathcal{P}} \mathcal{N}$ then $\mathcal{M} \prec_{\Lambda_e^{\mathcal{P}}} \mathcal{N}$.

Proof A straightforward induction on the structural complexity of a $\Delta_0^{\mathcal{P}}$ -formula ϕ .

The next definition captures an important notion that plays a key role in this study.

Definition 2.6 Let \mathcal{M} and \mathcal{N} be \mathcal{L} -structures. We say that \mathcal{N} is a topless end-extension of \mathcal{M} and write $\mathcal{M} \subseteq_{\mathsf{topless}} \mathcal{N}$ if

- (i) $\mathcal{M} \subseteq_e \mathcal{N}$,
- (ii) $M \neq N$,
- (iii) if $C \in N$ and $C^* \subseteq \operatorname{Ord}^{\mathcal{M}}$ then $C \in M$.

If $\mathcal{M} \subseteq_{\text{topless}} \mathcal{N}$ and $\mathcal{M} \subseteq_{e}^{\mathcal{P}} \mathcal{N}$ then we say that \mathcal{N} is a topless powerset preserving end-extension of \mathcal{M} and write $\mathcal{M} \subseteq_{\text{topless}}^{\mathcal{P}} \mathcal{N}$.

Note that if the \mathcal{M} and \mathcal{N} in Definition 2.6 satisfy MOST then condition (iii) is a paraphrasing of the assertion that there is no least ordinal in $Ord^{\mathcal{N}} \setminus Ord^{\mathcal{M}}$.

This paper studies the largest transitive initial segment of a model of MOST that is pointwise fixed by a non-trivial automorphism.

Definition 2.7 Let \mathcal{M} be an \mathcal{L} -structure with $\mathcal{M} \models \text{MOST}$ and let $j : \mathcal{M} \longrightarrow \mathcal{M}$ be a non-trivial automorphism. We use fix(j) to denote the class of fixed points of j. *I.e.*

$$fix(j) = \{x \in M \mid j(x) = x\}.$$

Define $\mathcal{I}_{fix}(j)$ to be the substructure of \mathcal{M} with domain

$$I_{\text{fix}}(j) = \{x \in M \mid \forall y (\mathcal{M} \models (y \in \text{TC}(\{x\})) \Rightarrow j(y) = y)\}.$$

Let \mathcal{M} be an \mathcal{L}' -structure. The structure \mathcal{M} is said to be *recursively saturated* if for all $a_1, \ldots, a_n \in \mathcal{M}$ and for all recursive finitely realised types $\Gamma(x_1, \ldots, x_m, a_1, \ldots, a_n)$ in the language \mathcal{L}' with parameters a_1, \ldots, a_n , \mathcal{M} realises $\Gamma(x_1, \ldots, x_m, a_1, \ldots, a_n)$. We refer the reader to [2, §2.4] for a detailed treatment of recursively saturated models. We will make use of the following nice feature of recursive saturation:

Theorem 2.6 (see [2]) Let \mathcal{L}' be a recursive language. If T is a consistent \mathcal{L}' -theory with an infinite model then T has a countable recursively saturated model. \Box

Fix a Gödel coding of \mathcal{L} in the theory MOST and use Form to denote the set of Gödel codes of well-formed \mathcal{L} -formulae. Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be an \mathcal{L} -structure with $\mathcal{M} \models \text{MOST}$. A satisfaction class for \mathcal{M} is a class $S \subseteq M$ such that S consists of ordered pairs $\langle a, b \rangle$ where $a \in (\text{Form}^{\mathcal{M}})^*$ and $b \in M$, and for all $n \in \omega$,



(S is n-correct) S satisfies Tarski's inductive conditions for truth for all Σ_n -formulae.

Let \mathcal{L}_X be the extension of \mathcal{L} obtained by adding a new unary predicate X. Use ZFC(X) to denote the \mathcal{L}_X -theory that extends ZFC with the schemes of \mathcal{L}_X -separation and \mathcal{L}_X -collection. The following result of Schlipf provides another characterisation of recursive saturation for countable models of ZFC; the theorem below is an immediate consequence of putting Theorem 3.4 and the remark following it in [20] with the well-known resplendence property of countable recursively saturated models (see [20, Theorem 1.3]).

Theorem 2.7 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be a countable ω -nonstandard model of ZFC. The structure \mathcal{M} is recursively saturated if and only if \mathcal{M} can be expanded to an \mathcal{L}_X -structure $\mathcal{M}_{Sat} = \langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$ such that

- (I) $\mathcal{M}_{Sat} \models ZFC(X)$,
- (II) $X^{\mathcal{M}}$ is a satisfaction class for \mathcal{M} .

Note that we cannot expect that a satisfaction class obtained from Theorem 2.7 will be n-correct for all internal natural numbers of \mathcal{M} , as this would prove the consistency of ZFC. However, if $\mathcal{M} = \langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$ is an \mathcal{L}_X -structure that satisfies (I) and (II) of Theorem 2.7 and is ω -nonstandard then there is a nonstandard $s \in (\omega^{\mathcal{M}})^*$ such that $X^{\mathcal{M}}$ is s-correct. This follows from overspill and the fact that there is an \mathcal{L}_X -formula with parameter n that expresses that X is n-correct.

3 The structure and first-order theory of $\mathcal{I}_{fix}(j)$

In this section we investigate the first-order theory of $\mathcal{I}_{\mathrm{fix}}(j)$ and the properties of this structure in relation to the domain of the automorphism j. We show that if $j:\mathcal{M}\longrightarrow\mathcal{M}$ is a non-trivial automorphism that hereditarily fixes $\omega^{\mathcal{M}}$ and \mathcal{M} satisfies MOST then $\mathcal{I}_{\mathrm{fix}}(j)$ satisfies MOST and \mathcal{M} is a topless powerset preserving endextension of $\mathcal{I}_{\mathrm{fix}}(j)$. The fact that \mathcal{M} is a topless powerset preserving end-extension of $\mathcal{I}_{\mathrm{fix}}(j)$ implies that $\mathcal{I}_{\mathrm{fix}}(j)$ also satisfies all instances of $\Delta_0^{\mathcal{P}}$ -collection. We capture the relationship that we will prove holds between $\mathcal{I}_{\mathrm{fix}}(j)$ and \mathcal{M} in the following definition:

Definition 3.1 *Let* \mathcal{M} *be an* \mathcal{L} -*structure with* $\mathcal{M} \models MOST$. *We say that* $I \subseteq M$ *is an* H-*cut of* \mathcal{M} *if*

- $(i) \ \langle I, \in^{\mathcal{M}} \rangle \subseteq_{\text{topless}}^{\mathcal{P}} \mathcal{M},$
- (ii) $\langle I, \in^{\mathcal{M}} \rangle \models MOST$.

Note that if I is an H-cut of a model \mathcal{M} of MOST then I is a union of sets H_{κ} in \mathcal{M} . We make this explicit with the following observation:

Lemma 3.1 *Let* \mathcal{M} *be an* \mathcal{L} -*structure with* $\mathcal{M} \models MOST$. *If* $I \subseteq M$ *is an* H-*cut of* \mathcal{M} *then for all* $\kappa \in I$,

$$\langle I, \in^{\mathcal{M}} \rangle \models (\kappa \text{ is a cardinal}) \text{ if and only if } \mathcal{M} \models (\kappa \text{ is a cardinal})$$

and if $\mathcal{M} \models (\kappa \text{ is a cardinal}) \text{ then } H_{\kappa}^{\mathcal{M}} = H_{\kappa}^{\langle I, \in^{\mathcal{M}} \rangle}.$



Proof Let $I \subseteq M$ be an H-cut of \mathcal{M} and let $\kappa \in I$. The expression that ' κ is a cardinal' can be written as a $\Delta_0^{\mathcal{P}}$ -formula which takes the parameter $\kappa \times \kappa$. This shows that

$$\langle I, \in^{\mathcal{M}} \rangle \models (\kappa \text{ is a cardinal}) \text{ if and only if } \mathcal{M} \models (\kappa \text{ is a cardinal}).$$

The formula $y = \mathrm{TC}(x)$ can be expressed as a $\Delta_0^{\mathcal{P}}$ -formula. The fact that $(H_{\kappa}^{\langle \mathcal{I}, \in^{\mathcal{M}} \rangle})^* \subseteq (H_{\kappa}^{\mathcal{M}})^*$ now follows from the fact that $x \in H_{\kappa}$ can be expressed as a $\Delta_0^{\mathcal{P}}$ -formula with parameters $x \times \kappa$ and $\mathrm{TC}(x)$. To get the reverse inclusion, let $x \in (H_{\kappa}^{\mathcal{M}})^*$. All we need to show is that $x \in I$. Work inside \mathcal{M} . Since $x \in H_{\kappa}$, there is an $A \subseteq \kappa$ and a well-founded and extensional $R \subseteq A \times A$ such that ϖ_R " $A = \mathrm{TC}(\{x\})$. Since $\langle I, \in^{\mathcal{M}} \rangle \subseteq_e^{\mathcal{P}} \mathcal{M}$, it follows that $A, R \in I$ and that $\langle I, \in^{\mathcal{M}} \rangle$ thinks R is well-founded and extensional. Therefore

$$\langle I, \in^{\mathcal{M}} \rangle \models (\varpi_R \text{``A is a transitive set}).$$

Set foundation can then be used to show that $(\varpi_R {}^{u}A)^{\mathcal{M}} = (\varpi_R {}^{u}A)^{\langle I, \in^{\mathcal{M}} \rangle}$. This shows that $x \in I$ and concludes the proof that $H_{\kappa}^{\mathcal{M}} = H_{\kappa}^{\langle I, \in^{\mathcal{M}} \rangle}$.

Togha [25, Definition 1] introduces the notion of a cut of the ordinals of a model of set theory.

Definition 3.2 Let \mathcal{M} be an \mathcal{L} -structure with $\mathcal{M} \models \text{MOST}$. We say that $K \subseteq \text{Ord}^{\mathcal{M}}$ is an ordinal cut of \mathcal{M} closed under exponentiation if $\langle K, \in^{\mathcal{M}} \rangle \subseteq_{topless} \langle \text{Ord}^{\mathcal{M}}, \in^{\mathcal{M}} \rangle$ and for all \mathcal{M} -cardinals $\kappa \in K$, $(2^{\kappa})^{\mathcal{M}} \in K$.

If $\mathcal{M} \models \text{MOST}$ and $I \subseteq M$ is an H-cut of \mathcal{M} then $\text{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}$ is an ordinal cut of \mathcal{M} that is closed under exponentiation and, moreover, $\text{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}$ completely determines I.

Lemma 3.2 Let \mathcal{M} be an \mathcal{L} -structure with $\mathcal{M} \models MOST$. If $I \subseteq M$ is an H-cut of \mathcal{M} then

- (I) $\operatorname{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}$ is an ordinal cut of \mathcal{M} closed under exponentiation,
- (II) $I = \bigcup \{ (H_{\kappa}^{\mathcal{M}})^* \mid (\mathcal{M} \models (\kappa \text{ is a cardinal})) \land (\kappa \in I) \}.$

Proof Let $I \subseteq M$ be an H-cut of \mathcal{M} . It is clear that $\langle \operatorname{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}, \in^{\mathcal{M}} \rangle \subseteq_{topless} \langle \operatorname{Ord}^{\mathcal{M}}, \in^{\mathcal{M}} \rangle$. Let $\kappa \in I$ be an \mathcal{M} -cardinal. It follows from Lemma 3.1 that κ is an $\langle I, \in^{\mathcal{M}} \rangle$ -cardinal and $(2^{\kappa})^{\mathcal{M}} = (2^{\kappa})^{\langle I, \in^{\mathcal{M}} \rangle}$. Therefore $(2^{\kappa})^{\mathcal{M}} \in I$. This shows that $\operatorname{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}$ is an ordinal cut of \mathcal{M} closed under exponentiation. Let $J = \bigcup \{(H_{\kappa}^{\mathcal{M}})^* \mid (\mathcal{M} \models (\kappa \text{ is a cardinal})) \wedge (\kappa \in I)\}$. Lemma 3.1 shows that $J \subseteq I$. Conversely, let $x \in I$. Let $\kappa = |\mathcal{P}^{\langle I, \in^{\mathcal{M}} \rangle}(\operatorname{TC}(x))| \in I$. Therefore, by Lemma 3.1, κ is an \mathcal{M} -cardinal and $x \in (H_{\kappa}^{\langle I, \in^{\mathcal{M}} \rangle})^* = (H_{\kappa}^{\mathcal{M}})^*$. This shows that I = J and completes the proof of the lemma.

The fact that an H-cut sits toplessly inside a powerset preserving end-extension that satisfies MOST implies that it satisfies all instances of $\Delta_0^{\mathcal{P}}$ -collection.



Lemma 3.3 *Let* \mathcal{M} *be an* \mathcal{L} -structure with $\mathcal{M} \models MOST$. *If* $I \subseteq M$ *is an* H-cut of \mathcal{M} *then*

$$\langle I, \in^{\mathcal{M}} \rangle \models \Delta_0^{\mathcal{P}}$$
-collection.

Proof Let $\phi(x, y, \vec{z})$ be a $\Delta_0^{\mathcal{P}}$ -formula and let $\vec{a}, b \in I$ such that

$$\langle I, \in^{\mathcal{M}} \rangle \models (\forall x \in b) \exists y \phi(x, y, \vec{a}).$$

Let $\kappa \in \operatorname{Ord}^{\mathcal{M}} \backslash \operatorname{Ord}^{\langle I, \in^{\mathcal{M}} \rangle}$ be a cardinal. Note that $I \subseteq (H_{\kappa^+}^{\mathcal{M}})^*$. Define $\theta(u, w, w', \vec{v})$ to be the formula

$$(u \text{ is an ordinal}) \land (\exists f, T, y \in w)(\exists x \in w')$$

$$\begin{pmatrix} (T = TC(y)) \land (f : u \longrightarrow T \text{ is a bijection}) \land \phi(x, y, \vec{v}) \land \\ (p \text{ is an ordinal}) \land (S = TC(q)) \land \\ (g : p \longrightarrow S \text{ is a bijection}) \land \phi(x, q, \vec{v}) \end{pmatrix} \land \psi \leq p$$

So $\theta(u, w, w', \vec{v})$ is a $\Delta_0^{\mathcal{P}}$ formula. Working inside \mathcal{M} , $\Delta_0^{\mathcal{P}}$ -separation ensures that

$$C = \{\lambda \in H_{\kappa^+} \mid \theta(\lambda, H_{\kappa^+}, b, \vec{a})\}\$$

is a set. Lemma 3.1 implies that $C^* \subseteq I$. Therefore, since $\langle I, \in^{\mathcal{M}} \rangle \subseteq_{\text{topless}}^{\mathcal{P}} \mathcal{M}, C \in I$. Working inside $\langle I, \in^{\mathcal{M}} \rangle$, let $\mu = \sup C$. Now, μ is a cardinal and

$$\langle I, \in^{\mathcal{M}} \rangle \models (\forall x \in b) (\exists y \in H_{\mu^+}) \phi(x, y, \vec{a}).$$

The remainder of this section is dedicated to proving that if $j: \mathcal{M} \longrightarrow \mathcal{M}$ is a non-trivial automorphism of a model of MOST that hereditarily fixes $\omega^{\mathcal{M}}$ then $I_{\mathrm{fix}}(j)$ is an H-cut of \mathcal{M} .

Theorem 3.4 Let \mathcal{M} be an \mathcal{L} -structure with $\mathcal{M} \models \text{MOST}$. If $j : \mathcal{M} \longrightarrow \mathcal{M}$ is a non-trivial automorphism with $\omega^{\mathcal{M}} \in I_{\text{fix}}(j)$ then $I_{\text{fix}}(j) \subseteq \mathcal{M}$ is an \mathcal{H} -cut of \mathcal{M} .

Combined with Lemma 3.3 this shows that if $j:\mathcal{M}\longrightarrow\mathcal{M}$ is an automorphism that satisfies the conditions of Theorem 3.4 then $\mathcal{I}_{\mathrm{fix}}(j)$ satisfies MOST plus the scheme of $\Delta_0^{\mathcal{P}}$ -collection.

Corollary 3.5 *Let* \mathcal{M} *be an* \mathcal{L} -structure with $\mathcal{M} \models MOST$. If $j : \mathcal{M} \longrightarrow \mathcal{M}$ is a non-trivial automorphism with $\omega^{\mathcal{M}} \in I_{fix}(j)$ then

$$\mathcal{I}_{\text{fix}}(j) \models \text{MOST} + \Delta_0^{\mathcal{P}}$$
-collection.



For the remainder of this section we let $j: \mathcal{M} \longrightarrow \mathcal{M}$ be a non-trivial automorphism where $\mathcal{M} \models \text{MOST}$ and j(x) = x for all $x \in (\omega^{\mathcal{M}})^*$. We first show that \mathcal{M} is a topless powerset preserving end-extension of $\mathcal{I}_{\text{fix}}(j)$.

Lemma 3.6 $\mathcal{I}_{fix}(j)$ satisfies extensionality and powerset, and $\mathcal{I}_{fix}(j) \subseteq_{topless}^{\mathcal{P}} \mathcal{M}$.

Proof It is clear from the definition of $\mathcal{I}_{\mathrm{fix}}(j)$ that $I_{\mathrm{fix}}(j)$ is a transitive proper subclass of \mathcal{M} . It follows that $\mathcal{I}_{\mathrm{fix}}(j) \models (\mathrm{extensionality})$ and $I_{\mathrm{fix}}(j) \neq M$. Let $y \in I_{\mathrm{fix}}(j)$. Let $x \in M$ be such that $\mathcal{M} \models (x \subseteq y)$. Since $y \in I_{\mathrm{fix}}(j)$, j fixes every element of x^* . Since j is automorphism it follows that j(x) = x and so $x \in I_{\mathrm{fix}}(j)$. This shows that for all $x \in \mathcal{P}^{\mathcal{M}}(y)^*$, $x \in I_{\mathrm{fix}}(j)$. And so $\mathcal{P}^{\mathcal{M}}(y) \in I_{\mathrm{fix}}(j)$. Since $I_{\mathrm{fix}}(j)$ is transitive, this shows that $\mathcal{I}_{\mathrm{fix}}(j) \models (\text{powerset})$ and $\mathcal{I}_{\mathrm{fix}}(j) \subseteq_e^{\mathcal{P}} \mathcal{M}$. To show that $\mathcal{I}_{\mathrm{fix}}(j)$ sits toplessly in \mathcal{M} , let $C \in M$ with $C^* \subseteq \mathrm{Ord}^{\mathcal{I}_{\mathrm{fix}}(j)}$. It immediately follows from the fact that j is an automorphism that j(C) = C. Therefore $C \in I_{\mathrm{fix}}(j)$. This completes the proof that $\mathcal{I}_{\mathrm{fix}}(j) \subseteq_{\mathrm{topless}}^{\mathcal{P}} \mathcal{M}$.

To complete the proof that $\mathcal{I}_{fix}(j)$ is an H-cut of \mathcal{M} we need to show that $\mathcal{I}_{fix}(j)$ satisfies MOST.

Lemma 3.7 $\mathcal{I}_{fix}(j) \models Mac$.

Proof Lemma 3.6 and the fact that $I_{\mathrm{fix}}(j)$ is a transitive subclass of \mathcal{M} implies that extensionality, union, pair, emptyset, powerset, Δ_0 -separation, set foundation and the axiom of choice hold in $\mathcal{I}_{\mathrm{fix}}(j)$. Since $\omega^{\mathcal{M}} \in I_{\mathrm{fix}}(j)$, $\mathcal{I}_{\mathrm{fix}}(j)$ satisfies the axiom of infinity. To see that transitive containment holds in $\mathcal{I}_{\mathrm{fix}}(j)$, let $x \in I_{\mathrm{fix}}(j)$. Since j fixes every element of $\mathrm{TC}^{\mathcal{M}}(x)^*$, j also fixes $\mathrm{TC}^{\mathcal{M}}(x)$. Therefore $\mathrm{TC}^{\mathcal{M}}(x) \in I_{\mathrm{fix}}(j)$ and $\mathcal{I}_{\mathrm{fix}}(j)$ believes that $\mathrm{TC}^{\mathcal{M}}(x)$ is a transitive set which contains x.

It remains to show that $\mathcal{I}_{fix}(j)$ satisfies axiom H.

Lemma 3.8 For all $R, X \in I_{fix}(j)$,

$$\mathcal{M} \models (R \subseteq X \times X \text{ is well-founded and extensional})$$
 if and only if $\mathcal{I}_{fix}(j) \models (R \subseteq X \times X \text{ is well-founded and extensional})$

and if

$$\mathcal{M} \models (R \subseteq X \times X \text{ is well-founded and extensional})$$

then $(\varpi_R "X)^{\mathcal{M}} \in I_{\mathrm{fix}}(j)$.

Proof Let $R, X \in I_{fix}(j)$. The fact that

$$\mathcal{M} \models (R \subseteq X \times X \text{ is well-founded and extensional})$$
 if and only if $\mathcal{I}_{fix}(j) \models (R \subseteq X \times X \text{ is well-founded and extensional}),$

follows from the fact that ' $R \subseteq X \times X$ is well-founded and extensional' can be expressed as a $\Delta_0^{\mathcal{P}}$ -formula. Assume that

$$\mathcal{M} \models (R \subseteq X \times X \text{ is well-founded and extensional}).$$



The fact that $(\varpi_R "X)^{\mathcal{M}} \in I_{fix}(j)$ follows from the fact that j is an automorphism and ϖ_R is definable.

Lemma 3.9 For all $\kappa \in I_{fix}(j)$,

 $\mathcal{M} \models (\kappa \text{ is a cardinal}) \text{ if and only if } \mathcal{I}_{fix}(j) \models (\kappa \text{ is a cardinal}).$

Proof As in the proof of Lemma 3.1, this follows from the fact that ' κ is a cardinal' can be expressed as a $\Delta_0^{\mathcal{P}}$ -formula with parameters from $I_{\text{fix}}(j)$.

Lemma 3.10 $\mathcal{I}_{fix}(j) \models (every well-ordering is isomorphic to an ordinal).$

Proof Let $R, X \in I_{fix}(j)$ be such that $\mathcal{I}_{fix}(j) \models (R \text{ is a well-ordering of } X)$. It follows from Lemma 3.8 (and the fact that being a linear ordering is absolute) that $\mathcal{M} \models (R \text{ is a well-ordering of } X)$. Now, $(\varpi_R ``X)^{\mathcal{M}}$ is an ordinal and, by Lemma 3.8, $(\varpi_R ``X)^{\mathcal{M}} \in I_{fix}(j)$. Let $\alpha = (\varpi_R ``X)^{\mathcal{M}}$. In \mathcal{M} there is a bijection $f \subseteq X \times \alpha$ witnessing that R is isomorphic to α . Therefore $f \in I_{fix}(j)$ and f witnesses that R is isomorphic to α in $\mathcal{I}_{fix}(j)$.

Lemma 3.11 $\mathcal{I}_{fix}(j) \models (for \ all \ cardinals \ \kappa, \ the \ cardinal \ 2^{\kappa} \ exists).$

Proof Let $\kappa \in I_{fix}(j)$ be a cardinal. Work inside $\mathcal{I}_{fix}(j)$. Let $R \subseteq \mathcal{P}(\kappa) \times \mathcal{P}(\kappa)$ be a well-ordering. By Lemma 3.10, R is isomorphic to an ordinal $\alpha \geq 2^{\kappa}$.

Lemma 3.12 If $\kappa \in I_{fix}(j)$ is a cardinal then $H_{\kappa}^{\mathcal{M}} \in I_{fix}(j)$.

Proof Let $\kappa \in I_{fix}(j)$ be a cardinal. Work inside \mathcal{M} . Note that $|H_{\kappa}| = 2^{\kappa}$ and $2^{\kappa} \in I_{fix}(j)$. Using a bijection $f: H_{\kappa} \longrightarrow 2^{\kappa}$ one can find an $R \subseteq 2^{\kappa} \times 2^{\kappa}$ such that ϖ_R " $2^{\kappa} = H_{\kappa}$. Therefore, it follows from Lemma 3.8 that $H_{\kappa}^{\mathcal{M}} \in I_{fix}(j)$.

Lemma 3.13 $\mathcal{I}_{fix}(j) \models Axiom H$.

Proof By Lemma 3.10 it is enough to show that for every cardinal κ , there exists a transitive set which contains all transitive sets with cardinality $\leq \kappa$. Let $\kappa \in I_{\text{fix}}(j)$ be a cardinal. Lemmas 3.11 and 3.12 show that $2^{\kappa} \in I_{\text{fix}}(j)$ and $H_{2^{\kappa}}^{\mathcal{M}} \in I_{\text{fix}}(j)$. The point $H_{2^{\kappa}}^{\mathcal{M}} \in I_{\text{fix}}(j)$ is transitive and we claim that every transitive set of size $\leq \kappa$ in $I_{\text{fix}}(j)$ is contained in $H_{2^{\kappa}}^{\mathcal{M}}$. To see this, let $z \in I_{\text{fix}}(j)$ be such that $\mathcal{I}_{\text{fix}}(j) \models (\bigcup z \subseteq z) \land (|z| \leq \kappa)$ and so $\mathcal{I}_{\text{fix}}(j) \models (z \subseteq H_{2^{\kappa}}^{\mathcal{M}})$. This shows that $\mathcal{I}_{\text{fix}}(j)$ satisfies axiom H.

This completes the proof of Theorem 3.4.

4 Realizing *H*-cuts as $\mathcal{I}_{fix}(j)$

In this section we will show if I is an H-cut of a countable model \mathcal{M} of MOST + $\Delta_0^{\mathcal{P}}$ -collection then there is $\mathcal{M} \subseteq_{\mathrm{cf}} \mathcal{N}$ with $\mathcal{N} \models \mathrm{MOST}$ and a non-trivial automorphism $j: \mathcal{N} \longrightarrow \mathcal{N}$ such that $\mathcal{I}_{\mathrm{fix}}(j) = \langle I, \in^{\mathcal{M}} \rangle$ and $\mathrm{fix}(j) = M$. More generally, we will prove the following result:



Theorem 4.1 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be countable such that $\mathcal{M} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection and let $I \subseteq M$ be an H-cut of \mathcal{M} . If \mathbb{L} is a linear order then there is $\mathcal{M} \subseteq_{\text{cf}} \mathcal{N}$ with $\mathcal{N} \models \text{MOST}$ and an embedding $j \mapsto \check{j}$ of $\text{Aut}(\mathbb{L})$ into $\text{Aut}(\mathcal{N})$ such that if j has no fixed points then $\mathcal{I}_{\text{fix}}(\check{j}) = \langle I, \in^{\mathcal{M}} \rangle$ and $\text{fix}(\check{j}) = M$.

By setting \mathbb{L} in Theorem 4.1 to be the linear ordering \mathbb{Z} and by letting $j: \mathbb{Z} \longrightarrow \mathbb{Z}$ be the automorphism that sends $i \mapsto i+1$ for all $i \in \mathbb{Z}$, we obtain:

Corollary 4.2 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be countable such that $\mathcal{M} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection and let $I \subseteq M$ be an H-cut of \mathcal{M} . Then there is $\mathcal{M} \subseteq_{\text{cf}} \mathcal{N}$ with $\mathcal{N} \models \text{MOST}$ and an automorphism $j : \mathcal{N} \longrightarrow \mathcal{N}$ such that $\mathcal{I}_{\text{fix}}(j) = \langle I, \in^{\mathcal{M}} \rangle$ and fix(j) = M.

Theorem 4.1 will be proved by adapting a construction, originally due to Jeff Paris and George Mills [19] and applied to non-standard models of PA, to non-standard models of MOST + $\Delta_0^{\mathcal{P}}$ -collection. This construction was also used by the first author in [4] to prove that every cut of a countable model of $I\Delta_0 + B\Sigma_1$ that is closed under exponentiation can be realised as the largest initial segment of a model of $I\Delta_0$ that is pointwise fixed by some non-trivial automorphism. Theorem 4.1 and Corollary 4.2 are the set-theoretic analogues of results proved in section 3 of [4].

Throughout this section we fix an \mathcal{L} -structure $\mathcal{M} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection with $|\mathcal{M}| = \aleph_0$, and $I \subseteq \mathcal{M}$ an H-cut of \mathcal{M} . We also fix a linear order \mathbb{L} . For simplicity we will use < for the order relation on \mathbb{L} . We will write \mathcal{I} for the substructure of \mathcal{M} with domain I. We also fix $\bar{k} \in \text{Ord}^{\mathcal{M}} \setminus \text{Ord}^{\mathcal{I}}$ such that $\mathcal{M} \models (\bar{k} \text{ is a regular cardinal})$. We begin by noting that there is no least cardinal in $\text{Ord}^{\mathcal{M}} \setminus \text{Ord}^{\mathcal{I}}$, and so, regardless of our choice of \bar{k} , there are infinitely many cardinals in $\text{Ord}^{\mathcal{M}} \setminus \text{Ord}^{\mathcal{I}}$ below \bar{k} .

Lemma 4.3 There is no least cardinal in the class $Ord^{\mathcal{M}} \setminus Ord^{\mathcal{I}}$.

Proof Suppose that $\kappa \in \text{Ord}^{\mathcal{M}} \setminus \text{Ord}^{\mathcal{I}}$ is least such that $\mathcal{M} \models (\kappa \text{ is a cardinal})$. Working inside \mathcal{M} , consider the set

$$C = \{\alpha \in \kappa \mid |\alpha| < \kappa\}.$$

Note that $C \in M$. Since there is no largest cardinal in \mathcal{I} , $C^* = \operatorname{Ord}^{\mathcal{I}} \subseteq I$. This shows that $C \notin I$, which contradicts the fact that I is an H-cut of \mathcal{M} .

Using an iterated ultrapower construction we will obtain a cofinal extension $\mathcal N$ of $\mathcal M$ that satisfies MOST and such that any fixed point free automorphism of $\mathbb L$ generates an automorphism $j: \mathcal N \longrightarrow \mathcal N$ with $\mathcal I_{\mathrm{fix}}(j) = \mathcal I$ and $\mathrm{fix}(j) = M$. The ultrafilter $\mathcal U$ used in this iterated ultrapower construction will be an ultrafilter on the subsets of $\bar \kappa$ in $\mathcal M$. The elements of $\mathcal U$ will be large in the following sense:

Definition 4.1 We say that $X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-large if $|X|^{\mathcal{M}} \notin I$.

If X is a set and λ is a cardinal then we use $[X]^{\lambda}$ to denote the set of all subsets of X of size λ . Recall the canonical partion relation $\kappa \to *(\lambda)^n$, first considered in [5], that generalises the classical partition relation $\kappa \to (\lambda)^n_u$:



Definition 4.2 Let κ and λ be cardinals and let $n \in \omega$. If f is a function with domain $[\kappa]^n$, $H \subseteq \kappa$, and $\Delta \subseteq n$ such that for all $s_0 < \cdots < s_{n-1}$ and $t_0 < \cdots < t_{n-1}$ in H we have

$$f(\{s_0, \ldots, s_{n-1}\}) = f(\{t_0, \ldots, t_{n-1}\})$$
 if and only if $(\forall i \in \Delta)(s_i = t_i)$,

then we say that H is f-canonical. We write $\kappa \to *(\lambda)^n$ if for all functions f with domain $[\kappa]^n$, there is $H \subseteq \kappa$ with $|H| = \lambda$ and H is f-canonical.

This notion allows us to make explicit the combinatorial properties we will require of the ultrafilter used to extend \mathcal{M} .

Definition 4.3 We say that a non-principal (n.p.) ultrafilter $\mathcal{U} \subseteq (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-complete if for all $f \in M$, if $\mathcal{M} \models (f : [\bar{\kappa}]^n \longrightarrow TC(A))$ where $n \in \omega$ and $A \in I$ then there is $X \in \mathcal{U}$ such that $\mathcal{M} \models (f \text{ is constant on } [X]^n)$.

Definition 4.4 We say that an n.p. ultrafilter $\mathcal{U} \subseteq (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is canonically Ramsey if for all $f \in \mathcal{M}$, if $\mathcal{M} \models (f \text{ is a function with domain } [\bar{\kappa}]^n)$ where $n \in \omega$ then there is $X \in \mathcal{U}$ such that $\mathcal{M} \models (X \text{ is } f\text{-canonical}).$

Definition 4.5 We say that an n.p. ultrafilter $\mathcal{U} \subseteq (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-tight if for all $f \in M$, if $\mathcal{M} \models (f \text{ is a function with domain } [\bar{\kappa}]^n)$ where $n \in \omega$ then there is $X \in \mathcal{U}$ such that either

- (a) $\mathcal{M} \models (f \text{ is constant on } [X]^n), \text{ or }$
- (b) there is an I-large cardinal θ such that

$$\mathcal{M} \models \forall x_1 \cdots \forall x_n \left(\bigwedge_{1 \leq i \leq n} (x_i \in X) \Rightarrow f(\{x_1, \dots, x_n\}) \notin H_\theta \right).$$

It is important to note that Definitions 4.3, 4.4 and 4.5 only refer to functions in \mathcal{M} whose domain is $[\bar{\kappa}]^n$ where n is a natural number of the meta-theory. We now turn to showing that $\mathcal{P}(\bar{\kappa})$ in \mathcal{M} is rich enough to ensure the existence of an (external) ultrafilter on $(\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ that is simultaneously I-complete, canonically Ramsey, I-tight and contains arbitrarily small I-large sets.

Lemma 4.4 If $X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-large, $\lambda \in I$ is such that $\mathcal{M} \models (\lambda \text{ is a cardinal})$ and $f \in M$ is such that $\mathcal{M} \models f : \bar{\kappa} \longrightarrow H_{\lambda}$, then there is $Y \in (\mathcal{P}(X)^{\mathcal{M}})^*$ such that Y is I-large and f is constant on Y.

Proof Work inside \mathcal{M} . Let $X \subseteq \bar{\kappa}$ be I-large and let λ be a cardinal in I. Let $f: \bar{\kappa} \longrightarrow H_{\lambda}$. The fact that I is an H-cut of \mathcal{M} ensures that H_{λ} is in I and $|H_{\lambda}|$ is not I-large. Let μ be an I-large regular cardinal such that $\mu < |X|$. So, $|H_{\lambda}| < \mu$. Suppose that for all $x \in H_{\lambda}$, $f^{-1}[x] \cap X$ is not I-large, in particular $|f^{-1}[x] \cap X| < \mu$. So,

$$|X| = \left| \bigcup_{x \in H_{\lambda}} f^{-1}[x] \cap X \right| \le \mu,$$

which is a contradiction.



In [1, Corollary 2] Baumgartner completely determines the canonical partition relation on infinite cardinals and reveals the following relationship between the canonical and classical partition relations:

Theorem 4.5 (Baumgartner) If κ and λ are infinite cardinals and $n \in \omega$ then

$$\kappa \to *(\lambda)^{n+1}$$
 if and only if for all $\mu < \lambda, \kappa \to (\lambda)_{\mu}^{n+1}$. (1)

The setting for Baumgartner's [1] is ZFC, however an examination of the proof of Theorem 4.5 reveals that for fixed κ and λ all the instances of separation and replacement appealed to in the proof of the equivalence described by (1) can be restricted to sets that provably exist in MOST. This means that for fixed κ and λ , the equivalence described by (1) is provable in MOST. In ZFC, Theorem 4.5 coupled with the Erdős–Rado Theorem [6] shows that for every infinite successor cardinal λ and $n \in \omega$, there exists a cardinal κ such that $\kappa \to *(\lambda)^{n+1}$. Specifically:

Theorem 4.6 (Erdős–Rado) If κ is an infinite cardinal and $n \in \omega$ then

$$\beth_n(\kappa)^+ \to (\kappa^+)^{n+1}_{\kappa}$$
.

If both κ^+ and $\beth_n(\kappa)^+$ exist then all instances of separation and replacement appealed to in the proof that $\beth_n(\kappa)^+ \to (\kappa^+)^{n+1}_{\kappa}$ can be bounded by sets that provably exist in MOST. Thus, as long as both κ^+ and $\beth_n(\kappa)^+$ exist, MOST proves that $\beth_n(\kappa)^+ \to (\kappa^+)^{n+1}_{\kappa}$. It should be noted, however, that MOST is incapable of proving that $\beth_n(\aleph_0)$ exists for all natural numbers n. Despite this, for any infinite cardinal κ , MOST proves that 2^{κ} exists. Therefore, if n is (externally) a standard natural number and κ is an infinite cardinal then MOST proves that $\beth_n(\kappa)^+$ exists. In the context of MOST Theorem 4.6 becomes:

Lemma 4.7 MOST proves the theorem scheme: for all $n \in \omega$,

$$\forall \kappa \left((\kappa \text{ is an infinite cardinal}) \Rightarrow \left(\exists_n (\kappa)^+ \rightarrow (\kappa^+)_{\kappa}^{n+1} \right) \right).$$

Therefore combining Theorems 4.5 and 4.6 in the context of the theory MOST we get:

Lemma 4.8 MOST proves the theorem scheme: for all $n \in \omega$,

$$\forall \kappa \left((\kappa \text{ is an infinite cardinal}) \Rightarrow \left(\beth_n(\kappa)^+ \rightarrow *(\kappa^+)^{n+1} \right) \right).$$

Lemma 4.9 Let $n \in \omega$. If $\lambda \in \bar{\kappa}^* \cup \{\bar{\kappa}\}$ is *I*-large then there exists an *I*-large $\mu \in \bar{\kappa}^*$ with $\mathcal{M} \models (\mu \text{ is a cardinal}) \text{ such that } \beth_n(\mu) < \lambda$.



Proof Work inside \mathcal{M} . Assume that $\lambda \leq \bar{\kappa}$ is *I*-large and for all *I*-large cardinals γ , $\beth_n(\gamma) \geq \lambda$. Consider

$$C = \{ \gamma \in \bar{\kappa} \mid (\gamma \text{ is a cardinal}) \land (\beth_n(\gamma) < \lambda) \}.$$

By bounding all the quantifiers in the defining formula of C by $H_{\bar{\kappa}^+}$ we can see that Δ_0 -separation implies that C is a set. Our assumption about λ implies that $C^* \subseteq I$. Therefore, since I is a proper H-cut, $C \in I$. But this means that C is bounded in I. But if $\gamma \in I$ is a cardinal with $\gamma \notin C^*$ then $\beth_n(\gamma) \notin I$, which contradicts the fact that $\mathcal{I} \models \text{MOST}$.

Lemma 4.10 Let $n \in \omega$. If $X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-large and $f \in M$ is such that $\mathcal{M} \models (f \text{ is a function with domain } [X]^{n+1})$ then there is an I-large $Y \in (\mathcal{P}(X)^{\mathcal{M}})^*$ that is f-canonical.

Proof Work inside \mathcal{M} . Let $X \subseteq \bar{\kappa}$ be *I*-large and let f be a function with domain $[X]^{n+1}$. Let $\lambda = |X|$. So, λ is *I*-large. Using Lemma 4.9 we can find an *I*-large cardinal μ such that $\beth_n(\mu) < \lambda$. Therefore, by Lemma 4.8, $\lambda \to *(\mu^+)^{n+1}$. And so there is an *I*-large $Y \subseteq X$ that is f-canonical.

Lemma 4.11 Let $n \in \omega$. If $X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ is I-large and $f \in M$ is such that $\mathcal{M} \models (f \text{ is a function with domain } [X]^{n+1})$ then there is an I-large $Y \in (\mathcal{P}(X)^{\mathcal{M}})^*$ such that either

- (a) $\mathcal{M} \models (f \text{ is constant on } [Y]^{n+1}), \text{ or }$
- (b) there is an I-large $\theta \in \bar{\kappa}^*$ with $\mathcal{M} \models (\theta \text{ is a cardinal})$ such that

$$\mathcal{M} \models (\forall A \in [Y]^{n+1})(f(A) \notin H_{\theta}).$$

Proof Work inside \mathcal{M} . Let $X \subseteq \bar{\kappa}$ be I-large and let f be a function with domain $[X]^{n+1}$. Let $\kappa = |X|$. Using Lemma 4.9 we can find an I-large cardinal λ such that $\kappa \to (\lambda)_2^{n+1}$. Applying Lemma 4.9 we can then obtain an I-large cardinal μ^+ such that $\lambda \to (\mu^+)_{\mu}^{n+1}$. Then, using Lemma 4.9 again, let θ be an I-large cardinal such that $2^{\theta} < \mu^+$. Therefore $|H_{\theta}| \le \mu$. Now, define $g: [X]^{n+1} \longrightarrow 2$ such that for all $\{x_0, \ldots, x_n\} \in [X]^{n+1}$,

$$g(\lbrace x_0, \dots, x_n \rbrace) = \begin{cases} 0 & \text{if } f(\lbrace x_0, \dots, x_n \rbrace) \in H_\theta \\ 1 & \text{otherwise} \end{cases}$$

Let $Z \subseteq X$ be such that $|Z| \ge \lambda$ and g is constant on $[Z]^{n+1}$. If $g''[Z]^{n+1} = \{1\}$ then we are done. If $g''[Z]^{n+1} = \{0\}$ then let f' be the restriction of f to $[Z]^{n+1}$. Since $|\operatorname{rng}(f')| \le |H_{\theta}| \le \mu$, there is a $Y \subseteq Z$ with $|Y| \ge \mu^+$ such that f' is constant on $[Y]^{n+1}$.

Equipped with Lemmas 4.4, 4.10 and 4.11 we now show that there exists an external non-principle ultrafilter on the subsets of $\bar{\kappa}$ in \mathcal{M} which is simultaneously *I*-complete, canonically Ramsey, *I*-tight and contains arbitrarily small *I*-large sets. We use \mathcal{F} to



denote the class of all points in \mathcal{M} which correspond to a function whose domain is $[\bar{\kappa}]^n$ for some external natural number n. I.e.

$$\mathcal{F} = \{ f \in M \mid (\exists n \in \omega) (\mathcal{M} \models (f \text{ is a function with domain } [\bar{k}]^n)) \}.$$

Theorem 4.12 There exists an n.p. ultrafilter $\mathcal{U} \subseteq (\mathcal{P}(\bar{k})^{\mathcal{M}})^*$ which is I-complete, canonically Ramsey, I-tight and is such that $\{|X|^{\overline{\mathcal{M}}} \mid X \in \mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \setminus \mathrm{Ord}^{\mathcal{I}}$.

Proof Let $\langle f_n \mid n \in \omega \rangle$ be an enumeration of \mathcal{F} . Let $\langle k_n \mid n \in \omega \rangle$ be a sequence of natural numbers such that for all $n \in \omega$, $\mathcal{M} \models (f_n \text{ is a function with domain } [\bar{\kappa}]^{k_n})$. Let $\langle \lambda_n \mid n \in \omega \rangle$ be a decreasing sequence of \mathcal{M} -cardinals that is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \backslash \mathrm{Ord}^{\mathcal{I}}$ with $\lambda_0 \leq \bar{\kappa}$. Using Lemmas 4.4, 4.10 and 4.11 inductively build sequences $\langle W_n \mid n \in \omega \rangle$, $\langle X_n \mid n \in \omega \rangle$, $\langle Y_n \mid n \in \omega \rangle$ and $\langle Z_n \mid n \in \omega \rangle$ of I-large elements of $(\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ such that for all $n \in \omega$,

- 1. $\mathcal{M} \models (W_n \supseteq X_n \supseteq Y_n \supseteq Z_n \supseteq W_{n+1}),$
- 2. $\mathcal{M} \models (W_n \text{ is } f_n\text{-canonical}),$
- 3. if $k_n = 1$ and there is an \mathcal{M} -cardinal $\mu \in I$ such that $\mathcal{M} \models (f_n : \bar{k} \longrightarrow H_{\mu})$ then f_n is constant on X_n , otherwise $X_n = W_n$,
- 4. $\mathcal{M} \models (f_n \text{ is constant on } [Y_n]^{k_n})$ or there is an I-large \mathcal{M} -cardinal $\mu \in \bar{\kappa}^*$ such that

$$\mathcal{M} \models \left(\forall A \in [Y_n]^{k_n} \right) (f_n(A) \notin H_\mu),$$

5. $\mathcal{M} \models (|Z_n| < \lambda_n)$.

Define $\mathcal{U} = \{X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^* \mid (\exists n \in \omega)(\mathcal{M} \models W_n \subseteq X)\}$. It is clear from the construction that \mathcal{U} is an n.p. ultrafilter that is I-complete, canonically Ramsey, I-tight and is such that $\{|X|^{\mathcal{M}} \mid X \in \mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \setminus \mathrm{Ord}^{\mathcal{I}}$. \square

Let $\mathcal{U} \subseteq (\mathcal{P}(\bar{k})^{\mathcal{M}})^*$ be an n.p. ultrafilter obtained from Theorem 4.12, so \mathcal{U} is I-complete, canonically Ramsey, I-tight and $\{|X|^{\mathcal{M}} \mid X \in \mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \setminus \mathrm{Ord}^{\mathcal{I}}$. Let $\mathcal{L}_{\mathcal{F}}$ be the language extending \mathcal{L} such that for all $f \in \mathcal{F}$ and $n \in \omega$, if $\mathcal{M} \models (f \text{ has domain } [\bar{k}]^n)$ then $\mathcal{L}_{\mathcal{F}}$ has a new n-ary function symbol \hat{f} . Let $\mathcal{M}_{\mathcal{F}}$ be the $\mathcal{L}_{\mathcal{F}}$ -structure obtained from \mathcal{M} by defining, for all $f \in \mathcal{F}$ with $\mathcal{M} \models (f \text{ has domain } [\bar{k}]^n)$,

$$\hat{f}^{\mathcal{M}_{\mathcal{F}}}(x_1, \dots, x_n) = \begin{cases} f(\{x_1, \dots, x_n\})^{\mathcal{M}} & \text{if } x_1 < \dots < x_n \in \bar{\kappa}^* \\ \emptyset & \text{otherwise} \end{cases}$$

For each $n \in \omega$, define the n + 1-partial type $\Gamma_n(x_0, \dots, x_n) \subseteq \mathcal{L}_{\mathcal{F}}$ by

$$\phi(x_0, \dots, x_n) \in \Gamma_n(x_0, \dots, x_n)$$
 if and only if $\exists X \in \mathcal{U}$ such that $\mathcal{M}_{\mathcal{F}} \models \phi(a_0, \dots, a_n)$ for all $a_0 < \dots < a_n \in X^*$

Define

$$T_{\mathcal{U}} = \bigcup_{n \in \omega} \Gamma_n(x_0, \dots, x_n).$$



Let $\mathcal{L}_{\mathcal{F},\mathbb{L}}$ be the extension of $\mathcal{L}_{\mathcal{F}}$ obtained by adding constant symbols c_i for each $i \in \mathbb{L}$. Define

$$T_{\mathcal{U},\mathbb{L}} = \{\phi(c_{i_0},\ldots,c_{i_n}) \mid \phi \in T_{\mathcal{U}} \text{ and } i_0 < \cdots < i_n \in \mathbb{L}\}.$$

The fact that for each $f \in \mathcal{F}$, the interpretation of the function symbol \hat{f} in $\mathcal{M}_{\mathcal{F}}$ is coded by the point $f \in M$ yields the following:

Lemma 4.13
$$\mathcal{M}_{\mathcal{F}} \models \Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$$
-separation $+ \Sigma_1^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -collection.

Lemma 4.14 $T_{\mathcal{U},\mathbb{L}}$ is consistent and is $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F},\mathbb{L}})$ -complete.

Proof The fact that $T_{\mathcal{U},\mathbb{L}}$ is finitely realisable implies that it is consistent. Let $\phi(x_0,\ldots,x_n)$ be a $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -formula. We work inside $\mathcal{M}_{\mathcal{F}}$. $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -separation ensures that the function $f:[\bar{\kappa}]^{n+1}\longrightarrow 2$ defined such that for all $\alpha_0<\cdots<\alpha_n\in\bar{\kappa}$,

$$f(\alpha_0, \dots, \alpha_n) = \begin{cases} 1 & \text{if } \phi(\alpha_0, \dots, \alpha_n) \\ 0 & \text{otherwise} \end{cases}$$

is a set. Now, there is an $X \in \mathcal{U}$ such that

$$\mathcal{M} \models (f \text{ is constant on } [X]^{n+1}).$$

Therefore, for all
$$i_0 < \cdots < i_n \in \mathbb{L}$$
, $T_{\mathcal{U}, \mathbb{L}}$ decides $\phi(c_{i_0}, \ldots, c_{i_n})$.

Remark 4.15 Lemma 4.14 generalises to show that for all $n \in \omega$, if $\mathcal{M} \models \Sigma_n$ -separation then $T_{\mathcal{U},\mathbb{L}}$ is $\Sigma_n(\mathcal{L}_{\mathcal{F},\mathbb{L}})$ -complete. In particular, if $\mathcal{M} \models \mathsf{ZFC}$ then $T_{\mathcal{U},\mathbb{L}}$ is $\mathcal{L}_{\mathcal{F},\mathbb{L}}$ -complete.

Let TERM $\subseteq \mathcal{L}_{\mathcal{F},\mathbb{L}}$ be the class of terms of the form $\hat{f}(c_{i_0},\ldots,c_{i_n})$ where $i_0<\cdots< i_n\in\mathbb{L}$. Define

$$\hat{f}(c_{i_0}, \dots, c_{i_n}) \sim \hat{g}(c_{j_0}, \dots, c_{j_m})$$
 if and only if $(\hat{f}(c_{i_0}, \dots, c_{i_n}))$
= $\hat{g}(c_{j_0}, \dots, c_{j_m})) \in T_{\mathcal{U}, \mathbb{L}}$.

Now, \sim is an equivalence class on TERM. If $\tau \in$ TERM then we use $[\tau]$ to denote the equivalence class to which τ belongs. We now turn to defining an \mathcal{L} -structure $\mathcal{N}_{\mathcal{U},\mathbb{L}} = \langle \mathcal{N}_{\mathcal{U},\mathbb{L}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{L}}} \rangle$. Let

$$N_{\mathcal{U},\mathbb{L}} = \{ [\tau] \mid \tau \in \text{TERM} \}.$$

Define

$$\left[\hat{f}\left(c_{i_0},\ldots,c_{i_n}\right)\right] \in \mathcal{N}_{\mathcal{U},\mathbb{L}}\left[\hat{g}\left(c_{j_0},\ldots,c_{j_m}\right)\right] \text{ if and only if}$$

$$\left(\hat{f}\left(c_{i_0},\ldots,c_{i_n}\right) \in \hat{g}\left(c_{j_0},\ldots,c_{j_m}\right)\right) \in T_{\mathcal{U},\mathbb{L}}.$$



Lemma 4.16 Let $\phi(x_0, \ldots, x_n)$ be a $\Delta_0(\mathcal{L}_{\mathcal{F}})$ -formula. If

$$\mathcal{M}_{\mathcal{F}} \models (\forall x_1, \dots x_n \in \bar{\kappa})((x_1 < \dots < x_n) \Rightarrow \exists y \phi(y, x_1, \dots, x_n))$$

then there exists $f \in \mathcal{F}$ such that

$$\mathcal{M}_{\mathcal{F}} \models (\forall x_1, \dots, x_n \in \bar{\kappa})((x_1 < \dots < x_n) \Rightarrow \phi \left(\hat{f}(x_1, \dots, x_n), x_1, \dots, x_n)\right).$$

Proof Assume that $\mathcal{M}_{\mathcal{F}} \models (\forall x_1, \dots x_n \in \bar{\kappa})((x_1 < \dots < x_n) \Rightarrow \exists y \phi(y, x_1, \dots, x_n))$. We work inside $\mathcal{M}_{\mathcal{F}}$. Consider $[\bar{\kappa}]^n$ and define

$$\phi'(y, x)$$
 if and only if $(\exists x_1, \dots, x_n \in x)((x_1 < \dots < x_n) \land \phi(y, x_1, \dots, x_n))$.

Note that $(\forall x \in [\bar{\kappa}]^n) \exists y \phi'(y, x)$. Therefore $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -collection ensures that there is a set A such that $(\forall x \in [\bar{\kappa}]^n)(\exists y \in A)\phi'(y, x)$. Let $\blacktriangleleft \subseteq A \times A$ be a well-ordering of A. Let $\psi(y, x)$ be the $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -formula

$$\phi'(y, x) \land (\forall z \in A)(\phi'(z, x) \Rightarrow \langle z, y \rangle \notin \blacktriangleleft).$$

Now, $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -separation ensures that $f:[\bar{\kappa}]^n \longrightarrow A$, defined such that f(x) is the unique y such that $\psi(y,x)$, is a set.

Lemma 4.17 If $\phi(x_0, \ldots, x_n)$ a $\Sigma_1^{\mathcal{P}}$ -formula then for all $[\tau_0], \ldots, [\tau_n] \in N_{\mathcal{U}, \mathbb{L}}$,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \phi([\tau_0],\ldots,[\tau_n])$$
 if and only if $\phi(\tau_0,\ldots,\tau_n) \in T_{\mathcal{U},\mathbb{L}}$.

Proof Let $\phi(x_0, \ldots, x_n)$ be a $\Sigma_1^{\mathcal{P}}$ -formula and let $[\tau_0], \ldots, [\tau_n] \in \mathcal{N}_{\mathcal{U}, \mathbb{L}}$. Without loss of generality we can assume that ϕ only uses the quantifier \exists and the logical connectives \land and \neg . We prove the result by induction on the structure of ϕ . It is clear that the result holds for all atomic formulae and conjunctions of atomic formulae. Assume that the result holds for a $\Delta_0^{\mathcal{P}}$ -formula $\theta(x_0, \ldots, x_n)$. Lemma 4.14 implies that the result also holds for $\neg \theta$.

We now turn to showing that the class of formula satisfying the result is closed under existential quantification. Assume that the result holds for a $\Delta_0^{\mathcal{P}}$ -formula $\theta(y, x_0, \ldots, x_n)$. Let $[\tau_0], \ldots, [\tau_n] \in \mathcal{N}_{\mathcal{U}, \mathbb{L}}$. Suppose that $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models \exists y \theta(y, [\tau_0], \ldots, [\tau_n])$. Let $[\tau] \in \mathcal{N}_{\mathcal{U}, \mathbb{L}}$ be such that $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models \theta([\tau], [\tau_0], \ldots, [\tau_n])$. Therefore $\theta(\tau, \tau_0, \ldots, \tau_n) \in T_{\mathcal{U}, \mathbb{L}}$. This implies that $\exists y \theta(y, \tau_0, \ldots, \tau_n) \in T_{\mathcal{U}, \mathbb{L}}$.

Conversely, suppose that $\exists y \theta(y, \tau_0, \dots, \tau_n) \in T_{\mathcal{U}, \mathbb{L}}$. Assume, without loss of generality, that for all $0 \leq j \leq n$, $\tau_j = \hat{f}_j(c_{i_0}, \dots, c_{i_m})$ for $i_0 < \dots < i_m$ in \mathbb{L} . Let $X \in \mathcal{U}$ be such that for all $a_0 < \dots < a_m \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \exists y \theta \left(y, \, \hat{f}_0(a_0, \ldots, a_m), \ldots, \, \hat{f}_n(a_0, \ldots, a_m) \right).$$



Let $\theta'(y, x_0, ..., x_m)$ be the $\Delta_0^{\mathcal{P}}(\mathcal{L}_{\mathcal{F}})$ -formula

$$\left(\bigwedge_{0 \le i \le m} (x_i \in X)\right) \land (x_0 < \dots < x_m)$$

$$\Rightarrow \theta \left(y, \hat{f}_0(x_0, \dots, x_m), \dots, \hat{f}_n(x_0, \dots, x_m)\right).$$

Using Lemma 4.16 we can find an $f \in \mathcal{F}$ such that for all $x_0 < \cdots < x_m \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \theta'(\hat{f}(x_0,\ldots,x_m),x_0,\ldots,x_m).$$

Let
$$\tau = \hat{f}(c_{i_0}, \dots, c_{i_m})$$
. Now, $\theta(\tau, \tau_0, \dots, \tau_n) \in T_{\mathcal{U}, \mathbb{L}}$. Therefore $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models \theta([\tau], [\tau_0], \dots, [\tau_n])$ and so $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models \exists y \theta(y, [\tau_0], \dots, [\tau_n])$.

Remark 4.18 The proof of Lemma 4.17 generalises to show that for all $n \in \omega$, if $\mathcal{M} \models \Sigma_n$ -collection then the conclusion of Lemma 4.17 holds for all Σ_n -formulae. In particular, if $\mathcal{M} \models \text{ZFC}$ then $\mathcal{M} \prec \mathcal{N}_{\mathcal{U}, \mathbb{L}}$.

The structure \mathcal{M} embeds into $\mathcal{N}_{\mathcal{U},\mathbb{L}}$. To see this observe that for all $m \in M$ there exists an $h_m \in \mathcal{F}$ such that

$$\mathcal{M}_{\mathcal{F}} \models (\forall x \in \bar{\kappa}) \left(\hat{h}_m(x) = m \right).$$

It is clear that for all $i, j \in \mathbb{L}$ and for all $m \in M$, $[\hat{h}_m(c_i)]$ is equal to $[\hat{h}_m(c_j)]$. In order to make it easier to refer to these elements of $N_{\mathcal{U}, \mathbb{L}}$ we will fix an element $0 \in \mathbb{L}$ so that we can write $[\hat{h}_m(c_0)]$. We can also see that the linear order \mathbb{L} embeds in the ordinals of $\mathcal{N}_{\mathcal{U}, \mathbb{L}}$. To see this consider the map $\mathrm{id} \in \mathcal{F}$ such that

$$\mathcal{M}_{\mathcal{F}} \models (\hat{\mathrm{id}} : \bar{\kappa} \longrightarrow \bar{\kappa}) \land (\forall x \in \bar{\kappa}) (\hat{\mathrm{id}}(x) = x).$$

For each $i \in \mathbb{L}$, the term $[\hat{id}(c_i)]$ is an ordinal in $\mathcal{N}_{\mathcal{U},\mathbb{L}}$.

Lemma 4.19 The class $\{[\hat{id}(c_i)] \mid i \in \mathbb{L}\} \subseteq N_{\mathcal{U},\mathbb{L}}$ is a class of $\Sigma_1^{\mathcal{P}}$ -indiscernibles.

Proof This follows immediately from Lemma 4.17.

The fact that the image of the embedding of \mathbb{L} into $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ is a class of $\Sigma_1^{\mathcal{P}}$ -indiscernibles means that any automorphism of \mathbb{L} raises to an automorphism of $\mathcal{N}_{\mathcal{U},\mathbb{L}}$. For all $j \in \operatorname{Aut}(\mathbb{L})$, define $\check{j} : \mathcal{N}_{\mathcal{U},\mathbb{L}} \longrightarrow \mathcal{N}_{\mathcal{U},\mathbb{L}}$ by

$$\check{j}\left(\left[\hat{f}\left(c_{i_1}, \dots c_{i_n}\right)\right]\right) = \left[\hat{f}\left(c_{j(i_1)}, \dots c_{j(i_n)}\right)\right] \text{ for all } f \in \mathcal{F} \text{ and } i_1 < \dots < i_n \in \mathbb{L}.$$

Lemma 4.19 implies that for all $j \in \operatorname{Aut}(\mathbb{L})$, $\check{j} : \mathcal{N}_{\mathcal{U},\mathbb{L}} \longrightarrow \mathcal{N}_{\mathcal{U},\mathbb{L}}$ is an automorphism. Moreover, the map from $\operatorname{Aut}(\mathbb{L})$ into $\operatorname{Aut}(\mathcal{N}_{\mathcal{U},\mathbb{L}})$ defined by $j \mapsto \check{j}$ for all $j \in \operatorname{Aut}(\mathbb{L})$,



is an injective group homomorphism. It is also immediate from (2) that for all $j \in Aut(\mathbb{L})$ and for all $x \in M$, \check{j} fixes $[\hat{h}_x(c_0)]$, so fix $(\check{j}) \supseteq M$.

Lemma 4.17 also implies that embedding of \mathcal{M} into $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ defined by $x \mapsto [\hat{h}_x(c_0)]$ preserves $\Pi_2^{\mathcal{P}}$ properties of tuples from \mathcal{M} .

Lemma 4.20 Let $\phi(x_0, ..., x_n)$ be a $\Pi_2^{\mathcal{P}}$ -formula. For all $a_0, ..., a_n \in M$, if $\mathcal{M} \models \phi(a_0, ..., a_n)$ then

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \phi\left(\left[\hat{h}_{a_0}(c_0)\right],\ldots,\left[\hat{h}_{a_n}(c_0)\right]\right).$$

Proof The formula $\phi(x_0, \dots, x_n)$ is in the form $\forall y_0 \dots \forall y_m \theta(y_0, \dots, y_m, x_0, \dots, x_n)$ where $\theta(y_0, \dots, y_m, x_0, \dots, x_n)$ is $\Sigma_1^{\mathcal{P}}$. Let $a_0, \dots, a_n \in M$. Assume that

$$\mathcal{M} \models \forall y_0 \cdots \forall y_m \theta(y_0, \dots, y_m, a_0, \dots, a_n).$$

Let $f_0, \ldots, f_m \in \mathcal{F}$. Without loss of generality assume that each $\hat{f}_0, \ldots, \hat{f}_m$ are k+1-ary function symbols in $\mathcal{L}_{\mathcal{F}}$. For all $x_0 < \cdots < x_k \in \bar{\kappa}^*$,

$$\mathcal{M}_{\mathcal{F}} \models \theta \left(\hat{f}_0(x_0, \dots, x_k), \dots, \hat{f}_m(x_0, \dots, x_k), a_0, \dots, a_n \right).$$

Therefore, for all $i_0 < \cdots < i_k \in \mathbb{L}$,

$$\theta\left(\hat{f}_0\left(c_{i_0},\ldots,c_{i_k}\right),\ldots,\hat{f}_m\left(c_{i_0},\ldots,c_{i_k}\right),\hat{h}_{a_0}(c_0),\ldots,\hat{h}_{a_n}(c_0)\right)\in T_{\mathcal{U},\mathbb{L}}.$$

And so, by Lemma 4.17,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \theta\left(\left[\hat{f}_0\left(c_{i_0},\ldots,c_{i_k}\right)\right],\ldots,\left[\hat{f}_m\left(c_{i_0},\ldots,c_{i_k}\right)\right],\left[\hat{h}_{a_0}(c_0)\right],\ldots,\left[\hat{h}_{a_n}(c_0)\right]\right).$$

Since $f_0, \ldots, f_m \in \mathcal{F}$ and $i_0 < \cdots < i_k \in \mathbb{L}$ were arbitrary, the Lemma follows. \square

Using this result we can show that $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ satisfies all of the axioms of Mac.

Lemma 4.21 $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \text{Mac.}$

Proof By Lemma 4.20 it is enough to show that every axiom of Mac can be written as a $\Pi_2^{\mathcal{P}}$ -sentence.

Extensionality: $\forall x \forall y (x = y \iff (\forall z \in x)(z \in y) \land (\forall z \in y)(z \in x))$

Emptyset: $\exists x (\forall y \in x) (y \neq y)$

Union: $\forall x \exists y (\forall z \in x) (\forall w \in z) (w \in y)$

Pairing: $\forall x \forall y \exists z ((x \in z) \land (y \in z) \land (\forall w \in z)(w = x \lor w = y))$

 Δ_0 -separation: for all Δ_0 -formulae $\phi(x, \vec{z})$,

$$\forall a \forall \vec{z} (\exists y \subseteq a) (\forall x \in a) (x \in y \iff \phi(x, \vec{z})).$$

Set Foundation: $\forall x ((\forall w \in x)(w \neq w) \lor (\exists y \in x)(\forall z \in y)(z \notin x))$



Infinity:

$$\exists S(\emptyset \in S \land (\forall x \in S)(\exists y \in S)(x \in y \land (\forall z \in x)(z \in y) \land (\forall z \in y)(z = x \lor z \in x)))$$

Powerset: $\forall x \exists y ((\forall z \subseteq x)(z \in y) \land (\forall z \in y)(\forall w \in z)(w \in x))$ Transitive Containment: $\forall x \exists y (\bigcup y \subseteq y \land (\forall z \in x)(z \in y))$ Axiom of Choice: $\forall x \exists R(R)$ is a well-ordering of x

Before showing that $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ satisfies Axiom H we first show that $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ is an end-extension of \mathcal{I} and a cofinal extension of \mathcal{M} .

Lemma 4.22 For all $x \in I$ and for all $[\tau] \in N_{\mathcal{U}, \mathbb{L}}$, if $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models ([\tau] \in [\hat{h}_x(c_0)])$ then there exists $y \in I$ such that

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left([\tau] = \left[\hat{h}_y(c_0) \right] \right) \text{ and } \mathcal{M} \models (y \in x).$$

Proof Let $x \in I$ and let $[\tau] \in N_{\mathcal{U}, \mathbb{L}}$. Suppose that $\tau = \hat{f}(c_{j_1}, \ldots, c_{j_n})$ where $f \in \mathcal{F}$ and $j_1 < \cdots < j_n \in \mathbb{L}$, and $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models ([\tau] \in [\hat{h}_x(c_0)])$. Therefore, by Lemma 4.17, there exists $X \in \mathcal{U}$ such that for all $x_1 < \cdots < x_n \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(x_1, \dots, x_n) \in x.$$

Since \mathcal{U} is I-complete, there is $Y \in \mathcal{U}$ and $y \in x^*$ such that for all $x_1 < \cdots < x_n \in Y^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(x_1, \dots, x_n) = y.$$

Therefore $(\hat{f}(c_{j_1},\ldots,c_{j_n})=\hat{h}_y(c_0))\in T_{\mathcal{U},\mathbb{L}}$. So by Lemma 4.17, $\mathcal{N}_{\mathcal{U},\mathbb{L}}\models([\tau]=[\hat{h}_y(c_0)])$; and since $y\in x^*$, $\mathcal{M}\models(y\in x)$.

Lemma 4.22 shows that for all $j \in Aut(\mathbb{L})$, $I_{fix(\check{j})} \supseteq I$.

Lemma 4.23 For all $[\tau] \in N_{\mathcal{U}, \mathbb{L}}$, there exists $x \in M$ such that,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models [\tau] \in \left[\hat{h}_x(c_0)\right].$$

Proof Let $[\tau] \in N_{\mathcal{U},\mathbb{L}}$. Suppose that $\tau = \hat{f}(c_{j_1},\ldots,c_{j_n})$ where $f \in M$ and $j_1 < \cdots < j_n \in \mathbb{L}$. Let $x \in M$ be such that $\mathcal{M} \models (\operatorname{rng}(f) \subseteq x)$. Now, for all $x_1 < \cdots < x_n \in \bar{\kappa}^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(x_1, \dots, x_n) \in x.$$

Therefore, $(\hat{f}(c_{j_1}, \dots, c_{j_n}) \in \hat{h}_x(c_0)) \in T_{\mathcal{U}, \mathbb{L}}$. So by Lemma 4.17,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models [\tau] \in \left[\hat{h}_{x}(c_{0})\right].$$



The fact that \mathcal{U} is *I*-tight also ensures that $\operatorname{Card}^{\mathcal{M}} \backslash \operatorname{Card}^{\mathcal{I}}$ is downward cofinal in $\operatorname{Card}^{\mathcal{N}_{\mathcal{U},\mathbb{L}}} \backslash \operatorname{Card}^{\mathcal{I}}$.

Lemma 4.24 If $\lambda \in \operatorname{Card}^{\mathcal{N}_{\mathcal{U},\mathbb{L}}} \backslash \operatorname{Card}^{\mathcal{I}}$ then there is $\mu \in \operatorname{Card}^{\mathcal{M}} \backslash \operatorname{Card}^{\mathcal{I}}$ with $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models ([\hat{h}_{\mu}(c_0)] \leq \lambda)$.

Proof Let $\lambda \in \operatorname{Card}^{\mathcal{N}_{\mathcal{U},\mathbb{L}}} \backslash \operatorname{Card}^{\mathcal{I}}$. Therefore $\lambda = [\hat{f}(c_{i_1}, \ldots, c_{i_n})]$ where $f \in \mathcal{F}$ and $i_1 < \cdots < i_n \in \mathbb{L}$, and $[\hat{f}(c_{i_1}, \ldots, c_{i_n})] \notin I$ and

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left(\left\lceil \hat{f}\left(c_{i_1},\ldots,c_{i_n}\right) \right\rceil \text{ is a cardinal} \right).$$

If there is $\mu \in M$ such that

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left(\left[\hat{f}\left(c_{i_1},\ldots,c_{i_n}\right) \right] = \left[\hat{h}_{\mu}(c_0) \right] \right)$$

then, by Lemma 4.20, $[\hat{f}(c_{i_1}, \ldots, c_{i_n})] \in \operatorname{Card}^{\mathcal{M}} \backslash \operatorname{Card}^{\mathcal{I}}$ and we are done. Therefore, assume that $[\hat{f}(c_{i_1}, \ldots, c_{i_n})] \in N_{\mathcal{U}, \mathbb{L}} \backslash M$. Since \mathcal{U} is I-tight and $[\hat{f}(c_{i_1}, \ldots, c_{i_n})] \notin I$, there is an I-large \mathcal{M} -cardinal $\mu \in M$ and $X \in \mathcal{U}$ such that for all $x_1 < \cdots < x_n \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \left(\hat{f}(x_1,\ldots,x_n) \notin H_{\mu}\right).$$

Since $[\hat{f}(c_{i_1}, \ldots, c_{i_n})]$ is an $\mathcal{N}_{\mathcal{U}, \mathbb{L}}$ -cardinal, this implies that for all $x < x_1 < \cdots < x_n \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \left(\hat{h}_{\mu}(x) \leq \hat{f}(x_1, \dots, x_n)\right).$$

And so, by Lemma 4.17, $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models (\hat{h}_{\mu}(c_0) \leq \hat{f}(c_{i_1},\ldots,c_{i_n})).$

We now turn to showing that $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ satisfies Axiom H.

Lemma 4.25 $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \forall x(|x| \ exists).$

Proof Let $\phi(x)$ be the formula

 $\exists f \exists \alpha ((\alpha \text{ is an ordinal}) \land (f : x \longrightarrow \alpha \text{ is an injection})).$

Now, $\phi(x)$ is Σ_1 and $\mathcal{M} \models \forall x \phi(x)$. Therefore, by Lemma 4.20, $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models \forall x \phi(x)$ and the Lemma follows.

Lemma 4.26 $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models Axiom H.$

Proof Let $u \in N_{\mathcal{U}, \mathbb{L}}$. Using Lemma 4.25, let $\lambda \in N_{\mathcal{U}, \mathbb{L}}$ be such that

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models (\lambda \text{ is a cardinal}) \land (|u| \leq \lambda).$$



П

By Lemma 4.23 there is $\mu \in M$ such that μ is an \mathcal{M} -cardinal and $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models (\lambda \leq \hat{h}_{\mu}(c_0))$. Let $T \in M$ be such that $\mathcal{M} \models (T = H_{\mu^+})$. Therefore

$$\mathcal{M} \models \forall z \left(\left(\bigcup z \subseteq z \land |z| \le \mu \right) \Rightarrow z \subseteq T \right).$$

So, by Lemma 4.20,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \forall z \left(\left(\bigcup z \subseteq z \land |z| \le [\hat{h}_{\mu}(c_0)] \right) \Rightarrow z \subseteq [\hat{h}_T(c_0)] \right)$$
 and
$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left(\bigcup [\hat{h}_T(c_0)] \subseteq [\hat{h}_T(c_0)] \right).$$

In particular

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \forall z \left(\left(\bigcup z \subseteq z \land |z| \le |u| \right) \Rightarrow z \subseteq [\hat{h}_T(c_0)] \right).$$

And this shows that

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \forall u \exists T \left(\bigcup T \subseteq T \land \forall z \left(\left(\bigcup z \subseteq z \land |z| \le |u| \right) \Rightarrow z \subseteq T \right) \right).$$

Combining Lemma 4.21 and Lemma 4.26 we get:

Theorem 4.27
$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \text{MOST}.$$

We now turn to showing that if $j \in \operatorname{Aut}(\mathbb{L})$ has no fixed points then $\mathcal{I}_{\operatorname{fix}}(\check{j}) = \mathcal{I}$ and $\operatorname{fix}(\check{j}) = M$.

Theorem 4.28 If $j \in Aut(\mathbb{L})$ has no fixed points then $fix(\check{j}) = M$.

Proof Let $j \in \operatorname{Aut}(\mathbb{L})$ have no fixed points. It follows immediately from (2) that for all $x \in M$, \check{j} fixes $[\hat{h}_x(c_0)]$. Therefore, we need to show that if $[\tau] \in N_{\mathcal{U},\mathbb{L}}$ and $\check{j}([\tau]) = [\tau]$ then there exists $x \in M$ such that $[\tau] = [\hat{h}_x(c_0)]$. Let $\tau = \hat{f}(c_{i_0}, \ldots, c_{i_n})$ where $f \in \mathcal{F}$ and $i_0 < \cdots < i_n \in \mathbb{L}$. Assume that $\check{j}([\tau]) = [\tau]$. Therefore $[\hat{f}(c_{i_0}, \ldots, c_{i_n})] = [\hat{f}(c_{j(i_0)}, \ldots, c_{j(i_n)})]$. Since \mathcal{U} is canonically Ramsey, there is $X \in \mathcal{U}$ and $\Delta \subseteq n+1$ such that for all $x_0 < \cdots < x_n$ and $x_0 < \cdots < x_n$ in $x_n \in \mathbb{L}$.

$$\left(\mathcal{M}_{\mathcal{F}} \models \hat{f}(x_0,\ldots,x_n) = \hat{f}(y_0,\ldots,y_n)\right)$$
 if and only if $(\forall m \in \Delta)(x_m = y_m)$.

Therefore, for all $j_0 < \cdots < j_n$ and $k_0 < \cdots < k_n$ in \mathbb{L} ,

$$[\hat{f}(c_{j_0}, \dots, c_{j_n})] = [\hat{f}(c_{k_0}, \dots, c_{k_n})]$$
 if and only if $(\forall m \in \Delta)(j_m = k_m)$.

Therefore, since j has no fixed points, $\Delta = \emptyset$. And so there exists $y \in M$ such that for all $x_0 < \cdots < x_n \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(x_0,\ldots,x_n) = y.$$



Therefore $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models ([\tau] = [\hat{h}_{v}(c_0)]).$

Lemma 4.29 If $x \in M$ then $\mathcal{N}_{\mathcal{U}, \mathbb{L}} \models (\text{TC}([\hat{h}_x(c_0)]) = [\hat{h}_{\text{TC}(x)}(c_0)]).$

Proof Let $x \in M$. Now, for all $y \in \bar{\kappa}^*$,

$$\mathcal{M}_{\mathcal{F}} \models \mathrm{TC}\left(\hat{h}_{x}(y)\right) = \hat{h}_{\mathrm{TC}(x)}(y).$$

Therefore $(TC(\hat{h}_x(c_0)) = \hat{h}_{TC(x)}(c_0)) \in T_{\mathcal{U},\mathbb{L}}$. Now, "Y is the transitive closure of X" can be expressed by a $\Delta_0^{\mathcal{D}}$ -formula:

$$X \subseteq Y \land \bigcup Y \subseteq Y \land (\forall x \subseteq Y) \left(X \subseteq x \land \bigcup x \subseteq x \Rightarrow Y \subseteq x \right).$$

Therefore, by Lemma 4.17,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \mathrm{TC}\left(\left[\hat{h}_{x}(c_{0})\right]\right) = \left[\hat{h}_{\mathrm{TC}(x)}(c_{0})\right].$$

Theorem 4.30 If $j \in Aut(\mathbb{L})$ has no fixed points then $\mathcal{I}_{fix}(\check{j}) = \mathcal{I}$.

Proof Let $j \in \operatorname{Aut}(\mathbb{L})$ have no fixed points. Lemma 4.22 implies that the embedding of \mathcal{M} into $\mathcal{N}_{\mathcal{U},\mathbb{L}}$ embeds \mathcal{I} into an initial segment of $\mathcal{N}_{\mathcal{U},\mathbb{L}}$. Therefore if $x \in I$ and $y \in \mathcal{N}_{\mathcal{U},\mathbb{L}}$ is such that $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models (y \in \operatorname{TC}(\{x\}))$ then $\check{j}(y) = y$. Therefore, we need to show that if $x \in \mathcal{N}_{\mathcal{U},\mathbb{L}} \setminus I$ then there is some $y \in (\operatorname{TC}(\{x\})^{\mathcal{N}_{\mathcal{U},\mathbb{L}}})^*$ that is moved by \check{j} . In light of Theorem 4.28 it is enough to show that for all $x \in \mathcal{N}_{\mathcal{U},\mathbb{L}} \setminus I$, there exists $y \in \mathcal{N}_{\mathcal{U},\mathbb{L}} \setminus M$ with $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models (y \in \operatorname{TC}(\{x\}))$. Let $x \in \mathcal{N}_{\mathcal{U},\mathbb{L}} \setminus I$. If $x \in \mathcal{N}_{\mathcal{U},\mathbb{L}} \setminus M$ then we are done, so assume that $x \in M \setminus I$. It follows that $|\operatorname{TC}(\{x\})|^{\mathcal{M}} \notin I$; hence $|\operatorname{TC}(\{x\})|^{\mathcal{M}}$ is an I-large \mathcal{M} -cardinal. Therefore there is $X \in \mathcal{U}$ such that $\mathcal{M} \models (|X| \leq |\operatorname{TC}(\{x\})|)$. Let $g \in M$ be such that

$$\mathcal{M} \models (g : X \longrightarrow TC(\{x\})) \land (g \text{ is an injection}).$$

Therefore there exists $f \in M$ such that $\mathcal{M} \models (f : \bar{\kappa} \longrightarrow TC(\{x\}))$ and for all $z_1 < z_2 \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(z_1) \neq \hat{f}(z_2).$$

Therefore, for all $w \in M$, $\mathcal{N}_{\mathcal{U},\mathbb{L}} \models ([\hat{f}(c_0)] \neq [\hat{h}_w(c_0)])$. We need to show that $[\hat{f}(c_0)]$ is in the transitive closure of $[\hat{h}_{\{x\}}(c_0)]$ in $\mathcal{N}_{\mathcal{U},\mathbb{L}}$. For all $z \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \left(\hat{f}(z) \in \hat{h}_{\mathrm{TC}(\{x\})}(z)\right).$$



Therefore, for all $i \in \mathbb{L}$, $(\hat{f}(c_i) \in \hat{h}_{TC(\{x\})}(c_i)) \in T_{\mathcal{U},\mathbb{L}}$. So, by Lemma 4.17,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left[\hat{f}(c_0)\right] \in \left[\hat{h}_{TC(\{x\})}(c_0)\right].$$

Therefore, by Lemma 4.29,

$$\mathcal{N}_{\mathcal{U},\mathbb{L}} \models \left[\hat{f}(c_0)\right] \in TC\left(\left[\hat{h}_x(c_0)\right]\right),$$

which completes the proof of the theorem.

This completes the proof of Theorem 4.1.

5 Realizing countable models as *H*-cuts

This section tackles the question of which countable \mathcal{L} -structures can be realised as an H-cut of a model of set theory. We show:

- every countable transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -collection can be realized as an H-cut of a model of MOST + $\Delta_0^{\mathcal{P}}$ -collection,
- every countable recursively saturated model of MOST + $\Delta_0^{\mathcal{P}}$ -collection can be realised as an H-cut of a model of MOST + $\Delta_0^{\mathcal{P}}$ -collection,
- every countable model of ZFC can be realised as an H-cut of a model of ZFC.

Combined with the results of Sect. 4 we then have:

- 1. every countable transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -collection can be realised as $\mathcal{I}_{\text{fix}}(j)$ for some non-trivial automorphism j of a model of MOST,
- 2. every countable recursively saturated model of MOST + $\Delta_0^{\mathcal{P}}$ -collection can be realised as $\mathcal{I}_{\text{fix}}(j)$ for some non-trivial automorphism j of a model of MOST,
- 3. every countable model of ZFC can be realised as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of ZFC.

Together with the results of Sect. 3, (1) and (2) yield a complete characterisation of the countable transitive and countable recursively saturated \mathcal{L} -structures satisfying infinity which can be realised as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of MOST. Since every consistent theory with an infinite model has a countable recursively saturated model (Theorem 2.6), (2) implies that the theory of the class of \mathcal{L} -structures satisfying infinity that appear as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of MOST is exactly MOST $+\Delta_{\mathcal{D}}^{\mathcal{D}}$ -collection.

5.1 Countable transitive models of MOST + $\Delta_0^{\mathcal{P}}$ -collection

The realisation of a countable transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -collection as an H-cut of a model of MOST + $\Delta_0^{\mathcal{P}}$ -collection will be achieved by applying a theorem, due to Harvey Friedman [8], which combined with a result of Mathias [17] shows that every



countable transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -collection can be toplessly end-extended to a model of MOST + $\Delta_0^{\mathcal{P}}$ -collection.

In [8] Friedman studies countable transitive models of a theory that he calls Power Admissible Set Theory (PAdm^s). This theory is axiomatised using classes of \mathcal{L} -formulae that Friedman calls $\Delta_0^s(\mathcal{P})$. Let $\mathcal{L}_{\mathcal{P}}$ be the extension of the \mathcal{L} obtained by adding a new unary function symbol \mathcal{P} . Define pseudo- $\Delta_0^s(\mathcal{P})$ to be the class of $\Delta_0(\mathcal{L}_{\mathcal{P}})$ formulae that contain quantification in the form $\exists x \in y \text{ or } \forall x \in y \text{ where } x$ and y are distinct *variables*. The class $\Delta_0^s(\mathcal{P})$ is obtained by translating the formulae in pseudo- $\Delta_0^s(\mathcal{P})$ into \mathcal{L} -formulae using translations generated by the defining axiom $x \in \mathcal{P}(y) \iff \forall z (z \in x \Rightarrow z \in y)$.

Definition 5.1 *Power Admissible Set Theory* (PAdm^s) *is obtained from* KP *by adding powerset,* $\Delta_0^s(\mathcal{P})$ *-collection and* \mathcal{L} *-foundation.*

Friedman [8, Section 1] also introduces a class theory that corresponds to Power Admissible Set Theory. We use \mathcal{L}_{Cl} to denote the two-sorted extension of \mathcal{L} with set variables x, y, z, \ldots and class variables X, Y, Z, \ldots . The well-formed formulae of \mathcal{L}_{Cl} are built inductively from atomic formulae in the form $x \in y, x = y, X = Y, x = Y$ and $y \in X$ using the connectives and quantifiers of first-order logic. The class Δ_0^c is the smallest class of \mathcal{L}_{Cl} -formulae that contains all atomic formulae in the form $x = y, x = Y, x \in y$ and $y \in X$, contains all compound formulae formed using the connectives of first-order logic, and is closed under quantification in the form $\exists x \in y$ and $\forall x \in y$ where x and y are distinct variables. The class Σ^c is the smallest class of \mathcal{L}_{Cl} -formulae that contains all Δ_0^c -formulae, contains all compound formulae formed using the connectives \wedge and \vee of first-order logic, and is closed under quantification in the form $\exists x$ and $\forall x \in y$ where x and y are distinct variables.

Definition 5.2 *Power Admissible Class Theory* (PAdm^c) *is the* \mathcal{L}_{Cl} *-theory with axioms: extensionality for both sets and classes,* $\forall x \exists Y (x = Y)$, *pairing and union for sets, powerset for sets,* \mathcal{L} *-foundation, and the following:*

 $(\Delta_0^c$ -separation) for all Δ_0^c -formulae $\phi(x, \vec{Z})$,

$$\forall \vec{Z} \forall w \exists y \forall x (x \in y \iff x \in w \land \phi(x, \vec{Z}))$$

 $(\Delta_0^c$ -collection) for all Δ_0^c -formulae $\phi(x, y, \vec{Z})$,

$$\forall \vec{Z} \forall w ((\forall x \in w) \exists y \phi(x, y, \vec{Z}) \Rightarrow \exists C (\forall x \in w) (\exists y \in C) \phi(x, y, \vec{Z}))$$

 $(\Delta^c$ -CA) for all Σ^c -formulae $\phi(x, \vec{Z})$ and $\psi(x, \vec{Z})$,

$$\forall \vec{Z} (\forall x (\phi(x, \vec{Z}) \iff \neg \psi(x, \vec{Z})) \Rightarrow \exists Y \forall x (x \in Y \iff \phi(x, \vec{Z})))$$

(Class Powerset)

$$\exists X \forall y (y \in X \iff \exists w \exists z (y = \langle w, z \rangle \land \forall v (v \in z \iff \forall u (u \in v \Rightarrow u \in w))))$$

¹ An explicit list of translations can be found in [17].



Friedman [8, Theorem 1.6] notes that PAdm^c is a conservative extension of PAdm^s:

Theorem 5.1 (Friedman) If $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ is an \mathcal{L} -structure with $\mathcal{M} \models PAdm^s$ then there exists C such that $\langle M, C, \in^{\mathcal{M}} \rangle \models PAdm^c$.

On the other hand, as shown by Mathias [17, Metatheorem 6.20] we have:

Theorem 5.2 (Mathias) The theories PAdm^s + Infinity and $KP^{\mathcal{P}}$ have the same transitive models.

Note that apart from instances of the $\Pi_1^{\mathcal{P}}$ -foundation scheme, every axiom of $KP^{\mathcal{P}}$ is also an axiom of MOST + $\Delta_0^{\mathcal{P}}$ -collection. Therefore, every transitive model of $MOST + \Delta_0^{\mathcal{P}}$ -collection is a model of $KP^{\mathcal{P}}$.

Using a version of the Barwise Compactness Theorem, Friedman [8, Theorem 2.3] shows that any countable transitive model of PAdm^s has a topless powerset-preserving end extension that is a model of PAdm^s.

Theorem 5.3 (Friedman) Let $\mathcal{M} = \langle M, C, \in \rangle$ be countable and transitive with $\mathcal{M} \models$ PAdm^c. If $T \in C$ is an \mathcal{L} -theory with $\langle M, \in \rangle \models T$ then there is an \mathcal{L} -structure $\mathcal{N} = \langle N, \in^{\mathcal{N}} \rangle$ such that

- (i) $N \neq M$,

(i)
$$N \neq M$$
,
(ii) $\langle M, \in \rangle \subseteq_{\text{topless}}^{\mathcal{P}} \mathcal{N}$,
(iii) $\mathcal{N} \models T$.

Combining Theorems 5.1, 5.2 and 5.3 yields:

Corollary 5.4 If $\langle I, \in \rangle$ is a countable transitive model of MOST + $\Delta_0^{\mathcal{P}}$ -collection then there is $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ with $\mathcal{M} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection + \mathcal{L} -foundation and $I \subseteq M$ is an H-cut of M.

Combining this with the results of Sects. 3 and 4 gives a characterisation of the countable transitive structures (I, \in) with $\omega \in I$ that can be realised as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of MOST.

Theorem 5.5 Let (I, \in) be a countable transitive structure with $\omega \in I$. The following are equivalent:

- (I) $\langle I, \in \rangle \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection
- (II) there is an \mathcal{L} -structure $\mathcal{M} \models \text{MOST}$ and a non-trivial automorphism $j: \mathcal{M} \longrightarrow$ M such that

$$\mathcal{I}_{\text{fix}}(j) = \langle I, \in \rangle.$$

5.2 Countable recursively saturated models of MOST + $\Delta_0^{\mathcal{P}}$ -collection

In this section we will show that every countable recursively saturated model of $MOST + \Delta_0^{\mathcal{P}}$ -collection can be realised as an H-cut of a model of MOST +



 $\Delta_0^{\mathcal{P}}$ -collection. This will be achieved by proving the following refined version of Friedman's Self-Embedding Theorem [8, Section 4] for non-standard models of set theory:

Theorem 5.6 If $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ is a countable recursively saturated model of MOST + $\Delta_0^{\mathcal{P}}$ -collection then there exists an embedding $h : \mathcal{M} \longrightarrow \mathcal{M}$ such that $\operatorname{rng}(h) \subseteq M$ is an H-cut of \mathcal{M} .

Combined with Sects. 3 and 4 Theorem 5.6 yields a characterisation of the countable recursively saturated structures satisfying infinity that can be realised as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of MOST.

Theorem 5.7 Let $\mathcal{I} = \langle I, \in^{\mathcal{I}} \rangle$ be a countable recursively saturated structure with $\mathcal{I} \models Infinity$. The following are equivalent:

- (I) $\mathcal{I} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection.
- (II) there is an \mathcal{L} -structure $\mathcal{M} \models \text{MOST}$ and a non-trivial automorphism $j: \mathcal{M} \longrightarrow \mathcal{M}$ such that

$$\mathcal{I}_{\text{fix}}(j) = \mathcal{I}.$$

Combined with the observation that every consistent theory T extending MOST + $\Delta_0^{\mathcal{P}}$ -collection has a countable recursively saturated model (Theorem 2.6), Theorem 5.6 also yields the first-order theory of the class of \mathcal{L} -structures that can appear as $\mathcal{I}_{\text{fix}}(j)$ for some non-trivial automorphism j of a model of MOST.

Theorem 5.8 *Let* T *be a complete, consistent* \mathcal{L} *-theory such that* $T \vdash Infinity.$ *The following are equivalent:*

- (I) $T \vdash MOST + \Delta_0^{\mathcal{P}}$ -collection.
- (II) there is an \mathcal{L} -structure $\mathcal{M} \models \text{MOST}$ and a non-trivial automorphism $j: \mathcal{M} \longrightarrow \mathcal{M}$ such that

$$\mathcal{I}_{\text{fix}}(j) \models T$$
.

We now turn to proving Theorem 5.6. For the remainder of this section let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be a countable recursively saturated structure with $\mathcal{M} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection. We need to build an embedding $h: \mathcal{M} \longrightarrow \mathcal{M}$ such that $\text{rng}(h) \subseteq M$ is an H-cut of \mathcal{M} . This will be achieved by a three-stage back-and-forth construction in which the 'back' and 'forth' steps are similar to the proofs of the classical versions of Friedman's Self-Embedding Theorem for set theory and arithmetic (e.g., as in [8, Section 4] and [13, Chapter 12]). The 'third' stage of the back-and-forth construction will be used to ensure that the range of the embedding sits toplessly inside \mathcal{M} . Let $\langle m_i \mid i \in \omega \rangle$ be an enumeration of M in which each element of M appears infinitely often. Let $\langle \lambda_i \mid i \in \omega \rangle$ be an enumeration of the class

$$\{\lambda \in M \mid \mathcal{M} \models (\lambda \text{ is a limit cardinal})\}.$$



Lemma 5.9 There exists $X \in M$ such that

$$\mathcal{M} \models \exists \kappa ((\kappa \text{ is a cardinal}) \land \forall y (|\text{TC}(y)| < \kappa \iff y \in X))$$

and for all $\Sigma_1^{\mathcal{P}}$ -sentences ϕ , if $\mathcal{M} \models \phi$ then $\mathcal{M} \models \phi^X$.

Proof We use the fact that \mathcal{M} is recursively saturated. Let $\Gamma(x)$ be the one-type that consists of the following formulae:

- (i) $\exists \kappa ((\kappa \text{ is a cardinal}) \land \forall y (|\text{TC}(y)| < \kappa \iff y \in x))$
- (ii) for all $\Sigma_1^{\mathcal{P}}$ -sentences ϕ ,

$$\phi \Rightarrow \phi^x$$
.

Note that $\Gamma(x)$ is a recursive type. We need to show that $\Gamma(x)$ is finitely realised. Suppose that $\Delta(x) \subseteq \Gamma(x)$ is finite and that the instances of (ii) mentioned in $\Delta(x)$ are exactly

$$\psi_i \Rightarrow \psi_i^x$$
 where ψ_i is a $\Sigma_1^{\mathcal{P}}$ -sentence for $0 \le i < k$.

Without loss of generality we may assume that for all $0 \le i < k$, $\mathcal{M} \models \psi_i$. Suppose that for all $0 \le i < k$, ψ_i is the sentence $\exists z \theta_i(z)$ where $\theta_i(z)$ is a $\Delta_0^{\mathcal{P}}$ -formula. Let $a_0, \ldots, a_{k-1} \in M$ be such that for all $0 \le i < k$,

$$\mathcal{M} \models \theta_i(a_i).$$

Now, work inside \mathcal{M} . Let $\kappa = \sup\{|\mathrm{TC}(\mathcal{P}(a_i))| \mid 0 \leq i < k\}$. The fact that we are working in a model of MOST ensures that H_{κ^+} exists. Since $a_0, \ldots, a_{k-1}, \mathcal{P}(a_0), \ldots, \mathcal{P}(a_{k-1}) \in H_{\kappa^+}$ and $\theta_0(z), \ldots, \theta_{k-1}(z)$ are $\Delta_0^{\mathcal{P}}$, it follows that

$$\langle H_{\kappa^+}, \in \rangle \models \theta_i(a_i) \text{ for all } 0 \le i < k.$$

Since $H_{\kappa^+}^{\mathcal{M}}$ also satisfies (i), $H_{\kappa^+}^{\mathcal{M}}$ realises $\Delta(x)$ in \mathcal{M} . This shows that $\Gamma(x)$ is finitely realised. Since \mathcal{M} is recursively saturated, there is $X \in \mathcal{M}$ that realises $\Gamma(x)$. This proves the lemma.

Lemma 5.9 allows us to initiate the back-and-forth proof of Theorem 5.6. Let $X_0 \in M$ be such that

$$\exists \kappa ((\kappa \text{ is a cardinal}) \land \forall y (|\text{TC}(y)| < \kappa \iff y \in X_0)) \text{ and}$$
 for all $\Sigma_1^{\mathcal{P}}$ -sentences ϕ , if $\mathcal{M} \models \phi \text{ then } \mathcal{M} \models \phi^{X_0}$.

We will construct an embedding $h: \mathcal{M} \longrightarrow \mathcal{M}$ by constructing sequences $\langle u_i \mid i \in \omega \rangle$, $\langle v_i \mid i \in \omega \rangle$ and $\langle X_i \mid i \in \omega \rangle$ of elements of M, such that for all $i, j \in \omega$: if i < 2j then $\mathcal{M} \models v_i \in X_j$, and if i < j then $\mathcal{M} \models X_j \subseteq X_i$. Then we define:

$$h(u_i) = v_i$$
 for all $i \in \omega$.

At stage j, after having defined $u_0, \ldots, u_{2i-1}, v_0, \ldots, v_{2i-1}$ and X_i , we will ensure that the following condition is maintained:

 (\dagger_i) for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0,\ldots,x_{2i-1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2j-1})$$
 then $\mathcal{M} \models \phi^{X_j}(v_0, \dots, v_{2j-1})$.

Suppose that we have chosen $u_0, \ldots, u_{2k-1}, v_0, \ldots, v_{2k-1}$ and X_k and maintained (\dagger_k) . Stage k of the construction comprises three steps:

Step 1 This step will ensure that the image of h sits toplessly inside \mathcal{M} . Our aim at stage k is to prevent $(H_{\lambda_k}^{\mathcal{M}})^*$ from being the image of h (where λ_k is as in the definition preceding Lemma 5.9). Consider the following conditions:

- (a) $\mathcal{M} \models (H_{\lambda_k} \subseteq X_k)$ (b) $\langle (H_{\lambda_k}^{\mathcal{M}})^*, \in^{\mathcal{M}} \rangle \models \text{MOST}$
- (c) $v_0, \ldots, v_{2k-1} \in (H_{\lambda_k}^{\mathcal{M}})^*$
- (d) for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k-1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k-1})$$
 then $\mathcal{M} \models \phi^{H_{\lambda_k}}(v_0, \dots, v_{2k-1})$.

If any of the conditions (a), (b), (c) or (d) fail then it is already impossible for $(H_{\lambda_k}^{\mathcal{M}})^*$ to be the image of h. Therefore, if any of (a), (b), (c) or (d) fail then let $X_{k+1} = X_k$. Since \dagger_k holds, for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k-1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k-1})$$
 then $\mathcal{M} \models \phi^{X_{k+1}}(v_0, \dots, v_{2k-1}).$ (3)

If (a), (b), (c) and (d) all hold then we will choose X_{k+1} so that condition (a) fails for X_k . This will prevent $(H_{\lambda_k}^{\mathcal{M}})^*$ from being the image of h.

Lemma 5.10 If conditions (a), (b), (c) and (d) all hold then there exists $X \in (H_{\lambda_L}^{\mathcal{M}})^*$ such that

$$\mathcal{M} \models \exists \kappa ((\kappa \text{ is a cardinal}) \land \forall y (|TC(y)| < \kappa \iff y \in X))$$

and for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0,\ldots,x_{2k-1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k-1})$$
 then $\mathcal{M} \models \phi^X(v_0, \dots, v_{2k-1})$.

Proof Assume that conditions (a), (b), (c) and (d) all hold. Let $\Gamma(x)$ be the one-type that contains the following formulae:

- (i) $x \in H_{\lambda_k}$,
- (ii) $\exists \kappa ((\kappa \text{ is a cardinal}) \land \forall y (|\text{TC}(y)| < \kappa \iff y \in x)),$
- (iii) for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \dots, x_{2k-1})$,

$$\phi(u_0, \ldots, u_{2k-1}) \Rightarrow \phi^x(v_0, \ldots, v_{2k-1}).$$



 $\Gamma(x)$ is a recursive type. We need to show that $\Gamma(x)$ is finitely realised. Suppose that $\Delta(x) \subseteq \Gamma(x)$ is finite and that the instances of (iii) mentioned in $\Delta(x)$ are exactly

$$\psi_i(u_0, \dots, u_{2k-1}) \Rightarrow \psi_i^x(v_0, \dots, v_{2k-1})$$
 where ψ_i is a $\Sigma_1^{\mathcal{P}}$ -formula for $0 \le i < m$.

Without loss of generality we may assume that for all $0 \le i < m$,

$$\mathcal{M} \models \psi_i(u_0, \ldots, u_{2k-1}).$$

Suppose that for each $0 \le i < m$, the formula $\psi_i(x_0, \ldots, x_{2k-1})$ is $\exists z \theta_i(z, x_0, \ldots, x_{2k-1})$ where $\theta_i(z, x_0, \ldots, x_{2k-1})$ is a $\Delta_0^{\mathcal{P}}$ -formula. It follows from (d) that for all $0 \le i < m$,

$$\mathcal{M} \models \psi_i^{H_{\lambda_k}}(v_0, \ldots, v_{2k-1}).$$

Let $a_0, \ldots, a_{m-1} \in (H_{\lambda_k}^{\mathcal{M}})^*$ be such that for all $0 \le i < m$,

$$\mathcal{M} \models \theta_i^{H_{\lambda_k}}(a_i, v_0, \dots, v_{2k-1}).$$

Work inside \mathcal{M} . Let $\mu_1 = \sup\{|\mathrm{TC}(\mathcal{P}(a_i))| \mid 0 \le i < m\}$ and let $\mu_2 = \sup\{|\mathrm{TC}(\mathcal{P}(v_i))| \mid 0 \le i < 2k\}$. Let $\kappa = \max\{\mu_1, \mu_2\}$. The fact that we are working in a model of MOST ensures that H_{κ^+} exists. It follows from condition (b) that $H_{\kappa^+} \in H_{\lambda_k}$. Now,

$$a_0, \ldots, a_{m-1}, \mathcal{P}(a_0), \ldots, \mathcal{P}(a_{m-1}), v_0, \ldots, v_{2k-1}, \mathcal{P}(v_0), \ldots, \mathcal{P}(v_{2k-1}) \in H_{\nu^+}$$

Therefore, since each $\theta_i(z, x_0, \dots, x_{2k-1})$ is $\Delta_0^{\mathcal{P}}$,

$$\langle H_{\kappa^+}, \in \rangle \models \theta_i(a_i, v_0, \dots, v_{2k-1}) \text{ for all } 0 \le i < m.$$

Therefore $H_{\kappa^+}^{\mathcal{M}}$ realizes $\Delta(x)$ in \mathcal{M} . This shows that $\Gamma(x)$ is finitely realized. Since \mathcal{M} is recursively saturated, it follows that there is an $X \in \mathcal{M}$ that realizes $\Gamma(x)$. This proves the lemma.

Let X_{k+1} be the point in \mathcal{M} guaranteed by Lemma 5.10. It follows that

$$\mathcal{M} \models \neg (H_{\lambda_k} \subseteq X_{k+1}).$$

This prevents $(H_{\lambda_k}^{\mathcal{M}})^*$ from being the image of h. We also have that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k-1})$, (3) holds. This completes Step 1.

Step 2 This is the usual 'fourth' step in the proof of Friedman's Embedding Theorem (see [13, Theorem 12.3]). Let $u_{2k} = m_k$. This choice will eventually ensure that the domain of h is all of M. We need to choose $v_{2k} \in X_{k+1}^*$ such that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k})$$
 then $\mathcal{M} \models \phi(v_0, \dots, v_{2k})$.



The following Lemma shows that we can successfully make this choice:

Lemma 5.11 There exists $v \in X_{k+1}^*$ such that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k})$$
 then $\mathcal{M} \models \phi(v_0, \dots, v_{2k-1}, v)$.

Proof We use the fact that \mathcal{M} is recursively saturated. Let $\Gamma(x)$ be the one-type that contains the following formulae:

- (i) $x \in X_{k+1}$, (ii) for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \dots, x_{2k})$,

$$\phi(u_0,\ldots,u_{2k}) \Rightarrow \phi^{X_{k+1}}(v_0,\ldots,v_{2k-1},x).$$

 $\Gamma(x)$ is a recursive type. We need to show that $\Gamma(x)$ is finitely realised. Suppose that $\Delta(x) \subseteq \Gamma(x)$ is finite and that the instances of (ii) mentioned in $\Delta(x)$ are exactly

$$\psi_i(u_0, \dots, u_{2k})$$

 $\Rightarrow \psi_i^{X_{k+1}}(v_0, \dots, v_{2k-1}, x)$, where ψ_i is a $\Sigma_1^{\mathcal{P}}$ -formula for $0 \le i < m$.

Without loss of generality we may assume that for all $0 \le i < m$,

$$\mathcal{M} \models \psi_i(u_0,\ldots,u_{2k}).$$

Suppose that for each $0 \le i < m$, the formulae $\psi_i(x_0, \dots, x_{2k})$ is $\exists z_i \theta_i(z_i, x_0, \dots, x_{2k})$ where $\theta_i(z_i, x_0, \dots, x_{2k})$ is a $\Delta_0^{\mathcal{P}}$ -formula and without loss of generality the z_i are distinct. We have

$$\mathcal{M} \models \exists y \exists z_0 \cdots \exists z_{m-1} \bigwedge_{0 \le i < m} \theta_i(z_i, u_0, \dots, u_{2k-1}, y).$$

Since for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0,\ldots,x_{2k-1})$, (3) holds, it follows that

$$\mathcal{M} \models (\exists y \in X_{k+1})(\exists z_0 \in X_{k+1}) \cdots (\exists z_{m-1} \in X_{k+1}) \bigwedge_{0 \le i < m} \theta_i^{X_{k+1}}(z_i, v_0, \dots, v_{2k-1}, y).$$

Let $v \in X_{k+1}^*$ be such that for all $0 \le i < m$,

$$\mathcal{M} \models \psi_i^{X_{k+1}}(v_0, \ldots, v_{2k-1}, v).$$

Therefore $v \in M$ realizes $\Delta(x)$. This shows that $\Gamma(x)$ is finitely realised. Since \mathcal{M} is recursively saturated, there is a $v \in M$ which realizes $\Gamma(x)$. This proves the lemma. \square

Let $v \in X_{k+1}^*$ be the point in \mathcal{M} guaranteed by Lemma 5.11. Let

$$v_{2k} = \begin{cases} v_i & \text{if } u_{2k} = u_i \text{ for some } 0 \le i < 2k, \\ v & \text{otherwise} \end{cases}$$



This ensures that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k})$$
 then $\mathcal{M} \models \phi^{X_{k+1}}(v_0, \dots, v_{2k})$. (4)

This completes Step 2.

Step 3 This is the usual 'back' step in the proof of Friedman's Embedding Theorem (see [13, Theorem 12.3]). This step will eventually ensure that \mathcal{M} is a powerset-preserving end-extension of the image of h. In this step we have two cases to consider:

Case 1 For all $0 \le i < 2k+1$, $m_k \nsubseteq v_i$. In this case let $v_{2k+1} = v_0$ and let $u_{2k+1} = u_0$. This choice clearly satisfies (\dagger_{k+1}) .

Case 2 There exists $0 \le i < 2k+1$ such that $m_k \subseteq v_i$. Note that in this case it immediately follows that $m_k \in X_{k+1}^*$. Let $v_{2k+1} = m_k$. We need to choose u_{2k+1} such that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k+1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k+1})$$
 then $\mathcal{M} \models \phi^{X_{k+1}}(v_0, \dots, v_{2k+1})$.

Lemma 5.12 For all $0 \le j < 2k + 1$, $\mathcal{P}^{\mathcal{M}}(v_j) \in X_{k+1}^*$.

Proof Let $0 \le j < 2k + 1$. Since $\mathcal{M} \models MOST$, it follows that

$$\mathcal{M} \models \exists y (\forall x \subseteq u_i(x \in y) \land \forall x \in y (x \subseteq u_i)).$$

By (4):

$$\mathcal{M} \models (\exists y \in X_{k+1})(\forall x \subseteq v_i(x \in y) \land \forall x \in y(x \subseteq v_i))^{X_{k+1}}.$$

Since every subset of v_j is a member of X_{k+1} , it follows that $\mathcal{P}^{\mathcal{M}}(v_j) \in X_{k+1}^*$. \square

The following lemma ensures that we can choose u_{2k+1} to satisfy (\dagger_{k+1}) :

Lemma 5.13 There exists $u \in M$ such that for all $\Sigma_1^{\mathcal{P}}$ -formulae $\phi(x_0, \ldots, x_{2k+1})$,

if
$$\mathcal{M} \models \phi(u_0, \dots, u_{2k}, u)$$
 then $\mathcal{M} \models \phi^{X_{k+1}}(v_0, \dots, v_{2k+1})$.

Proof We use the fact that \mathcal{M} is recursively saturated. Let $\Gamma(x)$ be the one-type that contains the following formulae:

- (i) $x \subseteq u_i$
- (ii) for all $\Delta_0^{\mathcal{P}}$ -formulae $\phi(z, x_0, \dots, x_{2k+1})$,

$$(\forall z \in X_{k+1})\phi^{X_{k+1}}(z, v_0, \dots, v_{2k+1}) \Rightarrow \forall z \phi(z, u_0, \dots, u_{2k}, x).$$



 $\Gamma(x)$ is a recursive type. We need to show that $\Gamma(x)$ is finitely realised. Suppose that $\Delta(x) \subseteq \Gamma(x)$ is finite and that the instances of (ii) mentioned in $\Delta(x)$ are exactly

$$(\forall z \in X_{k+1}) \psi_j^{X_{k+1}}(z, v_0, \dots, v_{2k+1}) \Rightarrow \forall z \psi_j(z, u_0, \dots, u_{2k}, x)$$
 where ψ_j is a $\Delta_0^{\mathcal{P}}$ -formula for $0 \le j < m$.

Without loss of generality we may assume that for all $0 \le j < m$,

$$\mathcal{M} \models (\forall z \in X_{k+1}) \psi_j^{X_{k+1}}(z, v_0, \dots v_{2k+1}).$$

Suppose, for a contradiction, that

$$\mathcal{M} \models (\forall x \subseteq u_i) \exists z \bigvee_{0 \le j < m} \neg \psi_j(z, u_0, \dots, u_{2k}, x).$$

Therefore

$$\mathcal{M} \models (\forall x \in \mathcal{P}(u_i)) \exists z \bigvee_{0 \leq j < m} \neg \psi_j(z, u_0, \dots, u_{2k}, x).$$

By applying $\Delta_0^{\mathcal{P}}$ -collection we can conclude that

$$\mathcal{M} \models \exists t (\forall x \in \mathcal{P}(u_i))(\exists z \in t) \bigvee_{0 \leq j < m} \neg \psi_j(z, u_0, \dots, u_{2k}, x).$$

Therefore

$$\mathcal{M} \models \exists t (\forall x \subseteq u_i)(\exists z \in t) \bigvee_{0 \le j < m} \neg \psi_j(z, u_0, \dots, u_{2k}, x).$$

By (4) and Lemma 5.12, we have

$$\mathcal{M} \models (\exists t \in X_{k+1})(\forall x \subseteq v_i)(\exists z \in t) \bigvee_{0 \le j < m} \neg \psi_j^{X_{k+1}}(z, v_0, \dots, v_{2k}, x). \tag{5}$$

But $v_{2k+1} \subseteq v_i$ and

$$\mathcal{M} \models (\forall z \in X_{k+1}) \bigwedge_{0 \le j < m} \neg \psi_j^{X_{k+1}}(z, v_0, \dots, v_{2k+1})$$

which contradicts (5). Therefore $\Gamma(x)$ is finitely satisfied. Since \mathcal{M} is recursively saturated, there is a $u \in M$ that realises $\Gamma(x)$. We claim that this $u \in M$ is the point that



is required by the lemma. Suppose that this is not the case and that $\phi(x_0, \dots, x_{2k+1})$ is a $\Sigma_1^{\mathcal{P}}$ -formula such that

$$\mathcal{M} \models \phi(u_0, \dots, u_{2k}, u)$$
 and $\mathcal{M} \models \neg \phi^{X_{k+1}}(v_0, \dots, v_{2k+1}).$

The formula $\phi(x_0, \dots, x_{2k+1})$ is equivalent to a formula $\exists z \theta(z, x_0, \dots, x_{2k+1})$ where $\theta(z, x_0, \dots, x_{2k+1})$ is a Δ_0^p -formula. Therefore

$$\mathcal{M} \models (\forall z \in X_{k+1}) \neg \theta^{X_{k+1}}(z, v_0, \dots, v_{2k+1}) \text{ and } \mathcal{M} \models \neg \forall z \neg \theta(z, u_0, \dots, u_{2k}, u).$$

But this contradicts the fact that u realizes $\Gamma(x)$. This proves the lemma. \square

Let $u \in M$ be the point guaranteed by Lemma 5.13. Let

$$u_{2k+1} = \begin{cases} u_j & \text{if } v_{2k+1} = v_j \text{ for some } 0 \le j < 2k+1, \\ u & \text{otherwise} \end{cases}$$

It follows from Lemma 5.13 that this choice of u_{2k+1} satisfies (\dagger_{k+1}) .

This completes the k^{th} stage of the back-and-forth construction. Continuing this process yields a map $h: \mathcal{M} \longrightarrow \mathcal{M}$. The fact that (\dagger_n) holds at the beginning of each stage $n \in \omega$ ensures that $h: \mathcal{M} \longrightarrow \mathcal{M}$ is an embedding. Step 2 of each stage ensures that the domain of $h: \mathcal{M} \longrightarrow \mathcal{M}$ is all of M. Step 3 of each stage ensures that

$$\operatorname{rng}(h) \subseteq_{a}^{\mathcal{P}} \mathcal{M}.$$

Step 1 of each stage ensures both that $rng(h) \neq M$ and that rng(h) sits toplessly in \mathcal{M} . This completes the proof of Theorem 5.6.

5.3 Countable models of ZFC

Sections 5.1 and 5.2 show that if \mathcal{I} is countable with $\mathcal{I} \models \text{MOST} + \Delta_0^{\mathcal{P}}$ -collection and \mathcal{I} is either transitive or recursively saturated then \mathcal{I} can be realised as $\mathcal{I}_{\text{fix}}(j)$ for some non-trivial automorphism j of a model of MOST. This raises the following:

Question 5.1 Can every countable model of MOST + $\Delta_0^{\mathcal{P}}$ -collection be realised as $\mathcal{I}_{\text{fix}}(j)$ for some non-trivial automorphism j of a model of MOST?

A positive answer to Question 5.1 would yield a complete classification of the countable \mathcal{L} -structures satisfying infinity that can be realised as $\mathcal{I}_{fix}(j)$ for some non-trivial automorphism j of a model of MOST. A result proved by John Hutchinson in [10, Theorem 3.1] shows that Question 5.1 has a positive answer if both MOST + $\Delta_0^{\mathcal{P}}$ -collection and MOST are replaced by ZFC.²

² Hutchinson's result was generalized to models of countable cofinality by the second author [12, Theorem 5.1].



Theorem 5.14 (Hutchinson) If $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ is countable with $\mathcal{M} \models \text{ZFC}$ then there is a countable \mathcal{L} -structure \mathcal{N} such that $\mathcal{M} \prec_{\text{topless}} \mathcal{N}$.

Note that if $\mathcal{M} \models \operatorname{ZFC}$ and $\mathcal{M} \prec_e \mathcal{N}$ then for all $\alpha \in \operatorname{Ord}^{\mathcal{M}}$, $V_{\alpha}^{\mathcal{M}} = V_{\alpha}^{\mathcal{N}}$, and so $\mathcal{M} \prec_e^{\mathcal{P}} \mathcal{N}$. Therefore Theorem 5.14 shows that every countable model of ZFC can be realised as an H-cut of a model of ZFC. Combined with the construction in Sect. 4 this shows:

Theorem 5.15 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be countable with $\mathcal{M} \models \text{ZFC}$. There exists a countable \mathcal{L} -structure $\mathcal{M} \prec \mathcal{N}$ and an automorphism $j : \mathcal{N} \longrightarrow \mathcal{N}$ such that $\mathcal{I}_{\text{fix}}(j) = \mathcal{M}$.

6 An extension of Togha's Theorem

Togha [25, Theorem 3] proves the following set-theoretic analogue of a result, due to Smoryński [23, Theorem A], about automorphisms of countable recursively saturated models of PA.

Theorem 6.1 (Togha) Let \mathcal{M} be a countable recursively saturated model of ZFC. If I is an ordinal cut of \mathcal{M} closed under cardinal exponentiation (Definition 3.2) then there is an automorphism $j: \mathcal{M} \longrightarrow \mathcal{M}$ such that I is the largest initial segment of $\operatorname{Ord}^{\mathcal{M}}$ that is pointwise fixed by j.

In light of the correspondence revealed by Lemmas 3.2 and 3.1, Togha's Theorem shows that if \mathcal{M} is a countable recursively saturated model of ZFC and $I \subseteq M$ is an H-cut of \mathcal{M} then there is $j \in \operatorname{Aut}(\mathcal{M})$ such that $\mathcal{I}_{\operatorname{fix}}(j) = \langle I, \in^{\mathcal{M}} \rangle$. In this section we will generalise Togha's Theorem by showing that if I is an H-cut of a countable recursively saturated model \mathcal{M} of ZFC then \mathcal{M} is endowed with continuum-many automorphisms j with the property that $\mathcal{I}_{\operatorname{fix}}(j)$ is exactly $\langle I, \in^{\mathcal{M}} \rangle$. This generalisation of Togha's Theorem is analogous to the generalisation of Smoryński's result proved by the first author in [4, Theorem B].

Theorem 6.2 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ be a countable recursively saturated model of ZFC. Let $I \subseteq M$ be an H-cut of \mathcal{M} . There is an embedding $j \mapsto \check{j}$ of $\operatorname{Aut}(\mathbb{Q})$ into $\operatorname{Aut}(\mathcal{M})$ such that if $j \in \operatorname{Aut}(\mathbb{Q})$ is non-trivial then $\mathcal{I}_{\operatorname{fix}}(\check{j}) = \langle I, \in^{\mathcal{M}} \rangle$.

Theorem 6.2 can also be viewed as a refinement of a result due to Schmerl [22] which shows that if \mathcal{M} is a countable recursively saturated model of ZFC (or indeed any theory equipped with enough coding) then there is a group embedding of Aut(\mathbb{Q}) into Aut(\mathcal{M}).

It should be noted that Lemma 3.2 shows that any H-cut is uniquely determined by its ordinal spine, which forms an ordinal cut that is closed under exponentiation and contains $\omega^{\mathcal{M}}$. Therefore Theorem 6.2 is equivalent to Theorem C mentioned in the abstract.

Before proving Theorem 6.2 we first need to prove a model-theoretic result that allows us to expand cofinal elementary extensions of models of ZFC. We show that if $\mathcal{M} = \langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$ is a structure with $X^{\mathcal{M}} \subseteq M$ such that \mathcal{M} satisfies ZFC plus the full collection scheme in the language of \mathcal{M} , and \mathcal{N} is an \mathcal{L} -structure with



 $\langle M, \in^{\mathcal{M}} \rangle \prec_{\mathrm{cf}} \mathcal{N}$, then \mathcal{N} can be expanded to a structure \mathcal{N}' in the language of \mathcal{M} that is a cofinal elementary extension of \mathcal{M} . This model-theoretic tool is the set-theoretic version of a result that was independently proved by Schmerl [21, Theorem 1.2] and Kotlarski [15, Theorem 8] for cofinal elementary extensions of models of PA. We use \mathcal{L}_X to denote the language obtained by extending \mathcal{L} with a new unary predicate X. Recall that ZFC(X) is obtained from ZFC by adding the schemes of \mathcal{L}_X -separation and \mathcal{L}_X -collection.

Theorem 6.3 Let $\mathcal{M} = \langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$ be an \mathcal{L}_X -structure with $\mathcal{M} \models \mathsf{ZFC}(X)$. If $\mathcal{N} = \langle N, \in^{\mathcal{N}} \rangle$ is such that $\langle M, \in^{\mathcal{M}} \rangle \prec_{\mathsf{cf}} \mathcal{N}$ then there exists $X^{\mathcal{M}} \subseteq X^{\mathcal{N}} \subseteq N$ such that $\mathcal{M} \prec \langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle$.

Proof Let $\mathcal{N} = \langle N, \in^{\mathcal{N}} \rangle$ be such that $\langle M, \in^{\mathcal{M}} \rangle \prec_{\mathrm{cf}} \mathcal{N}$. Note that this immediately implies that $\mathcal{N} \models \mathrm{ZFC}$. We begin by defining $X^{\mathcal{N}}$. Define $F: \mathrm{Ord}^{\mathcal{M}} \longrightarrow M$ such that for all $\alpha \in \mathrm{Ord}^{\mathcal{M}}$,

$$\mathcal{M} \models \forall y (y \in F(\alpha) \iff y \in V_{\alpha} \land X(y)).$$

The fact that $\mathcal{M} \models \operatorname{ZFC}(X)$ ensures that $F(\alpha)$ exists for all $\alpha \in \operatorname{Ord}^{\mathcal{M}}$. Also note that since $M \subseteq N$, for all $\alpha \in \operatorname{Ord}^{\mathcal{M}}$, $F(\alpha) \in N$. Define $X^{\mathcal{N}} \subseteq N$ by: for all $x \in N$,

 $x \in X^{\mathcal{N}}$ if and only if there exists $\alpha \in \operatorname{Ord}^{\mathcal{M}}$ such that $\mathcal{N} \models (x \in F(\alpha))$.

The fact that $\langle M, \in^{\mathcal{M}} \rangle \prec_{\mathrm{cf}} \mathcal{N}$, implies that $X^{\mathcal{M}} \subseteq X^{\mathcal{N}}$. We will prove by induction on $n \in \omega$ that $\mathcal{M} \prec_n \langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle$. Firstly, note that since $\langle M, \in^{\mathcal{M}} \rangle \prec \mathcal{N}$, for all $\alpha \in \mathrm{Ord}^{\mathcal{M}}$, $V_{\alpha}^{\mathcal{M}} = V_{\alpha}^{\mathcal{N}}$. Therefore, for all $\alpha \in \mathrm{Ord}^{\mathcal{M}}$,

$$\mathcal{M} \models \forall x (x \in F(\alpha) \iff x \in V_{\alpha} \land X(x)) \text{ and } \mathcal{N}$$
$$\models \forall x (x \in F(\alpha) \iff x \in V_{\alpha} \land X(x)).$$

And, for all $\Delta_0(\mathcal{L}_X)$ -formulae $\phi(x_1, \ldots, x_m)$,

(I) for all $a_1, \ldots, a_m \in N$, if $\alpha \in \text{Ord}^{\mathcal{M}}$ is such that $a_1, \ldots, a_m \in (V_{\alpha}^{\mathcal{N}})^*$, then

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \phi(a_1, \dots, a_m) \text{ if and only if } \mathcal{N}$$

 $\models (\langle V_{\alpha}, \in, F(\alpha) \rangle \models \phi(a_1, \dots, a_m)).$

(II) for all $a_1, \ldots, a_m \in M$, if $\alpha \in \text{Ord}^{\mathcal{M}}$ is such that $a_1, \ldots, a_m \in (V_{\alpha}^{\mathcal{M}})^*$, then

$$\mathcal{M} \models \phi(a_1, \dots, a_m)$$
 if and only if $\langle M, \in^{\mathcal{M}} \rangle$
 $\models (\langle V_{\alpha}, \in, F(\alpha) \rangle \models \phi(a_1, \dots, a_m)).$



Let $\phi(x_1, \ldots, x_m)$ be a $\Delta_0(\mathcal{L}_X)$ -formula and let $a_1, \ldots, a_m \in M$. Let $\alpha \in \text{Ord}^{\mathcal{M}}$ be such that $a_1, \ldots, a_m \in (V_{\alpha}^{\mathcal{M}})^*$. Now,

$$\mathcal{M} \models \phi(a_1, \dots, a_m)$$
 if and only if $\langle M, \in^{\mathcal{M}} \rangle \models (\langle V_{\alpha}, \in, F(\alpha) \rangle \models \phi(a_1, \dots, a_m))$ if and only if $\mathcal{N} \models (\langle V_{\alpha}, \in, F(\alpha) \rangle \models \phi(a_1, \dots, a_m))$ if and only if $\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \phi(a_1, \dots, a_m)$.

This shows that $\mathcal{M} \prec_0 \langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle$ and completes the base case of the induction. Now, let $n \in \omega$ and suppose that $\mathcal{M} \prec_n \langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle$; we prove this relation for n+1 in place of n. We leave the case n=0 as an exercise, as it requires a simpler version of the argument below. So assume n>0. Let $\phi(x_1,\ldots,x_m)$ be an \mathcal{L}_X -formula $\exists y \forall z \psi(y,z,x_1,\ldots,x_m)$ where $\psi(y,z,x_1,\ldots,x_m)$ is $\sum_{n-1}(\mathcal{L}_X)$. Let $a_1,\ldots,a_m \in \mathcal{M}$. It follows immediately from the fact that $\mathcal{M} \prec_n \langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle$ that if $\mathcal{M} \models \phi(a_1,\ldots,a_m)$, then $\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \phi(a_1,\ldots,a_m)$. Conversely, suppose that

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \phi(a_1, \dots, a_m) \text{ and } \mathcal{M} \models \neg \phi(a_1, \dots, a_m).$$

Let $\alpha \in \operatorname{Ord}^{\mathcal{M}}$ be such that

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models (\exists y \in V_{\alpha}) \forall z \psi(y, z, a_1, \dots, a_m).$$

Note that

$$\mathcal{M} \models (\forall y \in V_{\alpha}) \exists z \neg \psi(y, z, a_1, \dots, a_m).$$

Work inisde \mathcal{M} . Since $\mathcal{M} \models \operatorname{ZFC}(X)$, we can find a set C such that for all $y \in V_{\alpha}$, there exists $z \in C$ such that $\neg \psi(y, z, a_1, \ldots, a_m)$ holds. Let $f : V_{\alpha} \longrightarrow C$ be such that for all $y \in V_{\alpha}$, $\neg \psi(y, f(y), a_1, \ldots, a_m)$ holds. Working in the meta-theory again, note that

$$\mathcal{M} \models (\forall y \in V_{\alpha}) \neg \psi(y, f(y), a_1, \dots, a_m).$$

The expression $(\forall y \in V_{\alpha}) \neg \psi(y, f(y), a_1, \dots, a_m)$ is $\Pi_n(\mathcal{L}_X)$ (indeed, $\Pi_{n-1}(\mathcal{L}_X)$) with parameters $a_1, \dots, a_m, f \in M$. Therefore, by the induction hypothesis

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models (\forall y \in V_{\alpha}) \neg \psi(y, f(y), a_1, \dots, a_m).$$

Now, let $y_0 \in N$ be such that

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \forall z \psi(y_0, z, a_1, \dots, a_m) \land (y_0 \in V_{\alpha}).$$

But,

$$\langle N, \in^{\mathcal{N}}, X^{\mathcal{N}} \rangle \models \neg \psi(y_0, f(y_0), a_1, \dots, a_m),$$

which is a contradiction. This completes the induction step and proves the theorem.



We now turn to proving Theorem 6.2. For the remainder of this section fix a countable recursively saturated structure $\mathcal{M} = \langle M, \in^{\mathcal{M}} \rangle$ with $\mathcal{M} \models \mathrm{ZFC}$, and fix an H-cut $I \subseteq M$. Again, we will write \mathcal{I} for the \mathcal{L} -structure $\langle I, \in^{\mathcal{M}} \rangle$. Let $\bar{k} \in \mathrm{Ord}^{\mathcal{M}} \backslash \mathrm{Ord}^{\mathcal{I}}$ be such that $\mathcal{M} \models (\bar{k} \text{ is a regular cardinal})$. Using the construction presented in Sect. 4 we will construct $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ such that $\mathcal{M} \prec_{\mathrm{cf}} \mathcal{N}_{\mathcal{U},\mathbb{Q}}$ and there is an embedding $j \mapsto \check{j}$ of $\mathrm{Aut}(\mathbb{Q})$ into $\mathrm{Aut}(\mathcal{N}_{\mathcal{U},\mathbb{Q}})$ such that for all fixed point free $j \in \mathbb{Q}$, $\mathcal{I}_{\mathrm{fix}}(\check{j}) = \mathcal{I}$. Carefully choosing the ultrafilter \mathcal{U} will allow us to build an isomorphism between \mathcal{M} and $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ that fixes I. In order to ensure that such an isomorphism exists we need \mathcal{M} and $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ to code the same subsets of I.

Definition 6.1 Let $\mathcal{N} = \langle N, \in^{\mathcal{N}} \rangle$ be an \mathcal{L} -structure and let $K \subseteq N$. Define

$$SSy_K(\mathcal{N}) = \{c^* \cap K \mid c \in N\}.$$

Definition 6.2 A filter $\mathcal{U} \subseteq (\mathcal{P}(\bar{\kappa}))^*$ is I-conservative if for every $n \in \omega$, and for every $f \in M$, if $\mathcal{M} \models (f : H_{\bar{\kappa}} \longrightarrow \mathcal{P}([\bar{\kappa}]^{n+1}))$ then there exists $X \in \mathcal{U}$ and an I-large \mathcal{M} -cardinal $\lambda \in \bar{\kappa}^*$ such that for all $x \in (H_{\lambda}^{\mathcal{M}})^*$,

$$\mathcal{M} \models ([X]^{n+1} \subseteq f(x)) \text{ or } \mathcal{M} \models ([X]^{n+1} \subseteq [\bar{\kappa}]^{n+1} \setminus f(x)).$$

As we did in Sect. 4, let

$$\mathcal{F} = \{ f \in M \mid (\exists n \in \omega) (\mathcal{M} \models f \text{ is a function with domain } [\bar{\kappa}]^{n+1}) \}$$

Let $\mathcal{L}_{\mathcal{F}}$ be the extension of \mathcal{L} defined in Sect. 4 that adds new function symbols \hat{f} for every $f \in \mathcal{F}$. Let $\mathcal{M}_{\mathcal{F}}$ be the expansion of \mathcal{M} to an $\mathcal{L}_{\mathcal{F}}$ -structure defined in Sect. 4. Since each new function symbol \hat{f} in $\mathcal{L}_{\mathcal{F}}$ is coded by a point in \mathcal{M} and $\mathcal{M} \models \mathsf{ZFC}$, we immediately get the following extension of Lemma 4.13:

Lemma 6.4
$$\mathcal{M}_{\mathcal{F}} \models \mathcal{L}_{\mathcal{F}}$$
-separation $+ \mathcal{L}_{\mathcal{F}}$ -collection.

We will extend Theorem 4.12 to show that we find an external non-principle ultrafilter on the subsets of \bar{k} in \mathcal{M} which is simultaneously I-complete, canonically Ramsey, I-tight, I-conservative, and contains arbitrarily small I-large sets. The fact that we can extend Theorem 4.12 will follow from the following lemma:

Lemma 6.5 Let $n \in \omega$. Let $X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ be I-large and let $f \in M$ be such that $\mathcal{M} \models (f : H_{\bar{\kappa}} \longrightarrow \mathcal{P}([X]^{n+1}))$. There exists an I-large \mathcal{M} -cardinal $\lambda \in \bar{\kappa}^*$ and an I-large $Y \subseteq X$ such that for all $x \in (H_{\lambda}^{\mathcal{M}})^*$,

$$\mathcal{M} \models ([Y]^{n+1} \subseteq f(x)) \text{ or } \mathcal{M} \models ([Y]^{n+1} \subseteq [\bar{\kappa}]^{n+1} \setminus f(x)).$$

Proof Work inside \mathcal{M} . Let $\mu = |X|$. Using Lemma 4.9 and Theorem 4.6 we can find *I*-large cardinals γ , $\lambda < \mu$ such that $\gamma = |H_{\lambda^+}|$ and $\mu \to ((2^{\gamma})^+)_{2^{\gamma}}^{n+1}$. Define $F: [X]^{n+1} \longrightarrow 2^{H_{\lambda^+}}$ such that for all $A \in [X]^{n+1}$,

$$F(A) = g_A$$
 where for all $x \in H_{\lambda^+}$, $g_A(x) = 1$ if and only if $A \in f(x)$.



Since $\mu \to ((2^{\gamma})^+)_{2^{\gamma}}^{n+1}$, we can find an I-large $Y \subseteq X$ such that F is monochromatic on $[Y]^{n+1}$. Therefore, for all $A, B \in [Y]^{n+1}$ and for all $x \in H_{\lambda^+}$, either $A, B \in f(x)$ or $A, B \notin f(x)$. Therefore, for all $x \in (H_{\lambda^+}^{\mathcal{M}})^*$,

$$\mathcal{M} \models ([Y]^{n+1} \subseteq f(x)) \text{ or } \mathcal{M} \models ([Y]^{n+1} \subseteq [\bar{\kappa}]^{n+1} \setminus f(x)).$$

Theorem 6.6 There exists an n.p. ultrafilter $\mathcal{U} \subseteq (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ which is I-complete, canonically Ramsey, I-tight, I-conservative and such that $\{|X|^{\mathcal{M}} \mid X \in \mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \setminus \mathrm{Ord}^{\mathcal{I}}$.

Proof We use exactly the same method as we used to prove Theorem 4.12. Let $\langle f_n \mid n \in \omega \rangle$ be an enumeration of \mathcal{F} and let $\langle k_n \mid n \in \omega \rangle$ be a sequence of natural numbers such that for all $n \in \omega$, $\mathcal{M} \models (f_n \text{ is a function with domain } [\bar{k}]^{k_n})$. Let $\langle \lambda_n \mid n \in \omega \rangle$ be a decreasing sequence of \mathcal{M} -cardinals that is downward cofinal in $\text{Ord}^{\mathcal{M}} \setminus \text{Ord}^{\mathcal{I}}$ with $\lambda_0 \in \bar{k}^*$. Let

$$\mathcal{G} = \left\{ g \in M \mid (\exists n \in \omega) \left(\mathcal{M} \models g : H_{\bar{\kappa}} \longrightarrow \mathcal{P}\left([\bar{\kappa}]^{n+1} \right) \right) \right\}.$$

Let $\langle g_n \mid n \in \omega \rangle$ be an enumeration of \mathcal{G} and let $\langle l_n \mid n \in \omega \rangle$ be a sequence of natural numbers such that for all $n \in \omega$, $\mathcal{M} \models (g_n : H_{\bar{\kappa}} \longrightarrow \mathcal{P}([\bar{\kappa}]^{l_n}))$. Using Lemmas 4.4, 4.10, 4.11 and 6.5 inductively build sequences $\langle W_n \mid n \in \omega \rangle$, $\langle Q_n \mid n \in \omega \rangle$, $\langle X_n \mid n \in \omega \rangle$, $\langle Y_n \mid n \in \omega \rangle$ and $\langle Z_n \mid n \in \omega \rangle$ of I-large elements of $(\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^*$ such that for all $n \in \omega$,

- 1. $\mathcal{M} \models (W_n \supseteq Q_n \supseteq X_n \supseteq Y_n \supseteq Z_n \supseteq W_{n+1}),$
- 2. $\mathcal{M} \models (W_n \text{ is } f_n\text{-canonical}),$
- 3. there exists an *I*-large \mathcal{M} -cardinal $\lambda \in \bar{\kappa}^*$ such that for all $x \in (H_{\lambda^+}^{\mathcal{M}})^*$,

$$\mathcal{M} \models ([Q_n]^{l_n} \subseteq g_n(x)) \text{ or } \mathcal{M} \models ([Q_n]^{l_n} \subseteq [\bar{k}]^{l_n} \setminus g_n(x)),$$

- 4. if $k_n = 1$ and there is an \mathcal{M} -cardinal $\mu \in I$ such that $\mathcal{M} \models (f_n : \bar{k} \longrightarrow H_{\mu})$ then f_n is constant on X_n , otherwise $X_n = Q_n$,
- 5. $\mathcal{M} \models (f_n \text{ is constant on } [Y_n]^{k_n}) \text{ or there is an } I\text{-large }\mathcal{M}\text{-cardinal }\mu \in \bar{\kappa}^* \text{ such that}$

$$\mathcal{M} \models \left(\forall A \in [Y_n]^{k_n} \right) (f_n(A) \notin H_\mu),$$

6. $\mathcal{M} \models (|Z_n| < \lambda_n)$.

Define $\mathcal{U} = \{X \in (\mathcal{P}(\bar{\kappa})^{\mathcal{M}})^* \mid (\exists n \in \omega)(\mathcal{M} \models W_n \subseteq X)\}$. It is clear from the construction that \mathcal{U} is an n.p. ultrafilter that is *I*-complete, canonically Ramsey, *I*-tight, *I*-conservative and is such that $\{|X|^{\mathcal{M}} \mid X \in \mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}} \setminus \mathrm{Ord}^{\mathcal{I}}$.



Let $\mathcal{U}\subseteq (\mathcal{P}(\bar{k})^{\mathcal{M}})^*$ be an n.p. ultrafilter obtained from Theorem 6.6, so \mathcal{U} is I-complete, canonically Ramsey, I-tight, I-conservative and $\{|X|^{\mathcal{M}}\mid X\in\mathcal{U}\}$ is downward cofinal in $\mathrm{Ord}^{\mathcal{M}}\setminus\mathrm{Ord}^{\mathcal{I}}$. Let $\mathcal{N}_{\mathcal{U},\mathbb{Q}}=\langle N_{\mathcal{U},\mathbb{Q}},\in^{\mathcal{N}}\rangle$ be the iterated ultrapower, constructed in Sect. 4, of \mathcal{M} and \mathcal{U} endowed with a class of indiscernibles of order-type \mathbb{Q} . The results proved in Sect. 4 and the fact $\mathcal{M}\models\mathrm{ZFC}$ imply that $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ has the following properties:

- 1. $\mathcal{M} \prec_{\mathrm{cf}} \mathcal{N}_{\mathcal{U},\mathbb{O}}$,
- 2. $|N_{\mathcal{U},\mathbb{O}}| = \aleph_0$,
- 3. there is an embedding $j \mapsto \tilde{j}$ of $\operatorname{Aut}(\mathbb{Q})$ into $\operatorname{Aut}(\mathcal{N}_{\mathcal{U},\mathbb{Q}})$ such that if $j \in \operatorname{Aut}(\mathbb{Q})$ has no fixed points then $\mathcal{I}_{\operatorname{fix}}(\tilde{j}) = \mathcal{I}$.

We will show below (Theorem 6.11) that (3) holds for all non-trivial $j \in \operatorname{Aut}(\mathbb{Q})$, not just the $j \in \operatorname{Aut}(\mathbb{Q})$ with no fixed points. First, however, we show that $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ can be identified with \mathcal{M} by an isomorphism that fixes I. The existence of this isomorphism will follow from Theorem 6.9. The fact that \mathcal{U} is I-conservative ensures that $\operatorname{SSy}_I(\mathcal{N}_{\mathcal{U},\mathbb{Q}}) = \operatorname{SSy}_I(\mathcal{M})$.

Lemma 6.7 $SSy_I(\mathcal{N}_{\mathcal{U},\mathbb{O}}) = SSy_I(\mathcal{M}).$

Proof Since $\mathcal{I} \subseteq_{\text{end}} \mathcal{N}_{\mathcal{U},\mathbb{O}}$, it follows that

$$SSy_I(\mathcal{N}_{\mathcal{U},\mathbb{O}}) \supseteq SSy_I(\mathcal{M}).$$

We need to show the reverse inclusion. Let $[\hat{f}(c_{i_0},\ldots,c_{i_n})] \in N_{\mathcal{U},\mathbb{Q}}$ where $f \in \mathcal{F}$ and $i_0 < \cdots < i_n \in \mathbb{Q}$. Let

$$A = \left\{ x \in I \mid \mathcal{N}_{\mathcal{U}, \mathbb{Q}} \models \left(x \in \left[\hat{f} \left(c_{i_0}, \dots, c_{i_n} \right) \right] \right) \right\}.$$

So $A \in SSy_I(\mathcal{N}_{\mathcal{U},\mathbb{Q}})$ and we need to show that $A \in SSy_I(\mathcal{M})$. Working inside $\mathcal{M}_{\mathcal{F}}$ define $g: H_{\bar{k}} \longrightarrow \mathcal{P}([\bar{k}]^{n+1})$ such that for all $x \in H_{\bar{k}}$,

$$g(x) = \left\{ \{\alpha_0 < \dots < \alpha_n\} \in \left[\bar{\kappa}\right]^{n+1} \mid x \in \hat{f}(\alpha_0, \dots, \alpha_n) \right\}.$$

So, $g \in M$. From Lemma 4.17 we have for all $x \in H_{\bar{\kappa}}^*$,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(x \in \left[\hat{f}\left(c_{i_0},\ldots,c_{i_n}\right)\right]\right) \text{ if and only if there exists } Y \in \mathcal{U} \text{ s.t. } \mathcal{M}_{\mathcal{F}} \right)$$

 $\models ([Y]^{n+1} \subseteq g(x)).$

Let $X \in \mathcal{U}$ and let $\lambda \in \bar{\kappa}^*$ be an *I*-large \mathcal{M} -cardinal such that for all $x \in (H_{\lambda^+}^{\mathcal{M}})^*$,

$$\mathcal{M}_{\mathcal{F}} \models ([X]^{n+1} \subseteq g(x)) \text{ or } \mathcal{M}_{\mathcal{F}} \models ([X]^{n+1} \subseteq [\bar{\kappa}]^{n+1} \setminus g(x)).$$

Working inside $\mathcal{M}_{\mathcal{F}}$, let

$$C = \left\{ x \in H_{\bar{\kappa}} \mid [X]^{n+1} \subseteq g(x) \right\}.$$

So, $C \in M$. Since $I \subseteq (H_{\lambda^+}^{\mathcal{M}})^*$, it follows that for all $x \in I$,

$$x \in C^*$$
 if and only if $\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(x \in \left[\hat{f}\left(c_{i_0},\ldots,c_{i_n}\right)\right]\right)$.

Therefore $A = C^* \cap I$ and $A \in SSy_I(\mathcal{M})$.

The next result shows that for any set in $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ there is a set in \mathcal{M} with the same \mathcal{M} -members on some initial segment of \mathcal{M} that contains I.

Lemma 6.8 Let $\lambda \in M$ be an I-large \mathcal{M} -cardinal. If $u \in (H_{\lambda}^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*$ then there is an I-large \mathcal{M} -cardinal $\mu \in M$ and $w \in M$ such that:

$$\left\{ x \in M \mid (\mathcal{M} \models (x \in H_{\mu})) \land \left(\mathcal{N}_{\mathcal{U}, \mathbb{Q}} \models \left(\left[\hat{h}_{x}(c_{0}) \right] \in u \right) \right) \right\}$$

$$= \left\{ x \in M \mid \mathcal{M} \models (x \in w) \right\}.$$
(6)

Proof Let $u \in (H_{\lambda}^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*$. Working inside $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$, define $f: \operatorname{Card} \cap \lambda \longrightarrow H_{\lambda}$ such that for all $\gamma \in \operatorname{Card} \cap \lambda$,

$$f(\gamma) = u \cap H_{\gamma}$$
.

Since $SSy_I(\mathcal{N}_{\mathcal{U},\mathbb{Q}}) = SSy_I(\mathcal{M})$, there is $g \in M$ such that g agrees with f on I. Work inside $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$. Define

$$C = \{ \delta \in \lambda \mid (\delta \in \text{Card}) \land (\forall \gamma \in \delta^+) ((\gamma \in \text{Card}) \\ \longrightarrow ((g \cap (2^{\gamma} \times H_{2^{\gamma}}) \text{ is a function}) \land (f(\gamma) = g(\gamma)))) \}.$$

Since I is a H-cut of $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$, there is an I-large cardinal $\eta \in C$. Using Lemma 4.24 we can find $\mu \in \operatorname{Card}^{\mathcal{M}} \backslash \operatorname{Card}^{\mathcal{I}}$ such that $\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models ([\hat{h}_{\mu}] \leq \eta)$. Since $\mathcal{M} \prec \mathcal{N}_{\mathcal{U},\mathbb{Q}}$,

$$\mathcal{M} \models (g \cap (2^{\mu} \times H_{2^{\mu}}) \text{ is a function}).$$

Let $w \in M$ be such that $\mathcal{M} \models (g(\mu) = w)$. Therefore

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(g\left(\left[\hat{h}_{\mu}(c_0) \right] \right) = \left[\hat{h}_{w}(c_0) \right] \right) \land \left(\left[\hat{h}_{w}(c_0) \right] = u \cap H_{\left[\hat{h}_{\mu}(c_0) \right]} \right).$$

And so $w, \mu \in M$ satisfy (6).

We are now in a position to show that \mathcal{M} and $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ can be identified by an isomorphism that fixes I.

Theorem 6.9 There exists an isomorphism $\Theta : \mathcal{M} \longrightarrow \mathcal{N}_{\mathcal{U},\mathbb{Q}}$ such that for all $x \in I$, $\Theta(x) = x$.



Proof We will construct Θ using a back-and-forth construction. We begin by endowing \mathcal{M} and $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ with satisfaction classes. By Theorem 2.7 there is an $X^{\mathcal{M}} \subseteq M$ such that

- (I) $\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle \models \mathsf{ZFC}(X),$
- (II) $X^{\mathcal{M}}$ is a satisfaction class for \mathcal{M} .

Since \mathcal{M} is recursively saturated, there is a non-standard $s \in (\omega^{\mathcal{M}})^* (= (\omega^{\mathcal{I}})^* = (\omega^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*)$ such that $X^{\mathcal{M}}$ is s-correct for \mathcal{M} . Throughout this proof we will identify formulae from the point of view of \mathcal{M} with their Gödel codes in $\omega^{\mathcal{M}}$. We will also abbrieviate our notation by identifying elements of M with the equivalence classes of their corresponding constant functions in $N_{\mathcal{U},\mathbb{Q}}$. By shortening $X^{\mathcal{M}}$ if necessary, we can assume without loss of generality that $\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$ satisfies:

$$\langle \phi, a \rangle \in X \Rightarrow \lceil \phi \rceil < s. \tag{7}$$

Using Theorem 6.3 we can expand $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ to an \mathcal{L}_X -structure $\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$ such that

$$\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle \prec \langle N, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle.$$

It immediately follows that $X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}$ is a satisfaction class that is s-correct for $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ and $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ is recursively saturated. Moreover, $\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$ satisfies (7). Let $\langle p_i \mid i \in \omega \rangle$ be an enumeration of M and let $\langle q_i \mid i \in \omega \rangle$ be an enumeration of $N_{\mathcal{U},\mathbb{Q}}$. We will construct $\Theta: \mathcal{M} \longrightarrow \mathcal{N}_{\mathcal{U},\mathbb{Q}}$ by constructing sequences $\langle u_i \mid i \in \omega \rangle$ and $\langle v_i \mid i \in \omega \rangle$, together with decreasing sequences $\langle \gamma_i \mid i \in \omega \rangle$ and $\langle r_i \mid i \in \omega \rangle$, such that

- (I) $\langle u_i \mid i \in \omega \rangle$ enumerates M,
- (II) $\langle v_i \mid i \in \omega \rangle$ enumerates N,
- (III) for all $i \in \omega$, γ_i is an *I*-large \mathcal{M} -cardinal,
- (IV) for all $i \in \omega$, $r_i \in (\omega^{\mathcal{M}})^*$ is non-standard.

We then define

$$\Theta(u_i) = v_i \text{ for all } i \in \omega.$$

We begin by letting $\gamma_0 \in \bar{\kappa}^*$ be an *I*-large \mathcal{M} -cardinal and $r_0 = s$. The fact that $\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle \prec \langle N, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$ ensures that if $\phi(y) \in (\text{Form}^{\mathcal{M}})^*$ with $\phi(y) < r_0$ then for all $a \in (H_{\gamma_0}^{\mathcal{M}})^*$,

$$\langle \phi(y), \langle a \rangle \rangle \in X^{\mathcal{M}}$$
 if and only if $\langle \phi(y), \langle a \rangle \rangle \in X^{\mathcal{N}_{\mathcal{U}, \mathbb{Q}}}$. (8)

At stage j>0, after having defined $u_0,\ldots,u_{j-1}\in M,\,v_0,\ldots,v_{j-1}\in N_{\mathcal{U},\mathbb{Q}},$ $r_j\in(\omega^{\mathcal{M}})^*$ and $\gamma_j\in M$ we will ensure that the following condition is maintained:



 (\dagger_j) if $\phi(x_0,\ldots,x_{j-1},y) \in (\text{Form}^{\mathcal{M}})^*$ with $\phi(x_0,\ldots,x_{j-1},y) < r_j$ then for all $a \in (H^{\mathcal{M}}_{\gamma_j})^*$,

$$\langle \phi(x_0,\ldots,x_{j-1},y), \langle u_0,\ldots,u_{j-1},a\rangle \rangle \in X^{\mathcal{M}}$$
 if and only if $\langle \phi(x_0,\ldots,x_{j-1},y), \langle v_0,\ldots,v_{j-1},a\rangle \rangle \in X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}$.

We construct the sequences $\langle u_i \mid i \in \omega \rangle$, $\langle v_i \mid i \in \omega \rangle$, $\langle \gamma_i \mid i \in \omega \rangle$ and $\langle r_i \mid i \in \omega \rangle$ by induction. Note that (8) corresponds to the condition \dagger_0 that forms the base case of our induction before any of the us or vs have been defined. Suppose that we have defined $u_0, \ldots, u_{2k-1} \in M$ and $v_0, \ldots, v_{2k-1} \in N_{\mathcal{U}, \mathbb{Q}}, r_{2k} \in (\omega^{\mathcal{M}})^*$ and γ_{2k} such that \dagger_{2k} holds.

STAGE 2k + 1: We need to choose $u_{2k} \in M$, $v_{2k} \in N$, $r_{2k+1} \in (\omega^{\mathcal{M}})^*$ and γ_{2k+1} such that \dagger_{2k+1} is maintained. Let $v_{2k} = q_k$. This choice will eventually ensure that Θ is onto $N_{\mathcal{U},\mathbb{Q}}$. Using Lemma 4.9, let λ be an I-large \mathcal{M} -cardinal such that $2^{\lambda} < \gamma_{2k}$. Working inside $\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$, define

$$u = \{ \langle \phi(x_0, \dots, x_{2k}, y), a \rangle \mid (\phi < r_{2k}) \land (a \in H_{\lambda}) \land (\langle \phi, \langle v_0, \dots, v_{2k}, a \rangle) \in X \} \}.$$

Note that $u \in (H_{2^{\lambda}}^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*$. The fact that $X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}$ is s-correct ensures that for all $l \in \omega$,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \exists x \left((\forall y \in H_{\lambda}) \left(\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k}, y), y \rangle \in u) \iff \phi(v_0, \dots, v_{2k-1}, x, y) \right) \right).$$

Using Lemma 6.8 we can find an *I*-large \mathcal{M} -cardinal $\gamma_{2k+1} \in M$ with $\gamma_{2k+1} \leq \lambda$ and $w \in M$ such that

$$\left\{x \in M \mid \left(\mathcal{M} \models \left(x \in H_{\gamma_{2k+1}}\right)\right) \land \left(\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{h}_x(c_0)\right] \in u\right)\right)\right\}$$
$$= \left\{x \in M \mid \mathcal{M} \models (x \in w)\right\}.$$

Therefore, for all $l \in \omega$,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \exists x \left(\left(\forall y \in H_{\gamma_{2k+1}} \right) \left(\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k}, y), y \rangle \in w \iff \phi(v_0, \dots, v_{2k-1}, x, y)) \right) \right).$$

Since $2^{\gamma_{2k+1}} < \gamma_{2k}$ and $w, \gamma_{2k+1} \in (H_{2^{\gamma_{2k}}}^{\mathcal{M}})^*, \dagger_{2k}$ implies that for all $l \in \omega$,

$$\mathcal{M} \models \exists x \left((\forall y \in H_{\gamma_{2k+1}}) \left(\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k}, y), y \rangle \in w \iff \phi(u_0, \dots, u_{2k-1}, x, y)) \right) \right).$$

For all $l \in \omega$, let $\psi_l(x, y, x_0, \dots, x_{2k-1}, z)$ be the \mathcal{L} -formula:

$$\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k}, y), y \rangle \in z \iff \phi(x_0, \dots, x_{2k-1}, x, y)).$$



And let $\Psi(l, u_0, \dots, u_{2k-1}, w, \gamma_{2k+1})$ be the \mathcal{L}_X -formula:

$$\exists x \left(\left(\forall y \in H_{\gamma_{2k+1}} \right) \left(\left\langle \psi_l(x, y, x_0, \dots, x_{2k-1}, z), \left\langle x, y, u_0, \dots, u_{2k-1}, w \right\rangle \right) \in X \right) \right).$$

Therefore, for all $l \in \omega$,

$$\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle \models \Psi(l, u_0, \dots, u_{2k-1}, w, \gamma_{2k+1}).$$

So, by overspill in $\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$, we can find a non-standard $r_{2k+1} \in (\omega^{\mathcal{M}})^*$ with $r_{2k+1} < r_{2k}$ such that

$$\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$$

$$\models \exists x ((\forall y \in H_{\gamma_{2k+1}})(\langle \psi_{r_{2k+1}}(x, y, x_0, \dots, x_{2k-1}, z), \langle x, y, u_0, \dots, u_{2k-1}, w \rangle) \in X)).$$

Let $u_{2k} \in M$ be such that

$$\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$$

 $\models (\forall y \in H_{\gamma_{2k+1}})(\langle \psi_{r_{2k+1}}(x, y, x_0, \dots, x_{2k-1}, z), \langle u_{2k}, y, u_0, \dots, u_{2k-1}, w \rangle) \in X).$

Therefore the choices of v_{2k} , u_{2k} , r_{2k+1} and γ_{2k+1} made in this stage maintain \dagger_{2k+1} .

STAGE 2k + 2: We need to choose $u_{2k+1} \in M$, $v_{2k+1} \in N_{\mathcal{U},\mathbb{Q}}$, $r_{2k+2} \in (\omega^{\mathcal{M}})^*$ and γ_{2k+2} such that \dagger_{2k+2} is maintained. Let $u_{2k+1} = p_k$. This choice will eventually ensure that Θ is defined on all of M. Using Lemma 4.9, let $\gamma_{2k+2} \in M$ be an I-large \mathcal{M} -cardinal such that $2^{\gamma_{2k+2}} < \gamma_{2k+1}$. Working inside $\langle M, \in^{\mathcal{M}}, X^{\mathcal{M}} \rangle$, define

$$u = \{ \langle \phi(x_0, \dots, x_{2k+1}, y), a \rangle \mid (\phi < r_{2k+1}) \land (a \in H_{\gamma_{2k+2}}) \land (\langle \phi, \langle u_0, \dots, u_{2k+1}, a \rangle) \in X \} \}.$$

Note that $u \in (H_{2^{1/2}k+2}^{\mathcal{M}})^*$. The fact that $X^{\mathcal{M}}$ is s-correct ensures that for all $l \in \omega$,

$$\mathcal{M} \models \exists x \left((\forall y \in H_{\gamma_{2k+2}}) \left(\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k+1}, y), y \rangle \in u) \iff \phi(u_0, \dots, u_{2k}, x, y) \right) \right).$$

Since $u, \gamma_{2k+2} \in (H^{\mathcal{M}}_{\gamma_{2k+1}})^*, \dagger_{2k+1}$ implies that

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \exists x \left((\forall y \in H_{\gamma_{2k+2}}) \left(\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k+1}, y), y \rangle \in u) \iff \phi(v_0, \dots, v_{2k}, x, y) \right) \right).$$

For all $l \in \omega$, let $\psi_l(x, y, x_0, \dots, x_{2k}, z)$ be the \mathcal{L} -formula:

$$\bigwedge_{\phi < l} (\langle \phi(x_0, \dots, x_{2k+1}, y), y \rangle \in z \iff \phi(x_0, \dots, x_{2k}, x, y)).$$



And let $\Psi(l, v_0, \dots, v_{2k}, u, \gamma_{2k+2})$ be the \mathcal{L}_X -formula:

$$\exists x ((\forall y \in H_{\gamma_{2k+2}})(\langle \psi_l(x, y, x_0, \dots, x_{2k}, z), \langle x, y, v_0, \dots, v_{2k}, u \rangle) \in X)).$$

Therefore, for all $l \in \omega$,

$$\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle \models \Psi(l, v_0, \dots, v_{2k}, u, \gamma_{2k+2}).$$

So, by overspill in $\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$, we can find a non-standard $r_{2k+2} \in (\omega^{\mathcal{M}})^*$ with $r_{2k+2} < r_{2k+1}$ such that

$$\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$$

$$\models \exists x ((\forall y \in H_{\gamma_{2k+2}})(\langle \psi_{r_{2k+2}}(x, y, x_0, \dots, x_{2k}, z), \langle x, y, v_0, \dots, v_{2k}, u \rangle) \in X)).$$

Let $v_{2k+1} \in N_{\mathcal{U},\mathbb{O}}$ be such that

$$\langle N_{\mathcal{U},\mathbb{Q}}, \in^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}}, X^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}} \rangle$$

$$\models (\forall y \in H_{\gamma_{2k+2}})(\langle \psi_{r_{2k+2}}(x, y, x_0, \dots, x_{2k-1}, z), \langle v_{2k+1}, y, v_0, \dots, v_{2k}, u \rangle) \in X).$$

Therefore the choices of v_{2k+1} , u_{2k+1} , r_{2k+2} and γ_{2k+2} made in this stage maintain \dagger_{2k+2} .

This completes the proof of the theorem.

Before concluding the proof of Theorem 6.2 we first show that the fact that \mathcal{U} is *I*-conservative can be used to demonstrate that Theorem 4.30 can be strengthened to show that *for every non-trivial* $j \in \operatorname{Aut}(\mathbb{Q})$, $\mathcal{I}_{\operatorname{fix}}(\tilde{j}) = \mathcal{I}$.

Lemma 6.10 If $[\tau] \in N_{\mathcal{U},\mathbb{Q}} \setminus I$ then there exists $x \in M \setminus I$ with

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\hat{h}_x(c_0) \in TC(\{[\tau]\})\right).$$

Proof Let $[\tau] \in N_{\mathcal{U},\mathbb{Q}} \setminus I$. Let $f \in \mathcal{F}$ and $i_1 < \cdots < i_n \in \mathbb{Q}$ be such that

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{f}\left(c_{i_1},\ldots,c_{i_n}\right) \right] = TC(\{[\tau]\}) \right).$$

Let $\lambda \in M$ be an *I*-large \mathcal{M} -cardinal such that

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{f}\left(c_{i_1}, \ldots, c_{i_n} \right) \right] \in H_{\left[\hat{h}_{\lambda}\left(c_0 \right) \right]} \right).$$

Using Lemma 6.8 we can find an *I*-large \mathcal{M} -cardinal $\mu \in M$ and $w \in M$ such that

$$\left\{x \in M \mid (\mathcal{M} \models (x \in H_{\mu})) \land \left(\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{h}_{x}(c_{0})\right] \in \left[\hat{f}\left(c_{i_{1}},\ldots,c_{i_{n}}\right)\right]\right)\right)\right\}$$

$$= \left\{x \in M \mid \mathcal{M} \models (x \in w)\right\}.$$



Working inside \mathcal{M} , define

$$C = \{ \gamma \in \mu \mid (\gamma \in \text{Card}) \land (w \cap (H_{\gamma^+} \backslash H_{\gamma}) \neq \emptyset) \}.$$

We will prove that $C \notin I$. Suppose that $C \in I$. Work inside \mathcal{M} . Let $\eta' = \sup C$ and let $\eta = (\eta')^+$. Therefore $\eta \in I$ is a cardinal and for all cardinals $\eta \leq \gamma < \mu$, $w \cap (H_{\gamma^+} \backslash H_{\gamma}) = \emptyset$. Let $a \in w \backslash H_{\eta}$ be a set with minimal rank. Therefore, for all $y \in a$, $y \in H_{\eta}$. But $\mathcal{P}(H_{\eta}) \subseteq H_{2^{2^{\eta}}}$, and so there exists a cardinal $\eta \leq \gamma < 2^{2^{\eta}}$ such that $a \in H_{\gamma^+} \backslash H_{\gamma}$. This contradicts the fact that $w \cap (H_{\gamma^+} \backslash H_{\gamma}) = \emptyset$ and shows that $C \notin I$. Therefore, since $I \subseteq M$ is a H-cut, there exists an I-large \mathcal{M} -cardinal $\xi \in M$ with $\xi \in \mu^*$ such that $w \cap (H_{\xi^+} \backslash H_{\xi}) \neq \emptyset$. Let $x \in (w \backslash (H_{\xi^+} \backslash H_{\xi}))$. Therefore

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\hat{h}_{x}(c_{0}) \in \left[\hat{f}\left(c_{i_{1}},\ldots,c_{i_{n}}\right)\right]\right).$$

Theorem 6.11 If $j \in Aut(\mathbb{Q})$ is non-trivial then $\mathcal{I}_{fix}(\tilde{j}) = \mathcal{I}$.

Proof Let $j \in \text{Aut}(\mathbb{Q})$ be non-trivial. The embedding of \mathcal{I} into an initial segment of $\mathcal{N}_{\mathcal{U},\mathbb{Q}}$ shows that $I \subseteq I_{\text{fix}}(\tilde{j})$. Therefore, we need to show that for every $[\tau] \in \mathcal{N}_{\mathcal{U},\mathbb{Q}} \setminus I$, there exists $y \in (\text{TC}(\{[\tau]\})^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*$ such that $\tilde{j}(y) \neq y$. Let $l, m \in \mathbb{Q}$ such that $j(l) = m \neq l$. Let $[\tau] \in \mathcal{N}_{\mathcal{U},\mathbb{Q}} \setminus I$. Using Lemma 6.10, let $x \in M \setminus I$ be such that

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\hat{h}_x(c_0) \in TC(\{[\tau]\})\right).$$

Therefore, by Lemma 4.29,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\hat{h}_{TC(x)}(c_0) \subseteq TC(\{[\tau]\})\right).$$

And, since $x \in M \setminus I$, $|TC(x)|^{\mathcal{M}} \notin I$. Therefore, there is $X \in \mathcal{U}$ with $\mathcal{M} \models (|X| \le |TC(x)|)$. Let $g \in M$ be such that

$$\mathcal{M} \models (g : X \longrightarrow TC(x)) \land (g \text{ is injective}).$$

Therefore, there is $f \in M$ such that $\mathcal{M} \models (f : \bar{k} \longrightarrow TC(x))$ and for all $z_1 < z_2 \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models \hat{f}(z_1) \neq \hat{f}(z_2).$$

So, by Lemma 4.17, $\tilde{j}([\hat{f}(c_l)]) \neq [\hat{f}(c_l)]$. Moreover, for all $z \in X^*$,

$$\mathcal{M}_{\mathcal{F}} \models (\hat{f}(z) \in TC(x)).$$

Therefore, by Lemma 4.17,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{f}(c_l) \right] \in \left[\hat{h}_{TC(x)}(c_0) \right] \right).$$



And so,

$$\mathcal{N}_{\mathcal{U},\mathbb{Q}} \models \left(\left[\hat{f}(c_l) \right] \in TC(\{[\tau]\}) \right),$$

which proves the theorem.

Let $\Theta: \mathcal{M} \longrightarrow \mathcal{N}_{\mathcal{U},\mathbb{Q}}$ be the isomorphism obtained from Theorem 6.9 such that for all $x \in I$, $\Theta(x) = x$. To complete the proof of Theorem 6.2 we use Θ to turn the embedding $j \mapsto \tilde{j}$ of $\operatorname{Aut}(\mathbb{Q})$ into $\operatorname{Aut}(\mathcal{N}_{\mathcal{U},\mathbb{Q}})$, into an embedding $j \mapsto \tilde{j}$ of $\operatorname{Aut}(\mathbb{Q})$ into $\operatorname{Aut}(\mathcal{M})$ such that if $j \in \operatorname{Aut}(\mathbb{Q})$ is non-trivial then $\mathcal{I}_{\operatorname{fix}}(\check{j}) = \mathcal{I}$. For all $j \in \operatorname{Aut}(\mathbb{Q})$, define $\check{j}: \mathcal{M} \longrightarrow \mathcal{M}$ such that

$$\check{j}(x) = \Theta^{-1}(\tilde{j}(\Theta(x)))$$
 for all $x \in M$.

It follows immediately from the fact that Θ is isomorphism that that the map $j \mapsto \check{j}$ is an injective group homomorphism of $\operatorname{Aut}(\mathbb{Q})$ into $\operatorname{Aut}(\mathcal{M})$.

Theorem 6.12 If $j \in Aut(\mathbb{Q})$ is non-trivial then $\mathcal{I}_{fix}(\check{j}) = \mathcal{I}$.

Proof Let $j \in \operatorname{Aut}(\mathbb{Q})$ be non-trivial. The fact that Θ fixes I immediately implies that $I \subseteq I_{\operatorname{fix}}(\check{j})$. We need to show that if $x \in M \setminus I$ then there exists $y \in (\operatorname{TC}(\{x\})^{\mathcal{M}})^*$ such that $\check{j}(y) \neq y$. Let $x \in M \setminus I$. Therefore $\Theta(x) \in N_{\mathcal{U},\mathbb{Q}} \setminus I$. By Theorem 6.11 there is $y \in (\operatorname{TC}(\{\Theta(x)\})^{\mathcal{N}_{\mathcal{U},\mathbb{Q}}})^*$ such that $\check{j}(y) \neq y$. Now, $\Theta^{-1}(y) \in (\operatorname{TC}(\{x\})^{\mathcal{M}})^*$ and

$$\check{j}(\Theta^{-1}(y)) = \Theta^{-1}(\tilde{j}(\Theta(\Theta^{-1}(y)))) = \Theta^{-1}(\tilde{j}(y)) \neq \Theta^{-1}(y).$$

This completes the proof of Theorem 6.2.

Open Access This article is distributed under the terms of the Creative Commons Attribution 4.0 International License (http://creativecommons.org/licenses/by/4.0/), which permits unrestricted use, distribution, and reproduction in any medium, provided you give appropriate credit to the original author(s) and the source, provide a link to the Creative Commons license, and indicate if changes were made.

References

- 1. Baumgartner, J.E.: Canonical partition relations. J. Symb. Log. 40(4), 541-554 (1975)
- Chang, C.C., Keisler, H.J.: Model Theory. Studies in Logic and the Foundations of Mathematics, vol. 73, 3rd edn. North-Holland, Amsterdam (1990)
- Enayat, A.: Automorphisms, Mahlo cardinals, and NFU. In: Enayat, A., Kossak, R. (eds.) Nonstandard Models of Arithmetic and Set Theory. Contemporary Mathematics, vol. 361. American Mathematical Society, pp. 37–59 (2004)
- 4. Enayat, A.: Automorphisms of models of bounded arithmetic, Fundam, Math. 192, 37-65 (2006)
- 5. Erdős, P., Rado, R.: A combinatorial theorem. J. Lond. Math. Soc. 25, 249–250 (1950)
- 6. Erdős, P., Rado, R.: A partition calculus in set theory. Bull. Am. Math. Soc. 62, 427–489 (1956)
- 7. Forster, T., Kaye, R.: End-extensions preserving power set. J. Symb. Log. 56(1), 323–328 (1991)



- Friedman, H.M.: Countable models of set theories. In: Mathias, A.R.D., Rogers Jr., H. (eds.) Cambridge Summer School in Mathematical Logic, August 1–21, 1971. Springer Lecture Notes in Mathematics, vol. 337. Spinger, Berlin, pp. 539–573 (1973)
- 9. Holmes, M.R.: Strong axioms of infinity in NFU. J. Symb. Log. 66, 87-116 (2001)
- Hutchinson, J.E.: Elementary extensions of countable models of set theory. J. Symb. Log. 41(1), 139– 145 (1976)
- Jensen, R.B.: On the consistency of a slight (?) modification of Quine's new foundations. Synthese 19, 250–263 (1969)
- Kaufmann, M.: Blunt and topless extensions of models of set theory. J. Symb. Log. 48(4), 1053–1073 (1983)
- Kaye, R.: Models of Arithmetic. Oxford Logic Guides. Oxford Science Publications, vol. 15. The Clarendon Press, Oxford (1991)
- Kossak, R., Schmerl, J.H.: The Structure of Models of Peano Arithmetic. Oxford Logic Guides. Oxford Science Publications, vol. 50. The Clarendon Press, Oxford (2006)
- 15. Kotlarski, H.: On cofinal extensions of models of arithmetic. J. Symb. Log. 48(2), 253-262 (1983)
- McKenzie, Z.: Automorphisms of models of set theory and extensions of NFU. Ann. Pure Appl. Log. 166(5), 601–638 (2015)
- 17. Mathias, A.R.D.: The strength of Mac Lane set theory. Ann. Pure Appl. Log. 110(1-3), 107-234 (2001)
- 18. Mitchell, W.: Boolean topoi and the theory of sets. J. Pure Appl. Algebra 2, 261–274 (1972)
- Paris, J., Mills, G.: Closure properties of countable nonstandard integers. Fundam. Math. 103(3), 205–215 (1979)
- 20. Schlipf, J.S.: Toward model theory through recursive saturation. J. Symb. Log. 43(2), 183–206 (1978)
- Schmerl, J.H.: Recursively saturated rather classless models of Peano arithmetic. Logic Year 1979-80.
 Lecture Notes in Mathematics, vol. 859. Springer, Berlin and New York, pp. 268–282 (1981)
- Schmerl, J.H.: Recursively saturated models generated by indiscernibles. Notre Dame J. Form. Log. 26(2), 99–105 (1985)
- Smoryński, C.: Back and forth inside a recursively saturated model of arithmetic. Logic Colloquium '80 (Prague, 1980). Studies in Logic and the Foundations of Mathematics, vol. 108. North-Holland, Amsterdam-New York, pp. 273–278 (1982)
- 24. Solovay, R.: The consistency strength of NFUB. arXiv:math/9707207 [math.LO] (preprint)
- Togha, A.A.: Automorphisms of countable recursively saturated models of set theory. In: Enayat, A., R. Kossak, R. (eds.) Nonstandard Models of Arithmetic and Set Theory. Contemporary Mathematics, vol. 361. American Mathematical Society, pp. 163–167 (2004)

