

Language Dependence in Philosophy of Science and Formal Epistemology

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- Suppose we have two false hypotheses H_1 and H_2 .
- Sometimes, we would like to be able to say that H_1 is *closer to the truth* than H_2 (e.g., Newton's hypothesis vs. Ptolemy's).
- Various accounts of “closeness to the truth” (*verisimilitude*) have been proposed in the literature (see [12], [10], [11]).
- One of the many accounts is a naive, syntactical explication, which “counts numbers of true conjuncts” ([15], [16]).
- **Problematic Example.** Suppose the truth about the weather is that it is *Hot* and *Rainy* and *Windy* ($T = H \& R \& W$).
- And, consider the following two false hypotheses:
 $(H_1) \sim H \& R \& W.$
 $(H_2) \sim H \& \sim R \& \sim W.$
- If we “count numbers of true conjuncts”, then we get the verdict that H_1 is closer to the truth than H_2 ($H_1 < H_2$).
- A problem with this account is that the orderings it imposes on hypotheses are *language dependent* [9]. Let me explain.

- In our example, we adopted the *HRW*-language (\mathcal{L}_1).
- Consider this *expressively equivalent HMA*-language (\mathcal{L}_2).
 - \mathcal{L}_2 has atoms H and M, A , where M, A are such that:
 - $M \models H \equiv R.$ [“Minnesotan” weather]
 - $A \models H \equiv W.$ [“Arizonan” weather]
- \mathcal{L}_1 and \mathcal{L}_2 can express exactly the same propositions. This is what it means for $\mathcal{L}_1, \mathcal{L}_2$ to be *expressively equivalent*.
- In \mathcal{L}_2 , the claims T, H_1 , and H_2 are expressed as follows:
 - $(T) H \& M \& A.$
 - $(H_1) \sim H \& \sim M \& \sim A.$
 - $(H_2) \sim H \& M \& A.$
- Thus, if we “count numbers of true conjuncts” — in \mathcal{L}_2 , then we get a *reversal* of the ordering we got in \mathcal{L}_1 . That is:
 - In \mathcal{L}_1 , we have $H_1 < H_2$.
 - In \mathcal{L}_2 , we have $H_2 < H_1$.

☞ This shows that verisimilitude — on this “counting true conjuncts” definition — is *language dependent* ([9], [10]).

- Let $\mathcal{P}(x) \stackrel{\text{def}}{=} \text{the set of predicates that an object } x \text{ falls under.}$
- We seek a measure $s(x, y)$ of the “degree to which x and y are similar”. A naive way to measure similarity is by “counting the predicates x and y share (both fall under)”.
- Here is the naive “shared predicate” similarity measure:

$$s(x, y) \stackrel{\text{def}}{=} |\mathcal{P}(x) \cap \mathcal{P}(y)|$$
- That is, $s(x, y)$ is just the size (cardinality) of the intersection of the sets of predicates that x and y (respectively) fall under. [Like “counting conjuncts” again!]
- Carnap *et. al.* ([3], [4]) say things which suggest the following “analogy by similarity” principle for inductive logic (where, it is not known whether a falls under the predicate ϕ):
 - (\mathcal{A}) If $n > m$, then $\Pr[\phi a \mid \phi b \& s(a, b) = n] > \Pr[\phi a \mid \phi b \& s(a, b) = m].$
- In words, (\mathcal{A}) says the following. Suppose that b falls under ϕ . Then, the probability that a also falls under ϕ *increases as the (known) similarity between a and b [$s(a, b)$] increases.*

- Unfortunately, (\mathcal{A}) leads to a *language dependent* theory of inductive probability. A Miller-style trick proves this.
 - The *ABCD*-language (\mathcal{L}_1) has four predicates *A, B, C, D*.
 - The *AXYD*-language (\mathcal{L}_2) also has four predicates *A, X, Y, D*, where *X* and *Y* are subject to the following Miller-esque semantic constraints: $Xx \models Ax \equiv Bx, Yx \models Bx \equiv Cx$.
 - Note: \mathcal{L}_1 and \mathcal{L}_2 are expressively equivalent!
 - Now, consider two objects *a* and *b* such that (speaking \mathcal{L}_1):

$$Aa \ \& \ Ba \ \& \ Ca$$

$$Ab \ \& \ \sim Bb \ \& \ Cb$$
 - In \mathcal{L}_2 , *a* and *b* have the following (equivalent) descriptions:

$$Aa \ \& \ Xa \ \& \ Ya$$

$$Ab \ \& \ \sim Xb \ \& \ \sim Yb$$
 - Now, we see that our measure s is *language dependent*:

$$s_{\mathcal{L}_1}(a, b) = 2 \neq 1 = s_{\mathcal{L}_2}(a, b)$$
- We can use this example to establish the language dependence of any probability theory that satisfies (\mathcal{A}).

- Recall the Carnapian principle of “analogy by similarity”:

$$(\mathcal{A}) \text{ If } n > m, \text{ then } \Pr[\phi a \mid \phi b \ \& \ s(a, b) = n] > \Pr[\phi a \mid \phi b \ \& \ s(a, b) = m].$$
- Applying this principle within *both* \mathcal{L}_1 and \mathcal{L}_2 yields:
 - $\Pr[Da \mid Db \ \& \ (Aa \ \& \ Ba \ \& \ Ca \ \& \ Ab \ \& \ \sim Bb \ \& \ Cb)] > \Pr[Da \mid Db \ \& \ (Aa \ \& \ Ba \ \& \ Ca \ \& \ Ab \ \& \ \sim Bb \ \& \ \sim Cb)]$
 - $\Pr[Da \mid Db \ \& \ (Aa \ \& \ Xa \ \& \ Ya \ \& \ \sim Ab \ \& \ Xb \ \& \ Yb)] > \Pr[Da \mid Db \ \& \ (Aa \ \& \ Xa \ \& \ Ya \ \& \ Ab \ \& \ \sim Xb \ \& \ \sim Yb)]$
 - It is useful here to adopt the following *abbreviations*:

$$x \stackrel{\text{def}}{=} Db \ \& \ (Aa \ \& \ Ba \ \& \ Ca \ \& \ Ab \ \& \ \sim Bb \ \& \ Cb) \quad [x \stackrel{\text{def}}{=} Db \ \& \ s_{\mathcal{L}_1}(a, b) = 2]$$

$$y \stackrel{\text{def}}{=} Db \ \& \ (Aa \ \& \ Xa \ \& \ Ya \ \& \ Ab \ \& \ \sim Xb \ \& \ \sim Yb) \quad [y \stackrel{\text{def}}{=} Db \ \& \ s_{\mathcal{L}_2}(a, b) = 1]$$

$$z \stackrel{\text{def}}{=} Db \ \& \ (Aa \ \& \ Xa \ \& \ Ya \ \& \ \sim Ab \ \& \ Xb \ \& \ Yb) \quad [z \stackrel{\text{def}}{=} Db \ \& \ s_{\mathcal{L}_2}(a, b) = 2]$$

$$u \stackrel{\text{def}}{=} Db \ \& \ (Aa \ \& \ Ba \ \& \ Ca \ \& \ Ab \ \& \ \sim Bb \ \& \ \sim Cb) \quad [u \stackrel{\text{def}}{=} Db \ \& \ s_{\mathcal{L}_1}(a, b) = 1]$$
 - Thus, we have $x \models y$ and $z \models u$, but (\mathcal{A}) implies *both*
 - $\Pr(Da \mid x) > \Pr(Da \mid u)$, and
 - $\Pr(Da \mid z) > \Pr(Da \mid y)$.
 - This violates the axioms of conditional probability.

☞ *Either* (\mathcal{A}) must go, or we need some *restriction on the choice of language* — in order to block inferring (1) and (2).

- Suppose we have a panel of three judges (J_1, J_2, J_3). This panel will vote on an *agenda*, which stems from:

Question. In the reunited Germany, should the German parliament and the seat of government move to Berlin or stay in Bonn?
- Suppose the panel votes on these two (atomic) *premises*:
 - $P \stackrel{\text{def}}{=} \text{the parliament should move.}$
 - $G \stackrel{\text{def}}{=} \text{the seat of government should move.}$
- There is also the following “*conclusion*” whose truth-value is *determined by* the truth-values of the premises:
 - $B \stackrel{\text{def}}{=} \text{both the parliament and the seat of government should move.}$
- Suppose the judges render the following judgments (votes):

	P?	G?	B?
J_1	yes	no	no
J_2	no	yes	no
J_3	yes	yes	yes

- For each judge, the conclusion column is *determined by* the premise columns (*i.e.*, we assume each judge is *consistent*).

- Example of *doctrinal paradox/discursive dilemma* ([6], [13]).

Doctrinal Paradox/Discursive Dilemma			
	P?	G?	B?
J_1	yes	no	no
J_2	no	yes	no
J_3	yes	yes	yes
Majority	yes	yes	yes & no?

- ☞ *Naive majority rule* for aggregating *all* judgments can lead to *inconsistent aggregations of premises + conclusions*.
- Various alternative aggregation procedures have been proposed, so as to ensure overall consistency. Example:
 - Premise-Based Procedure.** Use majority rule on the premises, and then just *enforce the logical conclusion*.
- The premise-based procedure seems reasonable (esp. if the premises make up the agenda that is explicitly voted on).

Premise-Based Procedure

	$P?$	$G?$	$B?$
J_1	yes	no	no
J_2	no	yes	no
J_3	yes	yes	yes
Majority:	yes	yes	(ignore)
Logical Conclusion:	yes		

- So, the premise-based procedure is a way of restoring logical consistency to a naive majority aggregation rule.
- But, the premise-based procedure (and many other procedures) faces another problem: *language dependence*.
- To see this, look at what happens when we move from the PG -language (\mathcal{L}_1) to the following PY -language (\mathcal{L}_2), where Y satisfies the following (Miller-style) semantic constraint:
 - $Y \models P \equiv G$.

Language Dependence of the Premise-Based Procedure

$\mathcal{L}_1:$	P	G	$P \equiv G$	$P \& G$
$\mathcal{L}_2:$	P	$P \equiv Y$	Y	$P \& Y$
J_1	yes	no	no	no
J_2	no	yes	no	no
J_3	yes	yes	yes	yes
Majority in $\mathcal{L}_1:$	yes	yes		yes
Majority in $\mathcal{L}_2:$	yes		no	no

- ☞ The aggregate conclusion judgment depends on which (expressively equivalent) language we use to express the premises (*i.e.*, the agenda).
 - This shows that the premise-based procedure is language dependent, in basically the same sense that our accounts of verisimilitude and inductive probability (above) were.
 - Which (consistent) procedures are language independent?

- A recent paper by Cariani (one of my former students, now at NU) *et. al.* gives a precise answer to this question [2].
- To understand their central theorem, I'll need to introduce a little more terminology. Two properties of procedures:
 - **Decisiveness.** An aggregation procedure is *decisive* iff it is *consistent and complete (and, hence, deductively closed)*.
 - **Anonymity.** An aggregation procedure is *anonymous* iff it is invariant under permutations of the (names of the) judges.
- Consistency and closure are basic logical constraints. Completeness is stronger than closure (given consistency).
- Anonymity basically says that *it shouldn't matter who the judges are (i.e., which judgment profiles belong to which judges)*. This is a substantive (non-logical) assumption. But, it's often plausible (*e.g.*, in our German Capital case).
- I will return to the completeness assumption, below.
- But, first, here is their new impossibility theorem [2].

Central Impossibility Theorem of Cariani *et. al.* [2]

Let n be the number of judges on the panel, and k be the number of atomic sentences in our language(s) \mathcal{L} . Provided that $k \geq \log_2(n + 2)$, *no* decisive aggregation procedure can be *both*:

- (i) anonymous,
- and*
- (ii) language independent.

- ☞ In other words, for languages with sufficiently many atomic sentences, *it is impossible for any aggregation procedure to be decisive, anonymous, and language independent*.
 - In fact, all decisive, language independent aggregation procedures must be *rolling dictatorships* [2].
 - \therefore *Most of the judgment aggregation procedures that have been proposed in the literature are language dependent.*
 - There are various ways to avoid this. *E.g.*, *incompleteness*.

- Incompleteness may make sense in some contexts [7]:
 If a group is prepared to refrain from making a collective judgment on some propositions ... then it may use an aggregation procedure such as the ‘unanimity procedure’, whereby the group makes a judgment on a proposition iff the group members unanimously endorse that judgment. Propositions judged to be true by all members are collectively judged to be true; and ones judged to be false by all members are collectively judged to be false; *no collective judgment is made on any other propositions.*
- Other incomplete/ \mathcal{L} -independent procedures (see [8] for survey):
 - A **conclusion-based** procedure, which returns no aggregate judgment for the premises. In our example, it returns “no” for the conclusion and it returns *nothing* for the premises.
 - **Approval voting** [1], which returns as the group judgment set the most popular of the voters’ judgment sets (where judges can vote on *any number* of agenda items that they like). In our example, it returns *all* the judgment sets (that is to say, *no definitive aggregate judgment whatsoever*).

- So far, we’ve seen “logical”/“syntactical” manifestations of LD. But, the underlying phenomenon is *much more general*.
- As Popper [14, Appendix 2] and Miller [10, Ch. 11] show, the problem also arises in the context of *quantitative* theories.
- For instance, suppose we have two (real-valued) parameters ϕ and ψ , and we have two hypotheses H_1 and H_2 , each of which implies estimates/predictions of their values. To wit:

	ϕ	ψ	α	β
H_1	0.150	1.225	0.925	2.000
H_2	0.100	1.000	0.800	1.700
T	0.000	1.000	1.000	2.000

- Here, H_2 is “closer to the truth” (T) *with respect to both* ϕ and ψ . But, this *reverses with respect to* α and β , where:

$$\alpha = \psi - 2\phi \quad \beta = 2\psi - 3\phi$$

$$\phi = \beta - 2\alpha \quad \psi = 2\beta - 3\alpha$$
- ☞ The ϕ/ψ and α/β languages are *expressively equivalent*. But, assessments of “closeness to T ” *depend on language*.

- Jim Joyce [5] has recently argued for probabilism, by appeal to the following sort of “verisimilitude theorem”.
 - **Theorem** (Joyce). If S ’s credence function b (construed as providing *estimates of truth-values*) fails to be probabilistic, then there exists a *probabilistic* estimate b' of truth-values that is *closer to the truth than* b , *in all possible worlds*.
- As it turns out, one can also prove the following (*via* translations slightly more complex than Popper-Miller’s).
 - **Theorem** (Me). For any pair of credence functions $\langle b, b' \rangle$ of the sort mentioned in Joyce’s theorem (and for any way of measuring “closeness” you like), there exists a *translation* τ such that $\tau(b)$ is *closer to the truth than* $\tau(b')$ — *with respect to the “ τ -truth-values” — in all possible worlds*.

- ☞ There are *some aspects of* “the truth” with respect to which (any) incoherent b is bound to be less accurate than (some) coherent b' , but — for any such pair $\langle b, b' \rangle$ — there are also *some aspects of* “the truth” on which *the opposite is the case*.
- Why can we *ignore some aspects of* the truth, but not others?

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