#### Means-end Relations and a Measure of Efficacy

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July 4, 2005

## Outline

#### Means-end relations

- Interest I: Practical syllogisms
- Interest II: Functional ascriptions
- Propositional Dynamic Logic

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#### 2 Efficacy via fuzzy logic

- Reliability as a fuzzy operator
- The resulting fuzzy logic

Means-end relations Efficacy via fuzzy logic Propositional Dynamic Logic

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This premise is a means-end relation.

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

### An example from von Wright



I want to make the hut habitable.

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I want to make the hut habitable. Unless I heat the hut, it will not be habitable.

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

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Therefore I must heat the hut.

Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

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Note: distinct premises

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- Expression of an agent's desire, Note: distinct premises
- A necessary means-end relation,<sup>4</sup>
- Concludes in a *necessary* action.

But necessary means-end relations are a bit tricky.

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

### An example from von Wright



I want to make the hut habitable. If I heat the hut, it will be habitable. Therefore, I have reason to heat the hut.

An alternative with a sufficient means-end relation.



• "The function of the heart is to pump blood."

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- "The function of the heart is to pump blood."
- "That switch mutes the television."

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We ascribe functions to biological stuff, artifacts, algorithms, personal roles...

Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

#### How functions relate to means and ends



"That switch mutes the television."

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#### How functions relate to means and ends



"That switch mutes the television." ↓ One can *use* the switch to mute the television.

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Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

#### How functions relate to means and ends



"That switch mutes the television."
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One can use the switch to mute the television.
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  - distinct of desirability
  - distinct from theory of practical reasoning

Interest I: Practical syllogisms Interest II: Functional ascriptions Propositional Dynamic Logic

## Initial analysis of means-end relations

• An end is some desirable condition – a proposition.

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- Ends-in-themselves?
- Objects as means?

## PDL syntax

Propositional Dynamic Logic is a logic of actions.

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Basic types: • a set **act** of *actions*,

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#### Intuitions:

•  $[\alpha]\varphi$ : after doing  $\alpha$ ,  $\varphi$  will hold.

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Basic types:

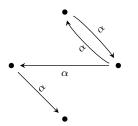
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Intuitions:

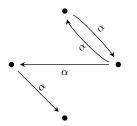
- $[\alpha]\varphi$ : after doing  $\alpha$ ,  $\varphi$  will hold.
- $\langle \alpha \rangle \varphi$ : after doing  $\alpha$ ,  $\varphi$  might hold.

#### PDL semantics



Possible world semantics with transition systems for each action  $\alpha$ .

## PDL semantics

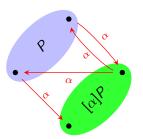


Possible world semantics with transition systems for each action  $\alpha$ .

 $w \xrightarrow{\alpha} w'$  means:

one can reach w' by doing  $\alpha$  in w.

#### PDL semantics



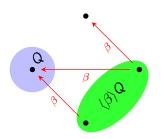
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 $\mathbf{w} \models [\alpha] \varphi \quad iff \quad \forall \ \mathbf{w} \xrightarrow{\alpha} \mathbf{w}' \quad \mathbf{w}' \models \varphi.$ 

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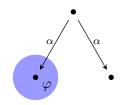
one can reach w' by doing  $\alpha$  in w.

$$w \models [\alpha]\varphi \quad iff \ \forall \ w \xrightarrow{\alpha} w' \ . \ w' \models \varphi.$$
$$w \models \langle \alpha \rangle \varphi \quad iff \ \exists \ w \xrightarrow{\alpha} w' \ . \ w' \models \varphi.$$

A means is an action  $\alpha$  that can realize one's end  $\varphi.$ 

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Two interpretations:



Weak:  $\alpha$  might realize  $\varphi$ .

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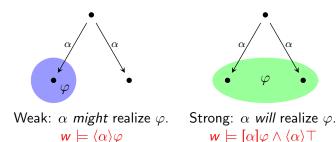


Weak:  $\alpha$  might realize  $\varphi$ . Strong:  $\alpha$  will realize  $\varphi$ .

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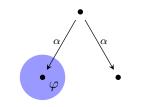
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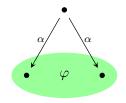
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A means is an action  $\alpha$  that can realize one's end  $\varphi$ .

Two interpretations:



Weak:  $\alpha$  might realize  $\varphi$ .  $w \models \langle \alpha \rangle \varphi$ 



Strong:  $\alpha$  will realize  $\varphi$ .  $w \models [\alpha] \varphi \land \underline{\langle \alpha \rangle} \top$  $\alpha$  can be done.

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End: Get 12 points with one dart.



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Three different means:

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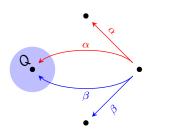
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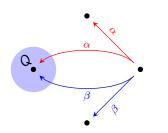


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Efficacy: The degree of reliability of a means to an end.



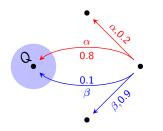
Efficacy is a measure of likelihoods.



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PDL includes non-determinism, not probabilities.

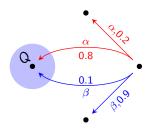


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*Fix (semantic):* use *probabilistic* transition structures.

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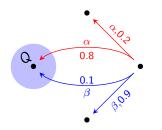
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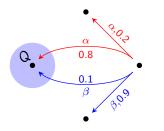
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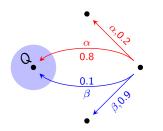
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<u>Write</u>:  $P(w \xrightarrow{\alpha} w') = x$ .



#### Syntactic fix?

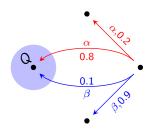
Hughes, Esterline, Kimiaghalam Means-end Relations and a Measure of Efficacy



Syntactic fix?

• Probabilistic Computation Tree Logic (pCTL)?

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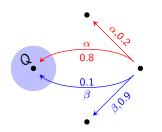


Syntactic fix?

- Probabilistic Computation Tree Logic (pCTL)?
  - Index dynamic operators, like  $[\alpha]$   $\langle \alpha \rangle$

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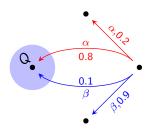
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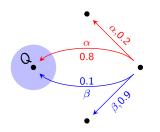
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• Probabilistic PDL?

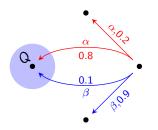


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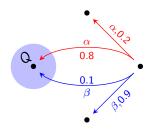
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  - Assigns values in [0, 1] to world-formula pairs.

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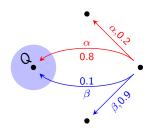


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### • Fuzzy PDL.

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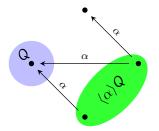
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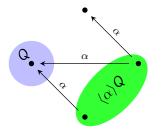
Hajek, et al., uses distributions on propositional formulas to define "Probably  $\varphi$  ".

Truth degree of "Probably  $\varphi$ " =  $P(\varphi)$ .

"Reliably", like "Probably", is a vague operator.

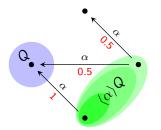


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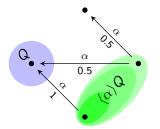
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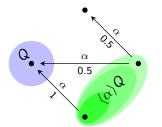
In PDL:  $\langle \alpha \rangle \varphi \Leftrightarrow \alpha \text{ will possibly realize } \varphi$ In <u>fuzzy</u> PDL:  $\langle \alpha \rangle \varphi \Leftrightarrow \alpha \text{ will probably realize } \varphi$ 

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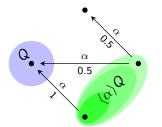
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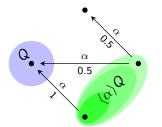
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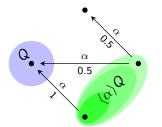
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• Like decision theory, we use means for expected outcomes.

- Unlike decision theory, there are no utilities involved.
- Elegant treatment of complex ends, like  $\langle \alpha \rangle \varphi \wedge \langle \beta \rangle \psi$ .

### Fuzzy ends An accidental advantage

### Weapons are for causing harm.

Means-end relations Efficacy via fuzzy logic Reliability as a fuzzy operator The resulting fuzzy logic

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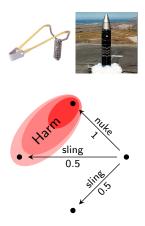
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Means-end relations Efficacy via fuzzy logic Reliability as a fuzzy operator The resulting fuzzy logic

### Fuzzy ends An accidental advantage



Weapons are for causing harm. Examples: slingshot, nuke This end is fuzzy. Fuzzy PDL allows for fuzzy ends. A nuke is more effective in causing harm than a slingshot. (Duh.)

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#### Reliability as a fuzzy operator The resulting fuzzy logic

## Extending the logic to other connectives

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### The resulting fuzzy logic

## Extending the logic to other connectives

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 $\llbracket \varphi \land \psi \rrbracket(w) = \min \{\llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \}$ 

### On formulas

$$\begin{split} \llbracket \langle \alpha \rangle \varphi \rrbracket(w) &= \sum_{w' \in \mathcal{W}} P(w \xrightarrow{\alpha} w') \cdot \llbracket \varphi \rrbracket(w') \\ \llbracket \varphi \land \psi \rrbracket(w) &= \min \{ \llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \} \\ \llbracket \varphi \lor \psi \rrbracket(w) &= \max \{ \llbracket \varphi \rrbracket(w), \llbracket \psi \rrbracket(w) \} \\ \llbracket \neg \varphi \rrbracket(w) &= 1 - \llbracket \varphi \rrbracket(w) \\ \llbracket \varphi \to \psi \rrbracket(w) &= \begin{cases} 1 & \text{if } \llbracket \varphi \rrbracket(w) \le \llbracket \psi \rrbracket(w), \\ \llbracket \psi \rrbracket(w) &= \text{else;} \end{cases} \end{split}$$

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 $\llbracket \alpha; \beta \rrbracket(w)(w') = \sum_{w'' \in \mathcal{W}} P(w \xrightarrow{\alpha} w'') \cdot P(w'' \xrightarrow{\beta} w')$  $\llbracket \varphi? \rrbracket(w)(w') = \begin{cases} \llbracket \varphi \rrbracket(w) & \text{if } w = w'; \\ 0 & \text{else.} \end{cases}$ 

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# The semantics of fuzzy PDL

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 $\left[ \begin{array}{c} [\varphi \cup \psi](w)(w') \\ [\varphi^*](w)(w') \end{array} \right\} \text{ undefined.}$ 

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  - Axioms:
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  - Test:  $[\psi?]\varphi \leftrightarrow (\psi \rightarrow \varphi)$

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#### Logical properties Completeness

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# I wish.

Hughes, Esterline, Kimiaghalam Means-end Relations and a Measure of Efficacy

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#### Logical properties Completeness

# I wish.

But not with these semantics.

Ongoing work...

• Include non-deterministic features (in paper).

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Thank you.

Concerns:

• Primary: Adding probabilities to transitions.

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Proposal: Interpret PDL as fuzzy logic.

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