# ON A HOMOMORPHISM PROPERTY OF HOOPS 

R. VEROFF AND M. SPINKS

Abstract. We present a syntactic proof that equation

$$
e \rightarrow(b \cdot c)=(e \rightarrow b) \cdot(e \rightarrow c)
$$

is satisfied in a hoop $\mathbf{A}$ for any idempotent $e \in A$ and all $b, c \in A$. The theorem both answers a question and generalizes a result of Ferreirim [6].

## 1. Introduction

A hoop is an algebra $\langle A ; \cdot, \rightarrow, 1\rangle$ of type $\langle 2,2,0\rangle$ that satisfies the identities:

$$
\begin{gather*}
x \cdot(y \cdot z) \approx(x \cdot y) \cdot z  \tag{M1}\\
x \cdot y \approx y \cdot x  \tag{M2}\\
x \cdot \mathbf{1} \approx x  \tag{M3}\\
x \rightarrow x \approx \mathbf{1}  \tag{M4}\\
(x \rightarrow y) \cdot x \approx(y \rightarrow x) \cdot y  \tag{M5}\\
(x \cdot y) \rightarrow z \approx x \rightarrow(y \rightarrow z) . \tag{M6}
\end{gather*}
$$

We denote the variety of hoops by HO .
Hoops were first investigated by Büchi and Owens in an unpublished manuscript [5] of 1975, and they have since been studied by Ferreirim [6, 7], Blok and Ferreirim [2, 3], Aglianó and Panti [1] and Blok and Pigozzi [4] among others. The study of hoops is motivated by their occurrence both in universal algebra and algebraic logic. Typical examples of hoops include both Brouwerian semilattices and the positive cones of lattice ordered Abelian groups, while hoops structurally enriched with normal multiplicative operators naturally generalize the

[^0]normal Boolean algebras with operators of Jónsson and Tarski $[8,9]$. For details, see in particular Blok and Pigozzi [4].

For a hoop $\mathbf{A}$, an idempotent is an element $e \in A$ with the property that $e \cdot e=e$. It is well known and easy to see [4, Lemma 1.9(v)] that for any hoop $\mathbf{A}$ and idempotent $e \in A$, the following equation is satisfied for all $b, c \in A$ :

$$
\begin{equation*}
e \rightarrow(b \rightarrow c)=(e \rightarrow b) \rightarrow(e \rightarrow c) . \tag{1}
\end{equation*}
$$

For any $\langle\rightarrow, \cdot, \mathbf{1}\rangle$-term $t:=t(\vec{x})$, the $k$-th iterated power $t^{k}, 0<k<$ $\omega$, is defined recursively by:

$$
\begin{gathered}
t^{0}:=\mathbf{1} \\
t^{k}:=t \cdot t^{k-1}
\end{gathered}
$$

A hoop is said to be $k$-potent if it satisfies the identity:

$$
x^{k} \approx x^{k-1}
$$

In her Ph.D. thesis [6, Chapter 3, Lemma 1.10], Ferreirim shows that in addition to (1), any $k$-potent hoop $\mathbf{A}$ satisfies the following equation for any idempotent $e \in A$ and all $b, c \in A$ :

$$
\begin{equation*}
e \rightarrow(b \cdot c)=(e \rightarrow b) \cdot(e \rightarrow c) \tag{2}
\end{equation*}
$$

Together, (1) and (2) assert that the map $a \mapsto(e \rightarrow a)$ is an endomorphism for any $k$-potent hoop $\mathbf{A}$ and fixed idempotent $e \in A$. This implies in particular that $\mathbf{H}(\mathbf{A}) \subseteq \mathbf{I S}(\mathbf{A})$ for any finite hoop $\mathbf{A}$ [6, Chapter 3, Lemma 1.11]. This strong property plays an important role in the proofs of several results in the theory of hoops, including Ferreirim's characterization of finitely based varieties of $k$-potent hoops [6, Chapter 3, Theorem 1.13], and her characterization of varieties of hoops in which every subquasivariety is itself a variety [6, Chapter 3, Theorem 2.13].

The proof of (2) given by Ferreirim in [6, Chapter 3, Lemma 1.10] relies on a sophisticated model-theoretic argument that exploits her characterization of the subdirectly irreducible $k$-potent hoops [6, Chapter 2, Theorem 3.12]. Immediately following her proof of (2) Ferreirim remarks [6, p. 58]: 'A syntactic proof of statement (2) ( $e \rightarrow b$ ) • $(e \rightarrow$ $c)=e \rightarrow(b \cdot c)$ in Lemma 1.10 would certainly be more elegant. We couldn't find one and propose it as an open problem.'

In this note, we present a solution to Ferreirim's problem by exhibiting a syntactic proof of her equation (2). Moreover, our solution generalizes Ferreirim's result to all hoops, since our proof does not assume $k$-potency. In particular, our proof holds for the subvariety $L$ of Eukasiewicz hoops, namely the class of all hoops satisfying the commutative identity $(x \rightarrow y) \rightarrow y \approx(y \rightarrow x) \rightarrow x$; the significance of L in the theory of HO lies in a result due to Ferreirim [6, Chapter 3, Corollary 3.4] that shows Eukasiewicz hoops are, in a precise technical sense, the building blocks of arbitrary hoops. We note that Ferreirim's proof of (2) does not extend to L, since the variety of Łukasiewicz hoops is not $k$-potent for any $k<\omega$ by [6, Chapter 2, Corollary 4.17] and [4, Corollary 5.5].

## 2. The Proof

In the following (machine-oriented) proof, the justification $[i \rightarrow j]$ indicates paramodulation from $i$ into $j$, that is, unifying the left-hand side of $i$ with a subterm of $j$, instantiating $j$ with the corresponding substitution, and replacing the subterm with the corresponding instance of the right-hand side of $i$.

| 1. $x \cdot(y \cdot z)=(x \cdot y) \cdot z$ | $[\mathrm{M} 1]$ |
| :--- | ---: |
| 2. $x \cdot y=y \cdot x$ | $[\mathrm{M} 2]$ |
| 3. $x \cdot 1=x$ | $[\mathrm{M} 3]$ |
| 4. $x \rightarrow x=1$ | $[\mathrm{M} 4]$ |
| 5. $(x \rightarrow y) \cdot x=(y \rightarrow x) \cdot y$ | $[\mathrm{M} 5]$ |
| 6. $(x \cdot y) \rightarrow z=x \rightarrow(y \rightarrow z)$ | $[\mathrm{M} 6]$ |
| 7. $e \cdot e=e$ | $[1 \rightarrow 2]$ |
| 8. $x \cdot(y \cdot z)=z \cdot(x \cdot y)$ | $[1 \rightarrow 2]$ |
| 9. $(x \cdot y) \cdot z=(y \cdot z) \cdot x$ | $[2 \rightarrow 1]$ |
| 10. $(x \cdot y) \cdot z=x \cdot(z \cdot y)$ | $[2 \rightarrow 3]$ |
| 11. $1 \cdot x=x$ | $[4 \rightarrow 4]$ |
| 12. $x \rightarrow x=y \rightarrow y$ | $[4 \rightarrow 3]$ |
| 13. $x \cdot(y \rightarrow y)=x$ | $[2 \rightarrow 5]$ |
| 14. $(x \rightarrow y) \cdot x=y \cdot(y \rightarrow x)$ | $[5 \rightarrow 6]$ |
| 15. $((x \rightarrow y) \cdot x) \rightarrow z=(y \rightarrow x) \rightarrow(y \rightarrow z)$ | $[2 \rightarrow 6]$ |
| 16. $(x \cdot y) \rightarrow z=y \rightarrow(x \rightarrow z)$ | $[7 \rightarrow 6]$ |
| 17. $e \rightarrow(e \rightarrow x)=e \rightarrow x$ | $[7 \rightarrow 1]$ |
| 18. $e \cdot(e \cdot x)=e \cdot x$ | $[6 \rightarrow 12]$ |
| 19. $((x \rightarrow y) \cdot x) \rightarrow y=z \rightarrow z$ |  |


| 20. $(x \cdot y) \rightarrow y=x \rightarrow(z \rightarrow z)$ | $[12 \rightarrow 6]$ |
| :---: | :---: |
| 21. $x \cdot(y \rightarrow(z \rightarrow(y \cdot z)))=x$ | $[6 \rightarrow 13]$ |
| 22. $(x \cdot y) \cdot z=y \cdot(z \cdot x)$ | $[2 \rightarrow 8]$ |
| 23. $x \cdot(x \rightarrow y)=y \cdot(y \rightarrow x)$ | [2 $\rightarrow$ 14] |
| 24. $(x \cdot(x \rightarrow y)) \rightarrow z=(y \rightarrow x) \rightarrow(y \rightarrow z)$ | [14 $\rightarrow 6$ ] |
| 25. $e \cdot((x \rightarrow e) \cdot x)=e \cdot(e \rightarrow x)$ | $[14 \rightarrow 18]$ |
| 26. $(x \cdot(x \rightarrow y)) \rightarrow z=y \rightarrow((y \rightarrow x) \rightarrow z)$ | [14 $\rightarrow$ 16] |
| 27. $x \rightarrow(y \rightarrow z)=y \rightarrow(x \rightarrow z)$ | $[6 \rightarrow 16]$ |
| 28. $x \cdot(y \rightarrow(z \rightarrow(z \cdot y)))=x$ | [16 $\rightarrow$ 13] |
| 29. $((x \rightarrow y) \cdot x) \rightarrow x=z \rightarrow z$ | [5 $\rightarrow$ 19] |
| 30. $x \rightarrow((x \rightarrow y) \rightarrow y)=z \rightarrow z$ | [16 $\rightarrow$ 19] |
| 31. $x \rightarrow(y \rightarrow y)=z \rightarrow(x \rightarrow z)$ | $[20 \rightarrow 16]$ |
| 32. $(x \rightarrow(y \cdot x)) \cdot((x \rightarrow(y \cdot x)) \rightarrow y)=y$ | $[21 \rightarrow 23]$ |
| 33. $e \rightarrow(x \rightarrow(e \rightarrow y))=x \rightarrow(e \rightarrow y)$ | $[17 \rightarrow 27]$ |
| 34. $(x \rightarrow y) \cdot((x \rightarrow y) \rightarrow z)=z \cdot(x \rightarrow(z \rightarrow y))$ | $[27 \rightarrow 23]$ |
| 35. $x \cdot(y \rightarrow((y \rightarrow z) \rightarrow z))=x$ | $[27 \rightarrow 13]$ |
| 36. $(x \rightarrow y) \rightarrow(z \rightarrow z)=u \rightarrow u$ | [20 $\rightarrow 29$ ] |
| 37. $((x \rightarrow y) \rightarrow y) \cdot(((x \rightarrow y) \rightarrow y) \rightarrow x)=x$ | $[23 \rightarrow 35]$ |
| 38. $x \cdot(y \rightarrow(z \rightarrow z))=x$ | $[36 \rightarrow 35]$ |
| 39. $x \cdot(y \rightarrow(z \rightarrow y))=x$ | $[31 \rightarrow 38$ ] |
| 40. $x \cdot((y \cdot z) \rightarrow z)=x$ | $[20 \rightarrow 38]$ |
| 41. $x \cdot((y \cdot z) \rightarrow y)=x$ | [16 $\rightarrow 38$ ] |
| 42. $((x \rightarrow y) \rightarrow y) \cdot(x \rightarrow y)=y$ | $[14 \rightarrow 39]$ |
| 43. $x \rightarrow(y \rightarrow x)=1$ | $[11 \rightarrow 39]$ |
| 44. $(x \cdot y) \rightarrow x=1$ | $[6 \rightarrow 43]$ |
| 45. $(x \cdot(y \cdot z)) \rightarrow(z \cdot x)=1$ | $[22 \rightarrow 44]$ |
| 46. $x \cdot(((y \cdot z) \cdot u) \rightarrow y)=x$ | [9 $\rightarrow 40]$ |
| 47. $x \cdot(x \rightarrow(y \cdot x))=y \cdot x$ | $[23 \rightarrow 40]$ |
| 48. $x \cdot(x \rightarrow(x \cdot y))=x \cdot y$ | $[23 \rightarrow 41]$ |
| 49. $(x \rightarrow x) \rightarrow x=x$ | $[13 \rightarrow 42]$ |
| 50. $((x \rightarrow y) \rightarrow y) \rightarrow((x \rightarrow y) \rightarrow z)=y \rightarrow z$ | [42 $\rightarrow 6$ ] |
| 51. $(x \rightarrow((x \rightarrow y) \rightarrow y)) \rightarrow z=z$ | [30 $\rightarrow$ 49] |
| 52. $(x \cdot y) \rightarrow((z \rightarrow y) \cdot x)=1$ | $[42 \rightarrow 45]$ |
| 53. $(x \rightarrow y) \rightarrow(x \rightarrow z)=(y \rightarrow x) \rightarrow(y \rightarrow z)$ | $[6 \rightarrow 15]$ |
| 54. $x \cdot((y \cdot z) \rightarrow(u \rightarrow y))=x$ | $[6 \rightarrow 46]$ |
| 55. $x \cdot(x \rightarrow(y \cdot x))=x \cdot y$ | $[2 \rightarrow 47]$ |
| 56. $x \rightarrow((x \rightarrow(x \cdot y)) \rightarrow z)=(x \cdot y) \rightarrow z$ | $[48 \rightarrow 6]$ |
| 57. $e \rightarrow((x \rightarrow e) \cdot e)=1$ | $[7 \rightarrow 52]$ |
| 58. $e \cdot(x \rightarrow e)=e \cdot 1$ | [57 $\rightarrow$ 55] |


| $\begin{aligned} & \text { 59. } e \cdot(x \rightarrow e)=e \\ & \text { 60. } e \cdot((x \rightarrow e) \cdot y)=e \cdot y \end{aligned}$ | $\begin{aligned} & {[3 \rightarrow 58]} \\ & {[59 \rightarrow 1]} \end{aligned}$ |
| :---: | :---: |
| 61. $e \cdot(e \rightarrow x)=e \cdot x$ | $[60 \rightarrow 25]$ |
| 62. $(e \rightarrow x) \rightarrow(e \rightarrow y)=(e \cdot x) \rightarrow y$ | $[61 \rightarrow 16]$ |
| 63. $(e \rightarrow x) \rightarrow(e \rightarrow y)=x \rightarrow(e \rightarrow y)$ | $[16 \rightarrow 62]$ |
| 64. $(e \rightarrow x) \rightarrow(e \rightarrow y)=e \rightarrow(x \rightarrow y)$ | $[6 \rightarrow 62$ ] |
| 65. $(e \rightarrow x) \cdot((e \rightarrow x) \rightarrow(e \rightarrow y))$ | $[63 \rightarrow 23]$ |
| 6. $x \cdot((y \rightarrow(z \cdot u)) \rightarrow(y \rightarrow z))=x$ | $[24 \rightarrow 46]$ |
| 67. $(x \rightarrow y) \cdot((x \rightarrow y) \rightarrow(x \rightarrow(y \cdot z)))=x \rightarrow(y \cdot z)$ | $[23 \rightarrow 66]$ |
| 68. $(x \rightarrow(x \cdot y)) \rightarrow(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow z)=y \rightarrow z$ | $[28 \rightarrow 26]$ |
| 69. $(x \rightarrow(y \cdot x)) \cdot((x \rightarrow(x \cdot y)) \rightarrow y)$ | [2 $\rightarrow 32$ ] |
| 70. $(x \rightarrow((y \rightarrow z) \cdot x)) \cdot(y \rightarrow((x \rightarrow((y \rightarrow z) \cdot x)) \rightarrow z))$ |  |
| $=y \rightarrow z$ | $[27 \rightarrow 32]$ |
| 71. $(x \rightarrow y) \rightarrow((z \rightarrow x) \rightarrow(z \rightarrow y))=1$ | $[53 \rightarrow 43]$ |
| 72. $(((x \rightarrow y) \rightarrow y) \rightarrow z) \rightarrow(x \rightarrow z)=$ | $[51 \rightarrow 71]$ |
| 73. $(((x \rightarrow y) \rightarrow y) \rightarrow y) \cdot 1=x \rightarrow y$ | $[72 \rightarrow 37]$ |
| 74. $((x \rightarrow y) \rightarrow y) \rightarrow y=x \rightarrow y$ | $[3 \rightarrow 73]$ |
| 75. $x \rightarrow((x \rightarrow(x \cdot y)) \rightarrow z)=x \rightarrow(y \rightarrow z)$ | $[6 \rightarrow 56]$ |
| 76. $(x \rightarrow y) \cdot(x \rightarrow((x \rightarrow y) \rightarrow(y \cdot z)))=x \rightarrow(y \cdot z)$ | [66 $\rightarrow 34$ ] |
| 77. $(x \rightarrow y) \cdot((((x \rightarrow y) \rightarrow y) \rightarrow y) \rightarrow(((x \rightarrow y) \rightarrow y) \rightarrow(y \cdot z)))$ |  |
| $=((x \rightarrow y) \rightarrow y) \rightarrow(y \cdot z)$ | [74 $\rightarrow 67$ ] |
| 78. $((x \rightarrow(x \cdot y)) \rightarrow y) \cdot(y \rightarrow(y \cdot z))$ |  |
| $(x \rightarrow(x \cdot y)) \rightarrow(y \cdot z)$ | $[68 \rightarrow 76]$ |
| 79. $(e \rightarrow x) \cdot(e \rightarrow(x \rightarrow y))=(e \rightarrow y) \cdot(y \rightarrow(e \rightarrow x))$ | $[64 \rightarrow 65]$ |
| 80. $(e \rightarrow(x \cdot e)) \cdot((x \cdot e) \rightarrow(e \rightarrow x))=e \rightarrow x$ | $[28 \rightarrow 79]$ |
| 81. $e \rightarrow(x \cdot e)=e \rightarrow x$ | $[54 \rightarrow 80]$ |
| 82. $e \rightarrow(x \cdot(e \cdot y))=e \rightarrow(y \cdot x)$ | $[22 \rightarrow 81]$ |
| 83. $e \rightarrow(x \cdot(e \cdot y))=e \rightarrow(x \cdot y)$ | [10 $\rightarrow 81$ ] |
| 84. $e \rightarrow(x \cdot(e \cdot y))=e \rightarrow(x \cdot(e \rightarrow y))$ | $[61 \rightarrow 83]$ |
| 85. $e \rightarrow(x \cdot(e \rightarrow y))=e \rightarrow(y \cdot x)$ | $[82 \rightarrow 84]$ |
| 86. $e \rightarrow((e \rightarrow x) \cdot y)=e \rightarrow(x \cdot y)$ | $[2 \rightarrow 85]$ |
| 87. $(e \rightarrow x) \rightarrow(e \rightarrow(y \cdot z))=e \rightarrow(x \rightarrow((e \rightarrow y) \cdot z))$ | $[86 \rightarrow 64]$ |
| 88. $e \rightarrow(x \rightarrow((e \rightarrow y) \cdot z))=e \rightarrow(x \rightarrow(y \cdot z))$ | [64 $\rightarrow 87$ ] |
| 89. $(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y)$ |  |
| $\cdot(x \rightarrow((((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow((x \rightarrow(y \cdot x))$ |  |
| $=x \rightarrow(y \cdot x)$ | $[69 \rightarrow 70]$ |
| 90. $(x \rightarrow y) \cdot(y \rightarrow(y \cdot z))=((x \rightarrow y) \rightarrow y) \rightarrow(y \cdot z)$ | [50 $\rightarrow 77$ ] |

$$
\begin{aligned}
& \text { 91. }(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y) \rightarrow(y \cdot z) \\
& =(x \rightarrow(x \cdot y)) \rightarrow(y \cdot z) \\
& {[78 \rightarrow 90]} \\
& \text { 92. }(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y) \\
& \cdot(x \rightarrow((((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y) \rightarrow(y \cdot x))) \\
& =x \rightarrow(y \cdot x) \\
& {[69 \rightarrow 89]} \\
& \text { 93. }(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y) \cdot(x \rightarrow((x \rightarrow(x \cdot y)) \rightarrow(y \cdot x))) \\
& =x \rightarrow(y \cdot x) \\
& {[91 \rightarrow 92]} \\
& \text { 94. }(((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y) \cdot(x \rightarrow(y \rightarrow(y \cdot x))) \\
& =x \rightarrow(y \cdot x) \\
& {[75 \rightarrow 93]} \\
& \text { 95. }((x \rightarrow(x \cdot y)) \rightarrow y) \rightarrow y=x \rightarrow(y \cdot x) \quad[28 \rightarrow 94] \\
& \text { 96. }((x \rightarrow(y \cdot x)) \rightarrow y) \rightarrow y=x \rightarrow(y \cdot x) \quad[2 \rightarrow 95] \\
& \text { 97. } e \rightarrow(((x \rightarrow((e \rightarrow y) \cdot x)) \rightarrow(e \rightarrow y)) \rightarrow(e \rightarrow y)) \\
& =x \rightarrow((e \rightarrow y) \cdot x) \\
& {[33 \rightarrow 96 \text { ] }} \\
& \text { 98. } e \rightarrow(x \rightarrow((e \rightarrow y) \cdot x))=x \rightarrow((e \rightarrow y) \cdot x) \quad[96 \rightarrow 97] \\
& \text { 99. } x \rightarrow((e \rightarrow y) \cdot x)=e \rightarrow(x \rightarrow(y \cdot x)) \quad[98 \rightarrow 88] \\
& \text { 100. } x \rightarrow((e \rightarrow y) \cdot x)=x \rightarrow(e \rightarrow(y \cdot x)) \quad[27 \rightarrow 99] \\
& \text { 101. } x \cdot(x \rightarrow(e \rightarrow(y \cdot x)))=(e \rightarrow y) \cdot x \quad[100 \rightarrow 47] \\
& \text { 102. }(e \rightarrow x) \cdot((e \rightarrow x) \rightarrow(e \rightarrow(x \cdot y))) \\
& =(e \rightarrow y) \cdot(e \rightarrow x) \quad[85 \rightarrow 101] \\
& \text { 103. }(e \rightarrow x) \cdot(e \rightarrow y)=e \rightarrow(y \cdot x) \quad[67 \rightarrow 102] \\
& \text { 104. }(e \rightarrow x) \cdot(e \rightarrow y)=e \rightarrow(x \cdot y) \quad[2 \rightarrow 103]
\end{aligned}
$$

This completes the proof.
Note. This proof was obtained with the assistance of the automated reasoning program Otter [10], using the method of proof sketches [12]. See [11] for examples of the application of automated reasoning to a wide range of problems in equational logic.

## References

[1] P. Aglianó and G. Panti, Geometrical methods in Wajsberg hoops, J. Algebra 256 (2002), 352-374.
[2] W. J. Blok and I. M. A. Ferreirim, Hoops and their implicational reducts (abstract), Algebraic Methods in Logic and Computer Science (C. Rauszer, ed.), Banach Centre Publications, vol. 28, Institute of Mathematics, Polish Academy of Sciences, Warszawa, 1993, pp. 219-230.
[3] _ On the structure of hoops, Algebra Universalis 43 (2000), 233-257.
[4] W. J. Blok and D. Pigozzi, On the structure of varieties with equationally definable principal congruences III, Algebra Universalis 32 (1994), 545-608.
[5] J. R. Büchi and T. M. Owens, Complemented monoids and hoops, Unpublished manuscript, 1975.
[6] I. M. A. Ferreirim, On varieties and quasivarieties of hoops and their reducts, Ph.D. thesis, University of Illinois at Chicago, 1992.
[7] __ On a conjecture by Andrzej Wroński for BCK-algebras and subreducts of hoops, Sci. Math. Japon. (2001), 119-132.
[8] B. Jónsson and A. Tarski, Boolean algebras with operators, Part I, Amer. J. Math. 73 (1951), 891-939.
[9] , Boolean algebras with operators, Part II, Amer. J. Math. 74 (1952), 127-162.
[10] W. McCune, Otter 3.3 Reference Manual, Tech. Memo ANL/MCS-TM-263, Mathematics and Computer Science Division, Argonne National Laboratory, Argonne, IL, August 2003.
[11] W. McCune and R. Padmanabhan, Automated deduction in equational logic and cubic curves, Lecture Notes in Computer Science (AI subseries), vol. 1095, Springer-Verlag, Berlin, 1996.
[12] R. Veroff, Solving open questions and other challenge problems using proof sketches, J. Automated Reasoning 27 (2001), no. 2, 157-174.

## Acknowledgements

This paper was written while the second author was a visitor at the School of Informatics and Engineering at the Flinders University of South Australia. The facilities and assistance provided by the University and the School are gratefully acknowledged.

## Authors’ Addresses

Robert Veroff<br>Department of Computer Science<br>University of New Mexico<br>Albuquerque, NM 87131<br>veroff@cs.unm.edu<br>Matthew Spinks<br>School of Informatics and Engineering<br>Flinders University of South Australia<br>mspinks@infoeng.flinders.edu.au


[^0]:    Date: January 19, 2004.

